Formal Semantics

Natural Language Processing CS 6120—Spring 2014 Northeastern University

David Smith some slides from Jason Eisner

Language as Structure

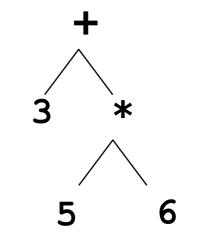
- So far, we've talked about structure
- What structures are more probable?
 - Language modeling: Good sequences of words/ characters
 - Text classification: Good sequences in defined contexts
- How can we recover hidden structure?
 - Tagging: hidden word classes
 - Parsing: hidden word relations

- Studying phonology, morphology, syntax, etc. independent of meaning is methodologically very useful
- We can study the structure of languages we don't understand
- We can use HMMs and CFGs to study protein structure and music, which don't bear meaning in the same way as language

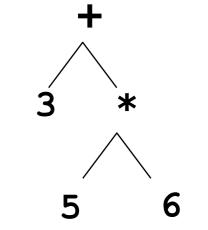
- How would you know if a computer "understood" the "meaning" of an (English) utterance (even in some weak "scarequoted" way)?
- How would you know if a person understood the meaning of an utterance?

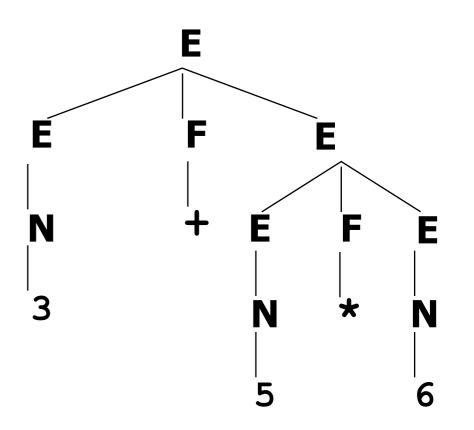
- Paraphrase, "state in your own words" (English to English translation)
- Translation into another language
- Reading comprehension questions
- Drawing appropriate inferences
- Carrying out appropriate actions
- Open-ended dialogue (Turing test)

- What is meaning of 3+5*6?
- First parse it into 3+(5*6)

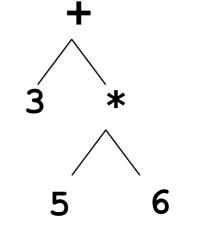


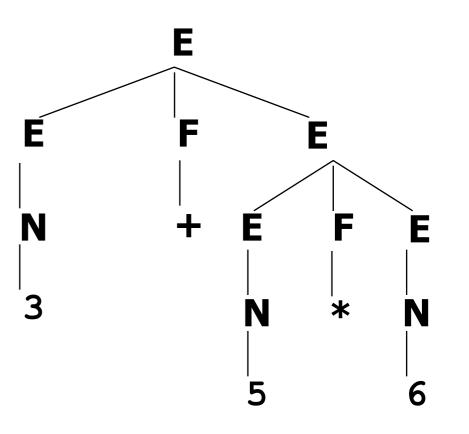
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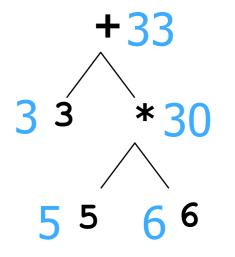


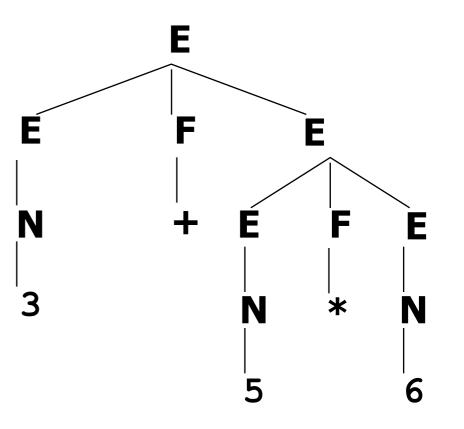
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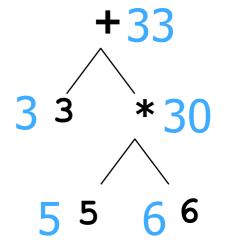


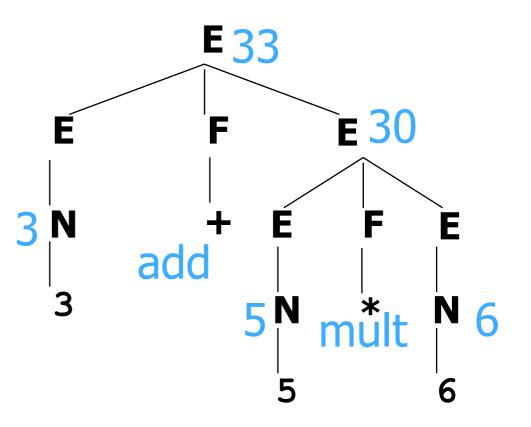
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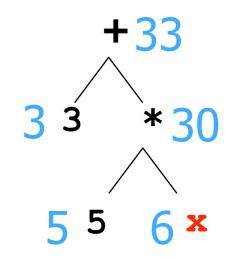


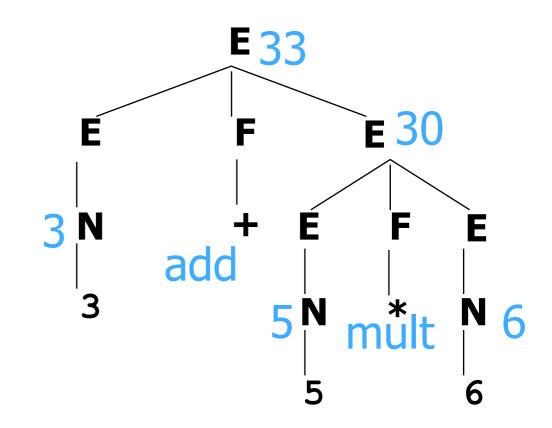


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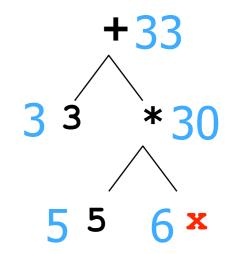


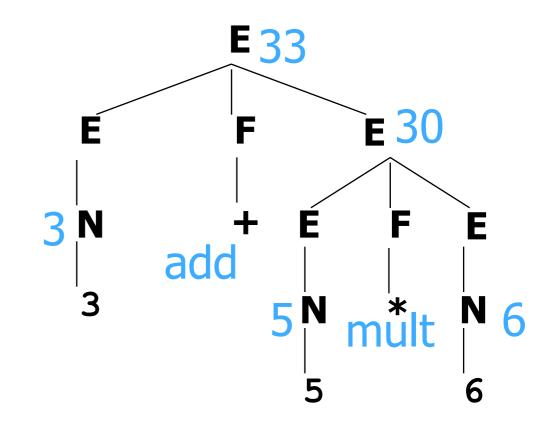




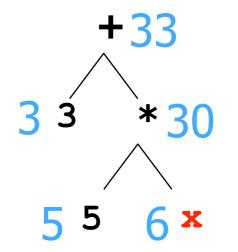


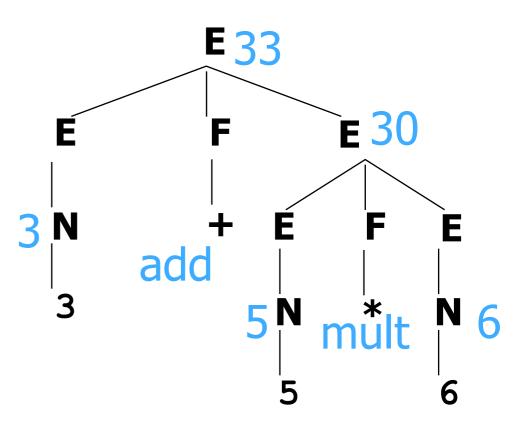
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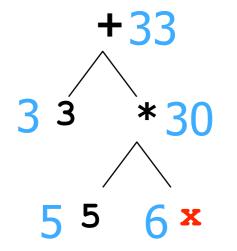


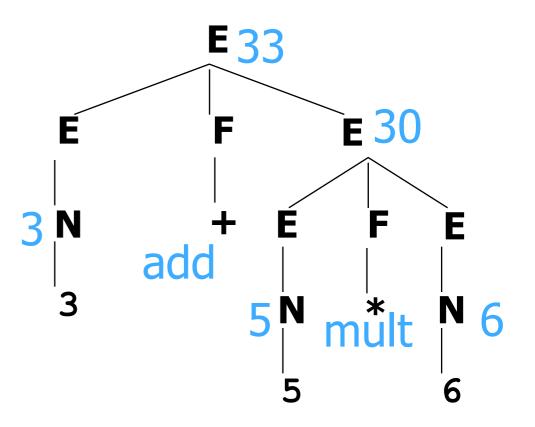
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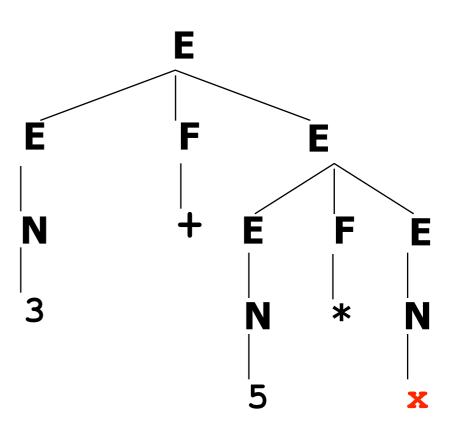




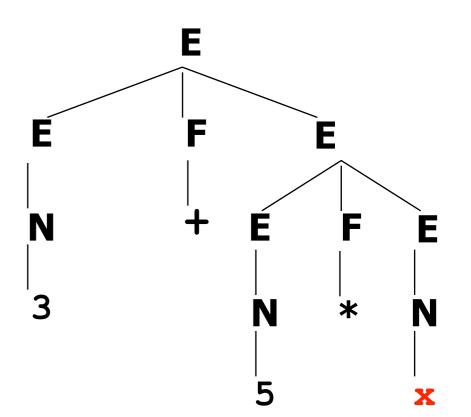
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- Analogies in language?





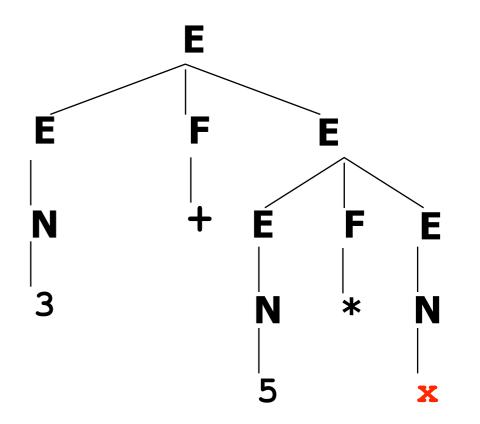


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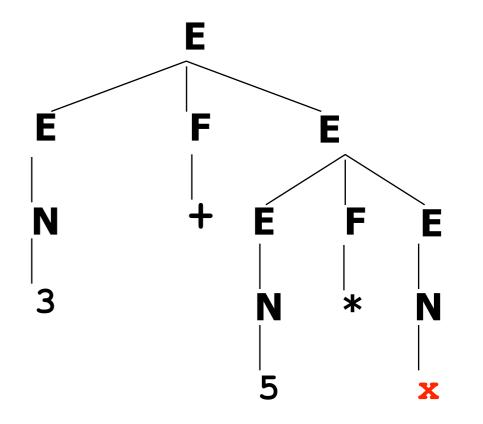


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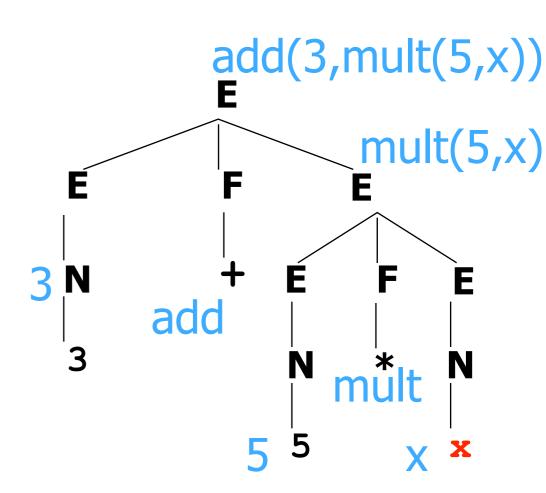
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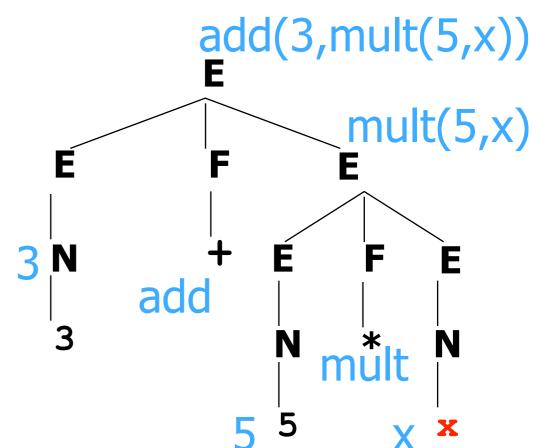


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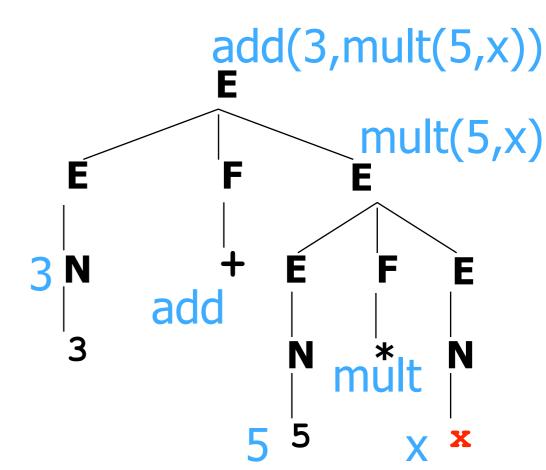
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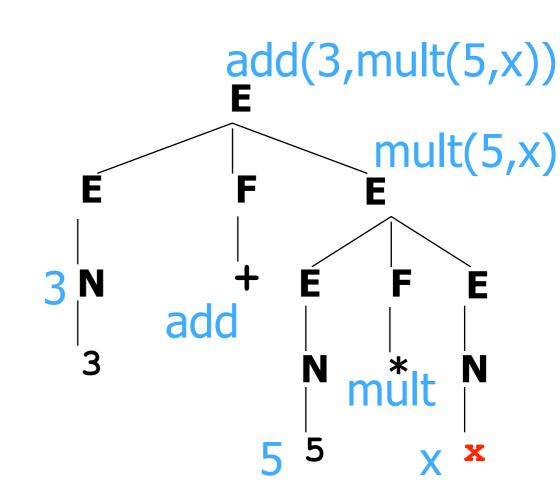
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5* (x+1) -2 is a different expression that produces equivalent code (can be converted to the previous code by optimization) Analogies in language?



- We understand if we can respond appropriately
 - ok for commands, questions (these demand response)
 - Computer, warp speed 5"
 - "throw axe at dwarf"
 - "put all of my blocks in the red box"
 - imperative programming languages
 - SQL database queries and other questions

We understand statement if we can determine its truth

- ok, but if you knew whether it was true, why did anyone bother telling it to you?
- comparable notion for understanding NP is to compute what the NP refers to, which might be useful

- We understand statement if we know how one could (in principle) determine its truth
 - What are exact conditions under which it would be true?
 - necessary + sufficient
 - Equivalently, derive all its consequences
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- We understand statement if we can use it to answer questions [very similar to above – requires reasoning]
 - Easy: John ate pizza. What was eaten by John?
 - Hard: White's first move is P-Q4. Can Black checkmate?
 - Constructing a procedure to get the answer is enough

- Paraphrase, "state in your own words" (English to English translation)
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- Translation to logical form that we can reason about

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 - Function might take other functions as arguments!

Logic: Lambda Terms

Lambda terms:

- A way of writing "anonymous functions"
 - No function header or function name
 - But defines the key thing: **behavior** of the function
 - Just as we can talk about 3 without naming it "x"
- Let square = $\lambda p p^* p$
- Equivalent to int square(p) { return p*p; }
- But we can talk about $\lambda p p^* p$ without naming it
- Format of a lambda term: λ variable expression

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- even(x) is true if x is even
- How about even(square(x))?
- λx even(square(x)) is true of numbers with even squares
 Just apply rules to get λx (even(x*x)) = λx (x*x mod 2 == 0)

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 - This happens to denote the same predicate as even does

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- Remember: square can be written as λx square(x)
 - And now times can be written as $\lambda x \lambda y$ times(x,y)

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- What is executed by loves(john, mary) ?

- Thus, have "constants" that name some of the entities and functions (e.g., *):
 - GeorgeWBush an entity
 - red a predicate on entities
 - •holds of just the red entities: red(x) is true if x is red!
 - Ioves a predicate on 2 entities
 - loves(GeorgeWBush, LauraBush)
 - Question: What does loves(LauraBush) denote?
- Constants used to define meanings of words
- Meanings of phrases will be built from the constants

most – a predicate on 2 predicates on entities

- most(pig, big) = "most pigs are big"
 - Equivalently, $most(\lambda x pig(x), \lambda x big(x))$
- returns true if most of the things satisfying the first predicate also satisfy the second predicate

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 - Equivalently, $most(\lambda x pig(x), \lambda x big(x))$
- returns true if most of the things satisfying the first predicate also satisfy the second predicate
- similarly for other quantifiers
 - all(pig,big) (equivalent to $\forall x \text{ pig}(x) \Rightarrow \text{big}(x)$)
 - exists(pig,big) (equivalent to ∃x pig(x) AND big(x))
 - can even build complex quantifiers from English phrases:

• "between 12 and 75"; "a majority of"; "all but the smallest 2"

A reasonable representation?

- Gilly swallowed a goldfish
 First attempt: swallowed(Gilly, goldfish)
- Returns true or false. Analogous to
 - prime(17)
 - equal(4,2+2)
 - Ioves(GeorgeWBush, LauraBush)
 - swallowed(Gilly, Jilly)
- ... or is it analogous?

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- goldfish isn't the name of a unique object the way Gilly is
- In particular, don't want Gilly swallowed a goldfish and Milly swallowed a goldfish
 to translate as swallowed(Gilly, goldfish) AND swallowed(Milly, goldfish) since probably not the same goldfish ...

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- Or using one of our quantifier predicates:
 - exists(λg goldfish(g), λg swallowed(Gilly,g))
 - Equivalently: exists(goldfish, swallowed(Gilly))
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 - Equivalently: exists(goldfish, swallowed(Gilly))
 - "In the set of goldfish there exists one swallowed by Gilly"
- Here goldfish is a predicate on entities
 - This is the same semantic type as red
 - But goldfish is noun and red is adjective .. #@!?



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(Simplify Notation)

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- Previous: It past(t) AND exists(goldfish, swallow(t,Gilly))
- Why stop at time? An event has other properties:
 - [Gilly] swallowed [a goldfish] [on a dare] [in a telephone booth] [with 30 other freshmen] [after many bottles of vodka had been consumed].
 - Specifies who what why when ...

- Gilly swallowed a goldfish

- Previous: It past(t) AND exists(goldfish, swallow(t,Gilly))
- Why stop at time? An event has other properties:
 - [Gilly] swallowed [a goldfish] [on a dare] [in a telephone booth] [with 30 other freshmen] [after many bottles of vodka had been consumed].
 - Specifies who what why when ...
- Replace time variable t with an event variable e
 - Be past(e), act(e,swallowing), swallower(e,Gilly), exists(goldfish, swallowee(e)), exists(booth, location(e)), ...
 - As with probability notation, a comma represents AND
 - Could define past as λe ∃t before(t,now), ended-at(e,t)

-Gilly swallowed a goldfish in \underline{a} booth

- Je past(e), act(e,swallowing), swallower(e,Gilly), exists(goldfish, swallowee(e)), <u>exists(booth, location(e)), ...</u>
- Gilly swallowed a goldfish in <u>every</u> booth
 - Be past(e), act(e,swallowing), swallower(e,Gilly), exists(goldfish, swallowee(e)), <u>all(booth, location(e)),</u> ...

Does this mean what we'd expect??

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Does this mean what we'd expect?? says that there's only one event with a single goldfish getting swallowed that took place in a lot of booths ...

Groucho Marx celebrates quantifier order ambiguity:

- In this country <u>a woman</u> gives birth <u>every 15 min</u>. Our job is to find that woman and stop her.
- ∃woman (∀15min gives-birth-during(woman, 15min))
- ► ¥15min (∃woman gives-birth-during(15min, woman))
- Surprisingly, both are possible in natural language!
- Which is the joke meaning (where it's always the same woman) and why?

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 ∃g goldfish(g), swallowee(e,g) ∀b booth(b)⇒location(e,b)
- Does this mean what we'd expect??
 - It's ∃e ∀b which means same event for every booth
 - Probably false unless Gilly can be in every booth during her swallowing of a single goldfish

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Other reading (∀b ∃e) involves <u>quantifier raising</u>:

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Other reading (Vb 3e) involves <u>quantifier raising</u>:

<u>all(booth, λb</u> [∃e past(e), act(e,swallowing), swallower (e,Gilly), exists(goldfish, swallowee(e)), location(e,b)])

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- Other reading $(\forall b \exists e)$ involves <u>quantifier raising</u>:
 - all(booth, λb [∃e past(e), act(e,swallowing), swallower (e,Gilly), exists(goldfish, swallowee(e)), location(e,b)])
 "for all booths b, there was such an event in b"

- Be act(e,wanting), wanter(e,Willy), exists(unicorn, λu wantee(e,u))
 - "there is a particular unicorn u that Willy wants"
 - In this reading, the wantee is an individual entity

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- Be act(e,wanting), wanter(e,Willy), wantee(e, λu unicorn(u))
 - "Willy wants any entity u that satisfies the unicorn predicate"
 - In this reading, the wantee is a <u>type</u> of entity
 - Sentence doesn't claim that such an entity exists

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Willy wants Lilly to get married

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- "Willy wants any event e' in which Lilly gets married"
- Here the wantee is a <u>type</u> of event
- Sentence doesn't claim that such an event exists
- Intensional verbs besides want: hope, doubt, believe,...

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- Traditional solution involves "possible-world semantics"
 - Can imagine other worlds where set of unicorn \neq set of dodos
 - Other worlds also useful for: You must pay the rent You can pay the rent If you hadn't, you'd be homeless



• Willy wants Lilly to get married

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 Be present(e), act(e,wanting), wanter(e,Willy), wantee(e, λf [act(f,marriage), marrier(f,Lilly)])

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 Be present(e), act(e,wanting), wanter(e,Willy), wantee(e, λf [act(f,marriage), marrier(f,Lilly)])

Willy wants to get marriedSame as Willy wants Willy to get married

- Willy wants Lilly to get married
 - Be present(e), act(e,wanting), wanter(e,Willy), wantee(e, λf [act(f,marriage), marrier(f,Lilly)])
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 - Same as Willy wants Willy to get married
 - Just as easy to represent as Willy wants Lilly ...

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- Willy wants to get married
 - Same as Willy wants Willy to get married
 - Just as easy to represent as Willy wants Lilly ...
 - The only trick is to construct the representation from the syntax. The empty subject position of "to get married" is said to be <u>controlled</u> by the subject of "wants."

- expert
 - λg expert(g)

expert

- λg expert(g)
- big fat expert
 - λg big(g), fat(g), expert(g)
 - But: bogus expert
 - Wrong: λg bogus(g), expert(g)
 - Right: λg (bogus(expert))(g) ... bogus maps to new concept

expert

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- Baltimore expert (white-collar expert, TV expert...)
 - Ag Related(Baltimore, g), expert(g) expert from Baltimore
 - Or with different intonation:
 - λg (Modified-by(Baltimore, expert))(g) expert on Baltimore
 - Can't use Related for this case: law expert and dog catcher
 - = λg Related(law,g), expert(g), Related(dog, g), catcher(g)
 - = dog expert and law catcher

the goldfish that Gilly swallowed

every goldfish that Gilly swallowed

three goldfish that Gilly swallowed

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 λg [goldfish(g), swallowed(Gilly, g)]

the goldfish that Gilly swallowed

every goldfish that Gilly swallowed

three goldfish that Gilly swallowed

λg [goldfish(g), swallowed(Gilly, g)]

like an adjective!
three swallowed-by-Gilly goldfish

Nouns and Their Modifiers

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three goldfish that Gilly swallowed

 λg [goldfish(g), swallowed(Gilly, g)]

like an adjective!
three swallowed-by-Gilly goldfish

Or for real: λg [goldfish(g), ∃e [past(e), act(e,swallowing), swallower(e,Gilly), swallowee(e,g)]]

Lili passionately wants Billy

- Wrong?: passionately(want(Lili,Billy)) = passionately(true)
- Better: (passionately(want))(Lili,Billy)
- Best: Je present(e), act(e,wanting), wanter(e,Lili), wantee(e, Billy), manner(e, passionate)

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- Best: Je present(e), act(e,wanting), wanter(e,Lili), wantee(e, Billy), manner(e, passionate)
- Lili often stalks Billy
 - often(stalk))(Lili,Billy)
 - many(day, λd ∃e present(e), act(e,stalking), stalker(e,Lili), stalkee(e, Billy), during(e,d))

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 - many(day, λd ∃e present(e), act(e,stalking), stalker(e,Lili), stalkee(e, Billy), during(e,d))
- Lili obviously likes Billy
 - (obviously(like))(Lili,Billy) one reading
 - obvious(like(Lili, Billy)) another reading

- What is the meaning of a full sentence?
 - Depends on the punctuation mark at the end. $\ensuremath{\textcircled{\circ}}$
 - Billy likes Lili.
 assert(like(B,L))
 - Billy likes Lili?

 ask(like(B,L))
 - or more formally, "Does Billy like Lili?"
 - Billy, like Lili!

 command(like(B,L))
 - or more accurately, "Let Billy like Lili!"

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or more accurately, "Let Billy like Lili!"

 Let's try to do this a little more precisely, using event variables etc.

• What did Gilly swallow?

ask(λx ∃e past(e), act(e,swallowing), swallower(e,Gilly), swallowee(e,x))

• Argument is identical to the modifier "that Gilly swallowed"

Is there any common syntax?

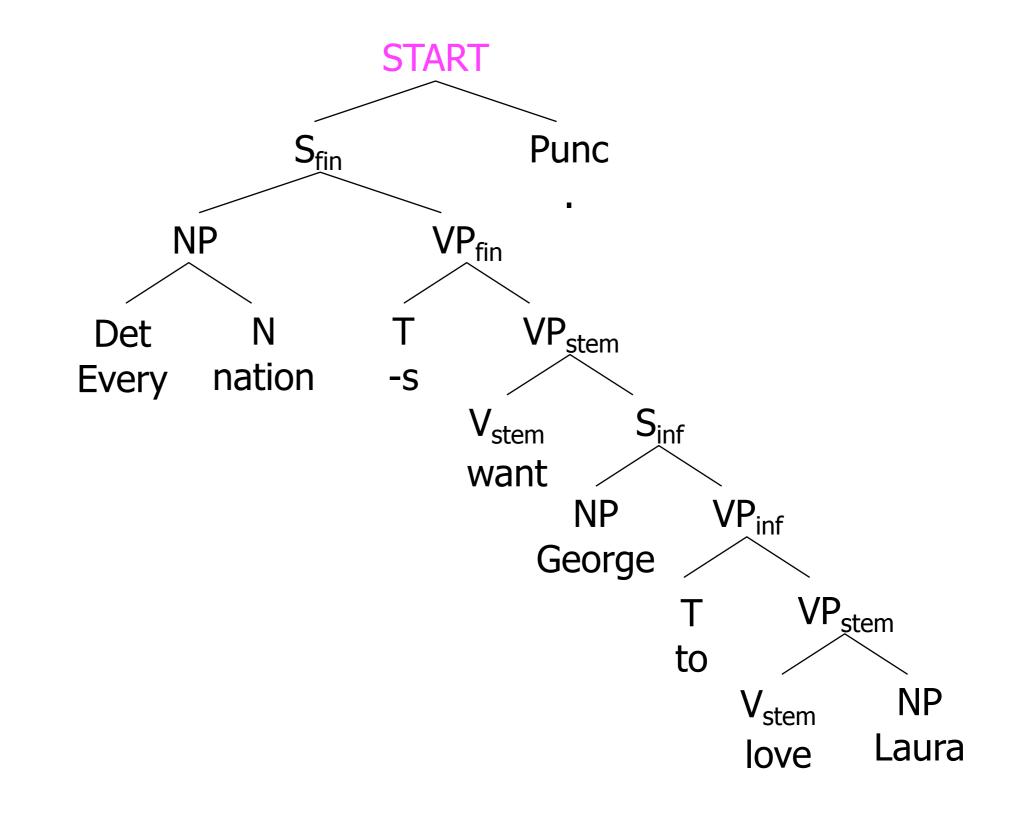
- What did Gilly swallow?
 - **ask**(λx ∃e past(e), act(e,swallowing), swallower(e,Gilly), swallowee(e,x))
 - Argument is identical to the modifier "that Gilly swallowed"
 - Is there any common syntax?
- Eat your fish!
 - command(λf act(f,eating), eater(f,Hearer), eatee(...))

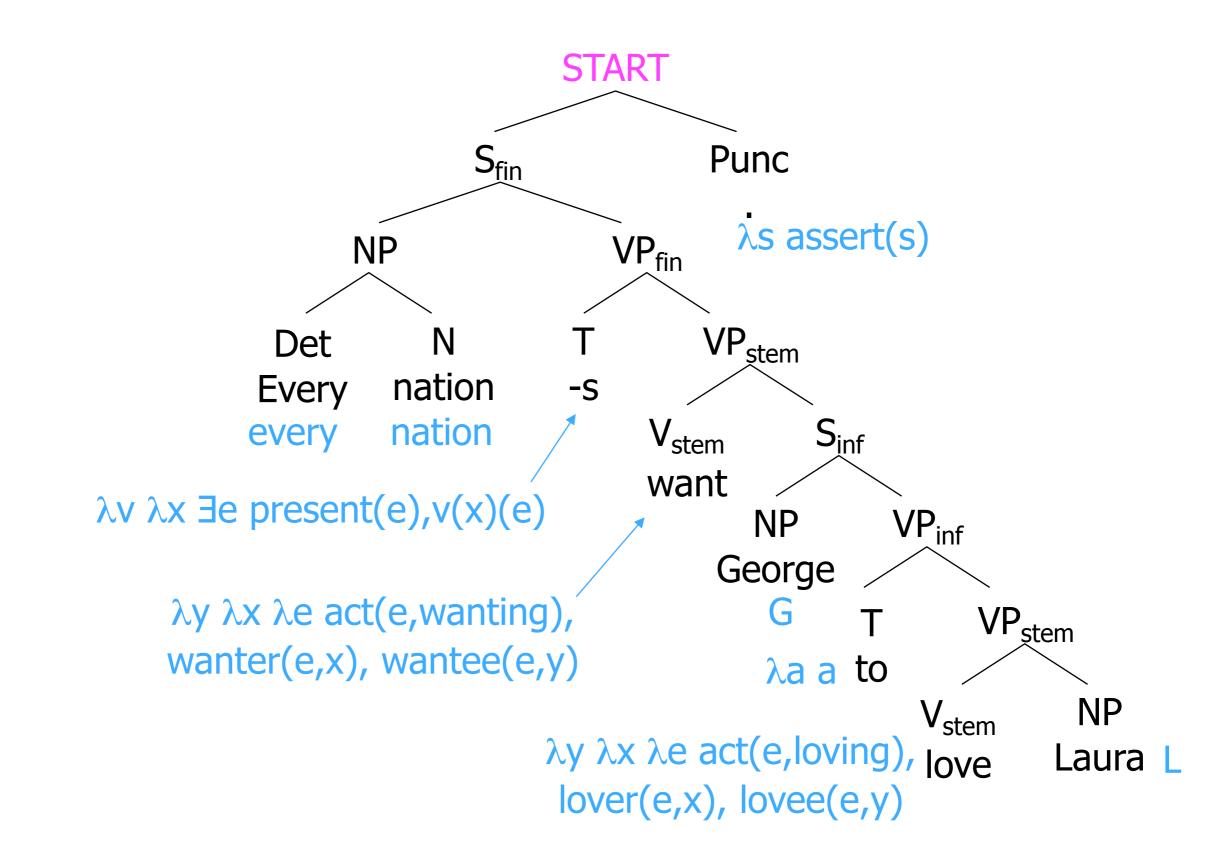
- What did Gilly swallow?
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 - Argument is identical to the modifier "that Gilly swallowed"
 - Is there any common syntax?
- Eat your fish!
 - command(λf act(f,eating), eater(f,Hearer), eatee(...))
- I ate my fish.
 - assert(∃e past(e), act(e,eating), eater(f,Speaker), eatee(...))

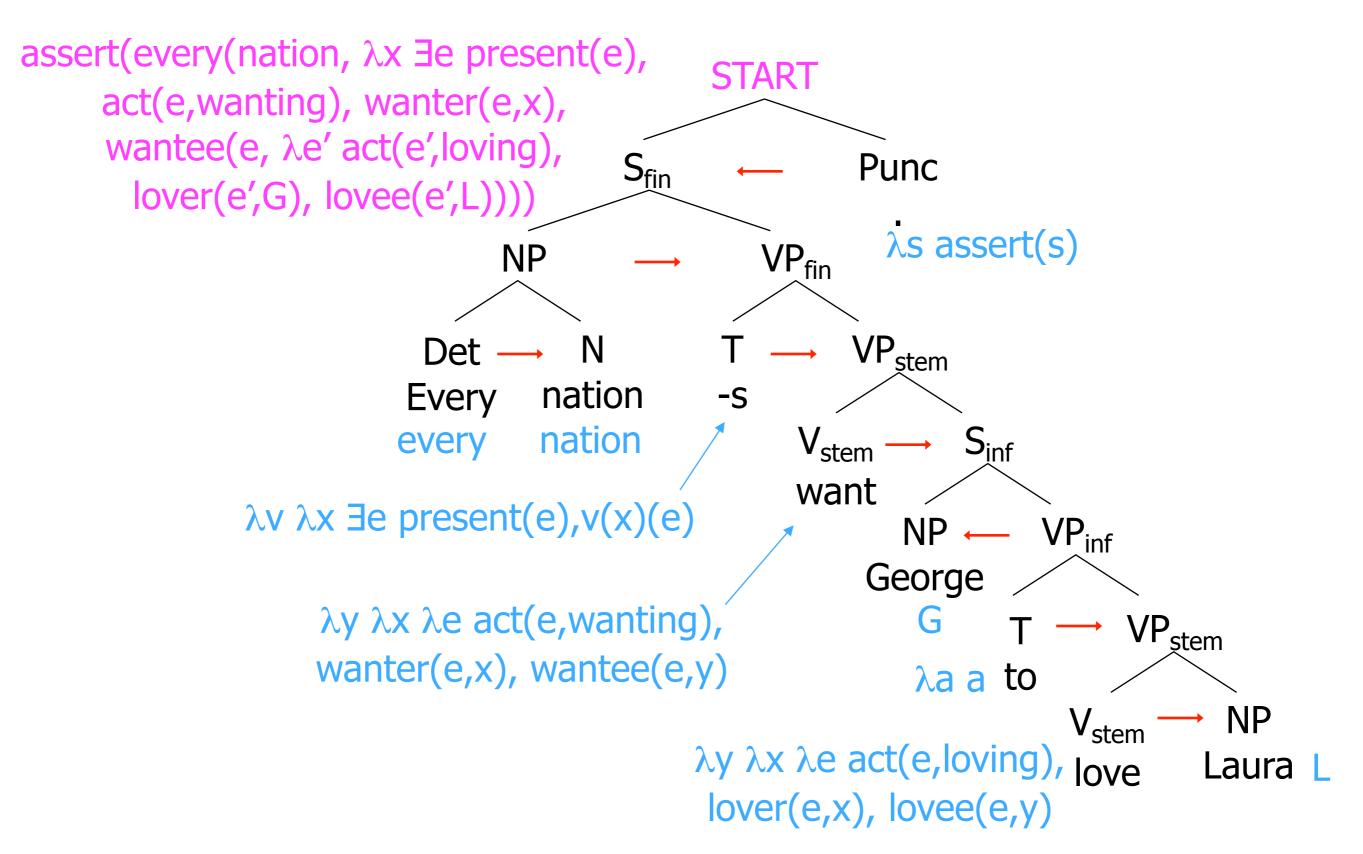
 We've discussed what semantic representations should look like.

But how do we get them from sentences???

- First parse to get a syntax tree.
- Second look up the semantics for each word.
- Third build the semantics for each constituent
 - Work from the bottom up
 - The syntax tree is a "recipe" for how to do it





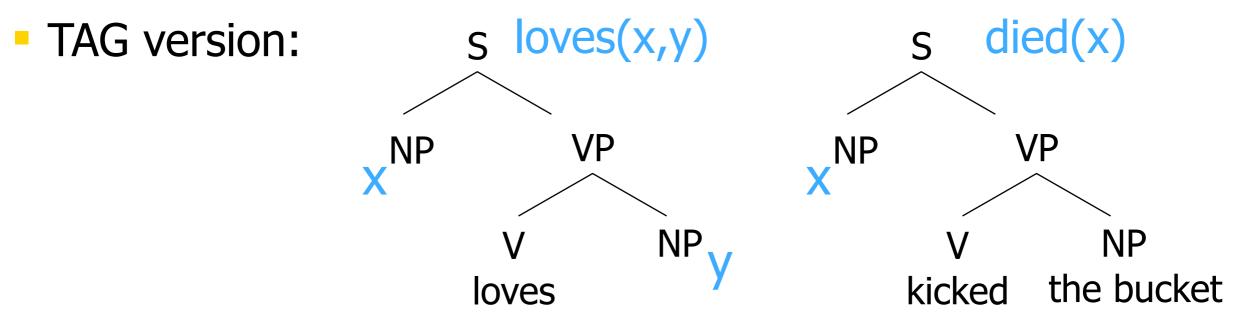


- Add a "sem" feature to each context-free rule
 - S \rightarrow NP loves NP
 - $S[sem=loves(x,y)] \rightarrow NP[sem=x] loves NP[sem=y]$
 - Meaning of S depends on meaning of NPs

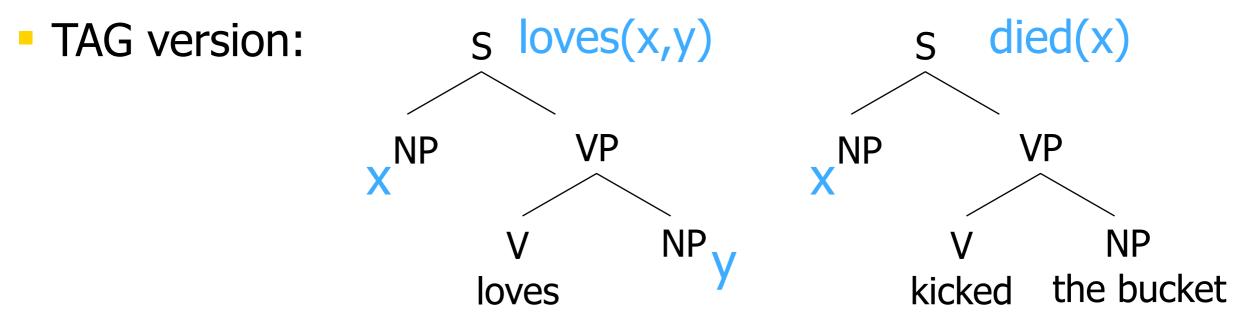
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- TAG version:

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 - $S[sem=loves(x,y)] \rightarrow NP[sem=x] loves NP[sem=y]$
 - Meaning of S depends on meaning of NPs
- TAG version:
 S loves(x,y)
 NP VP
 V VP
 V NP
 V

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 Template filling: S[sem=showflights(x,y)] → I want a flight from NP[sem=x] to NP[sem=y]

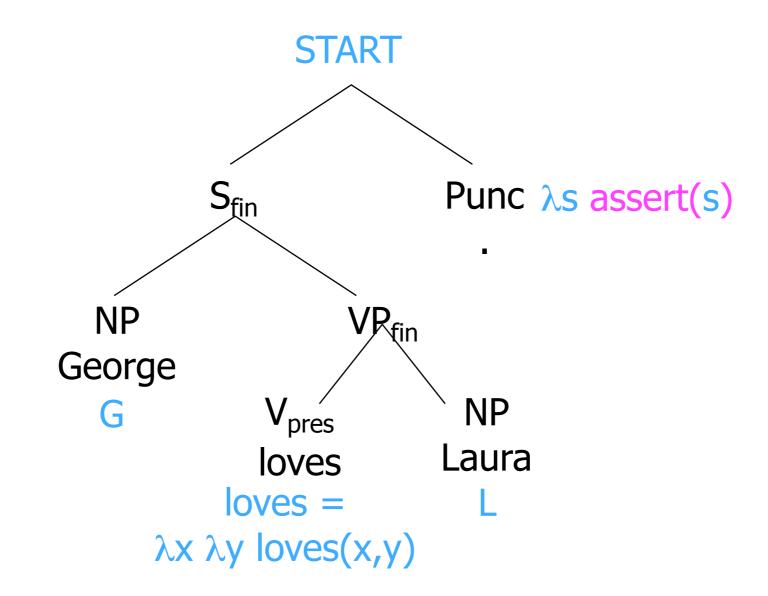
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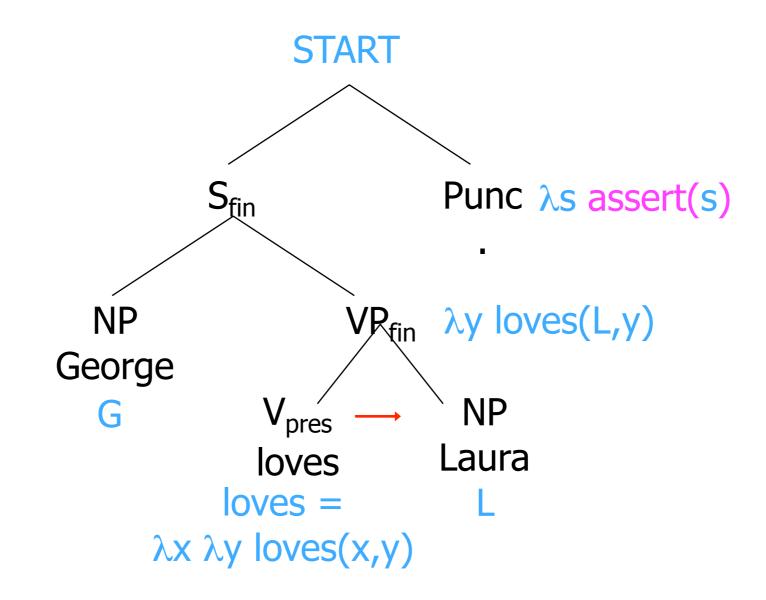
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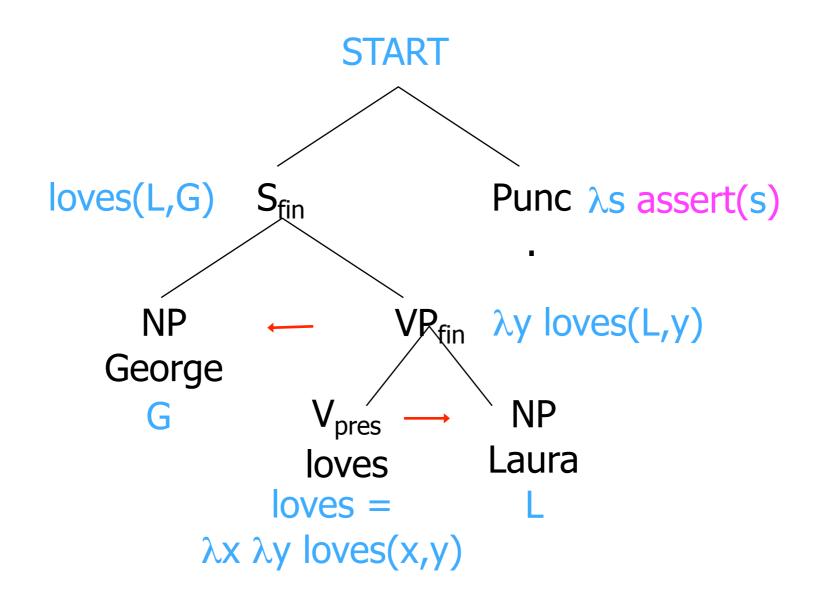
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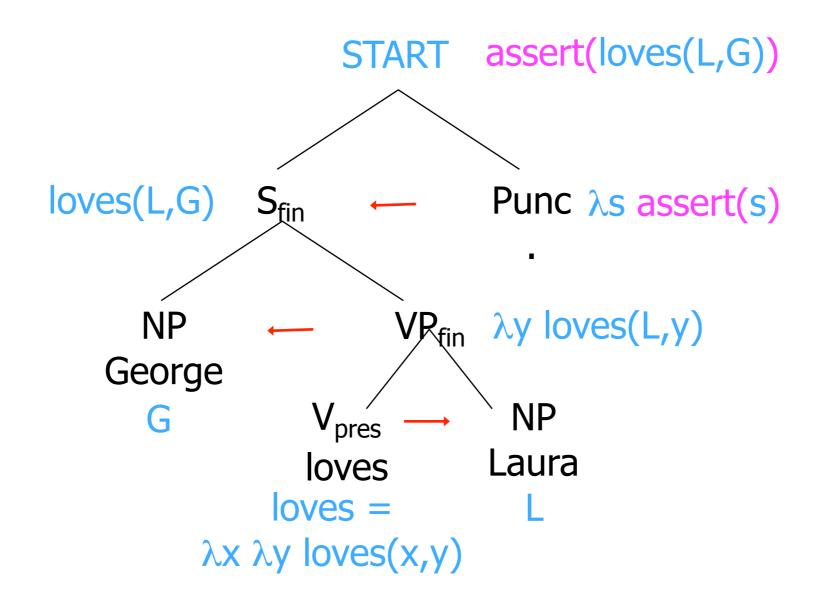
NOW George loves Laura has sem=loves(Laura)(George)

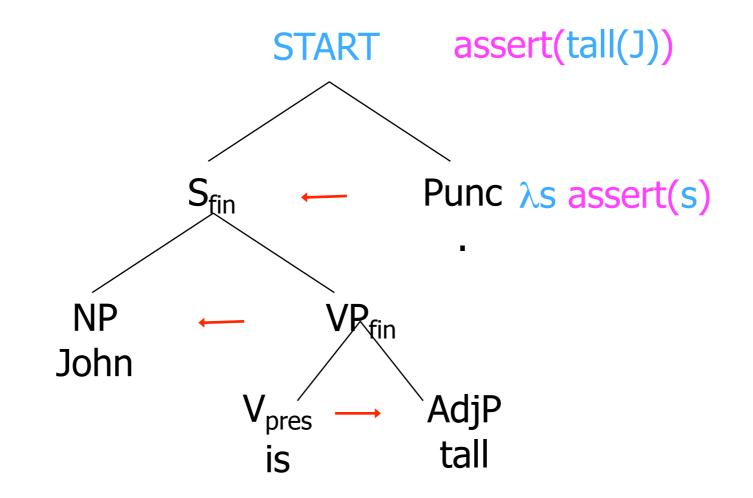
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- NOW George loves Laura has sem=loves(Laura)(George)
- In this manner we'll sketch a version where
 - Still compute semantics bottom-up
 - Grammar is in Chomsky Normal Form
 - So each node has 2 children: 1 function & 1 argument
 - To get its semantics, apply function to argument!

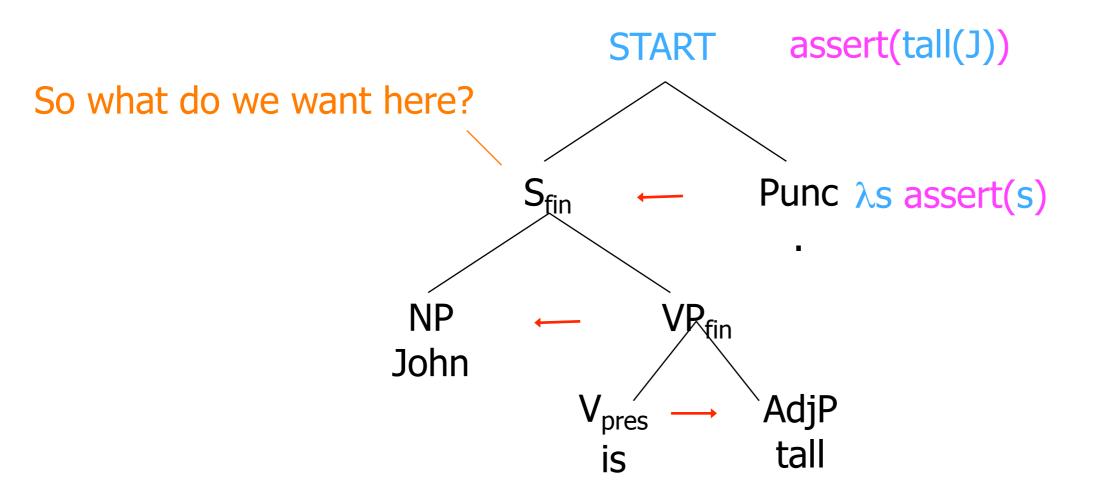


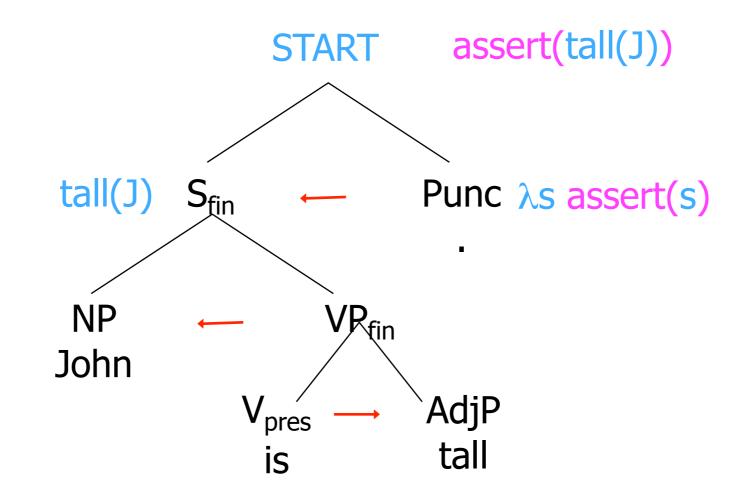


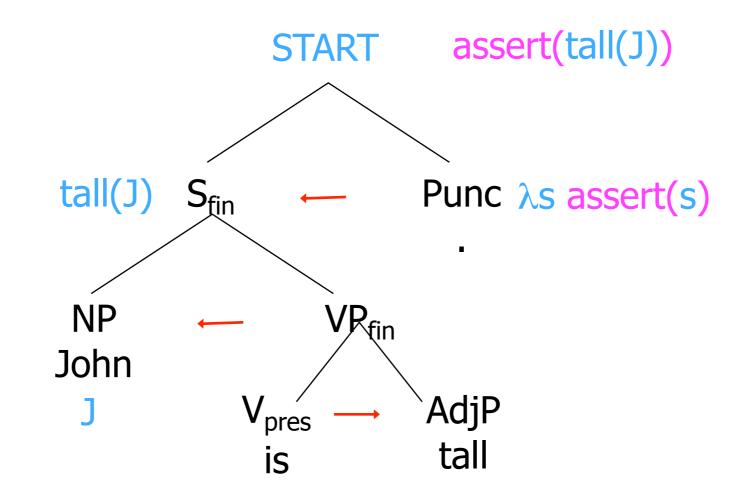


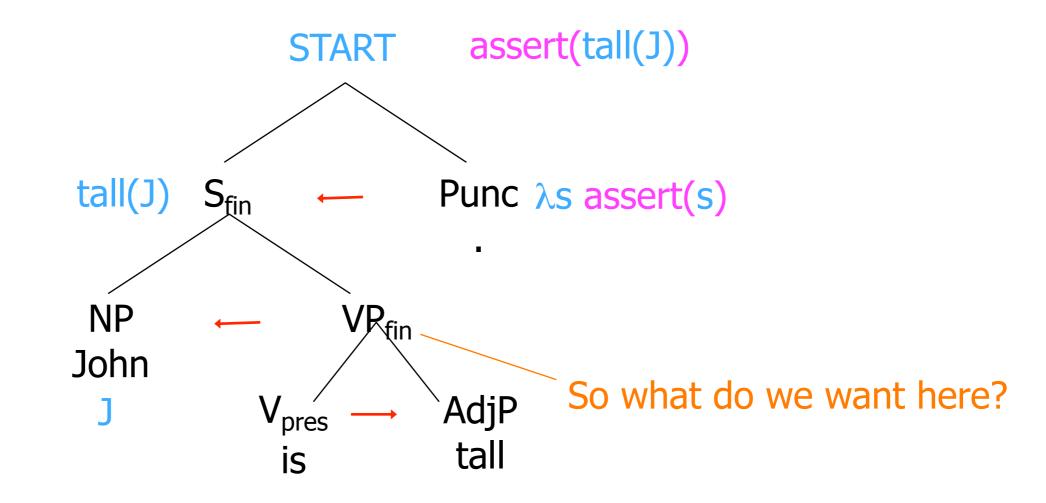


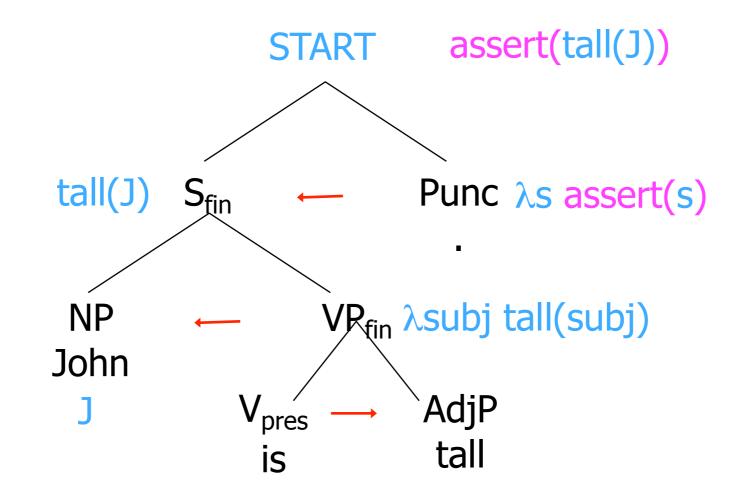


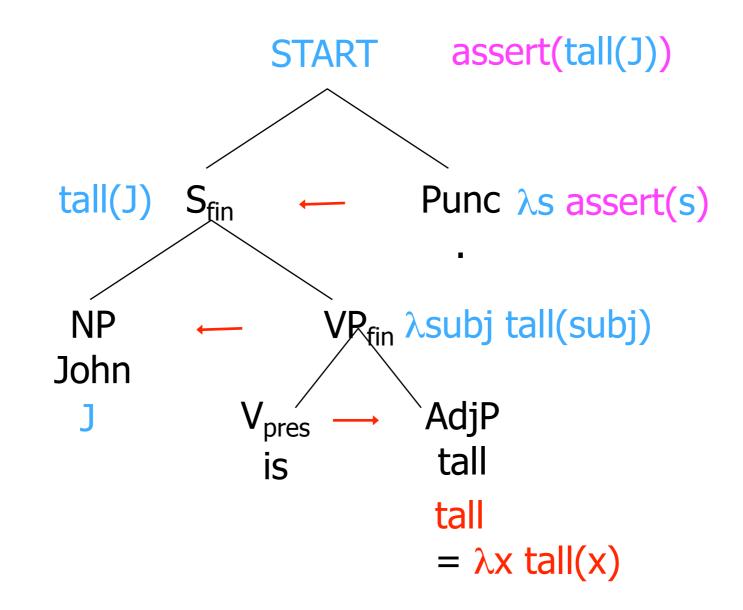


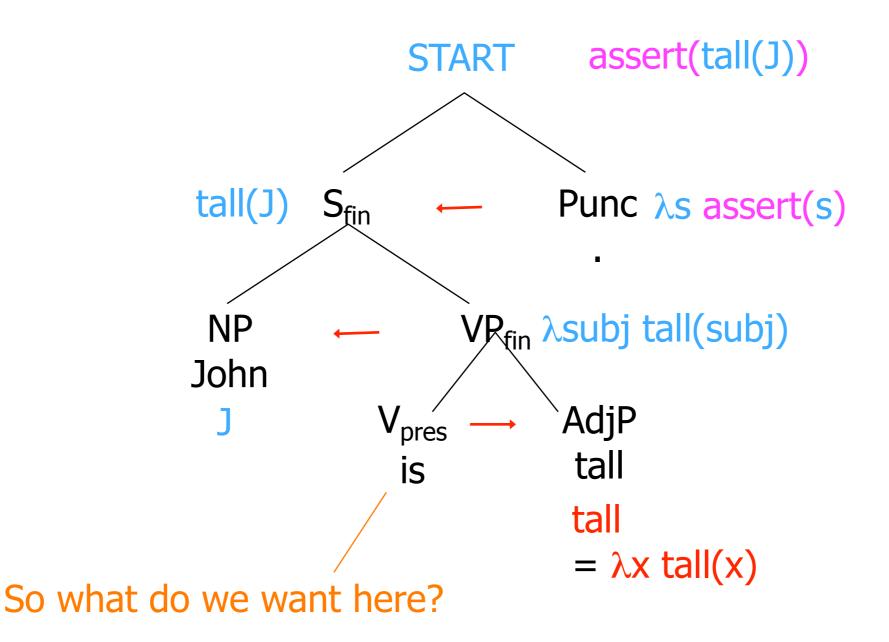


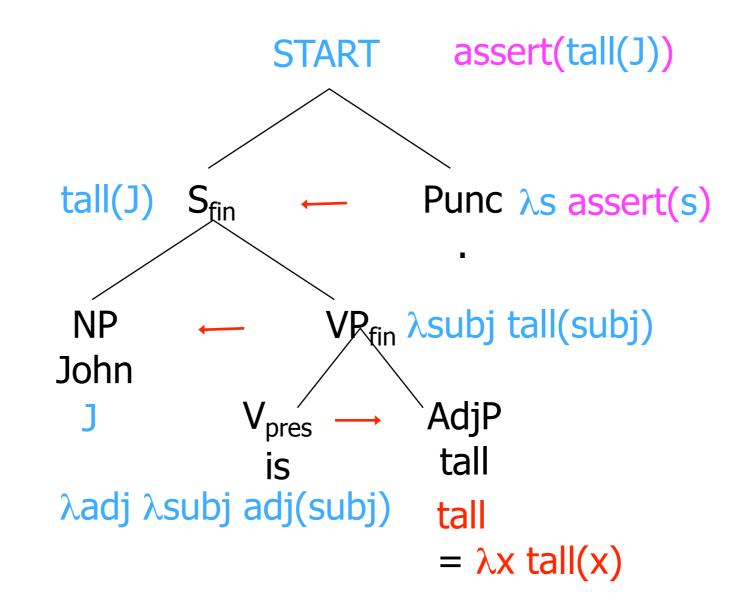


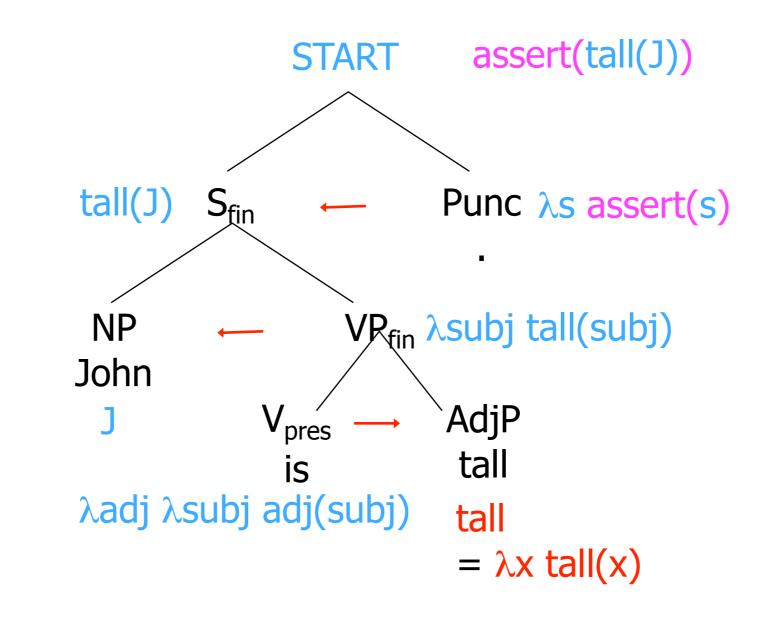




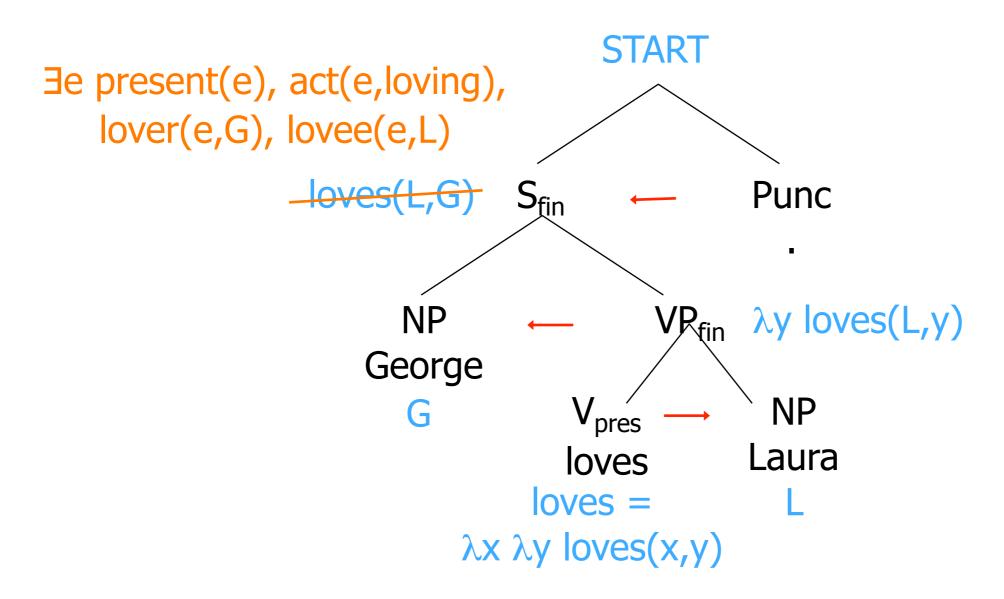


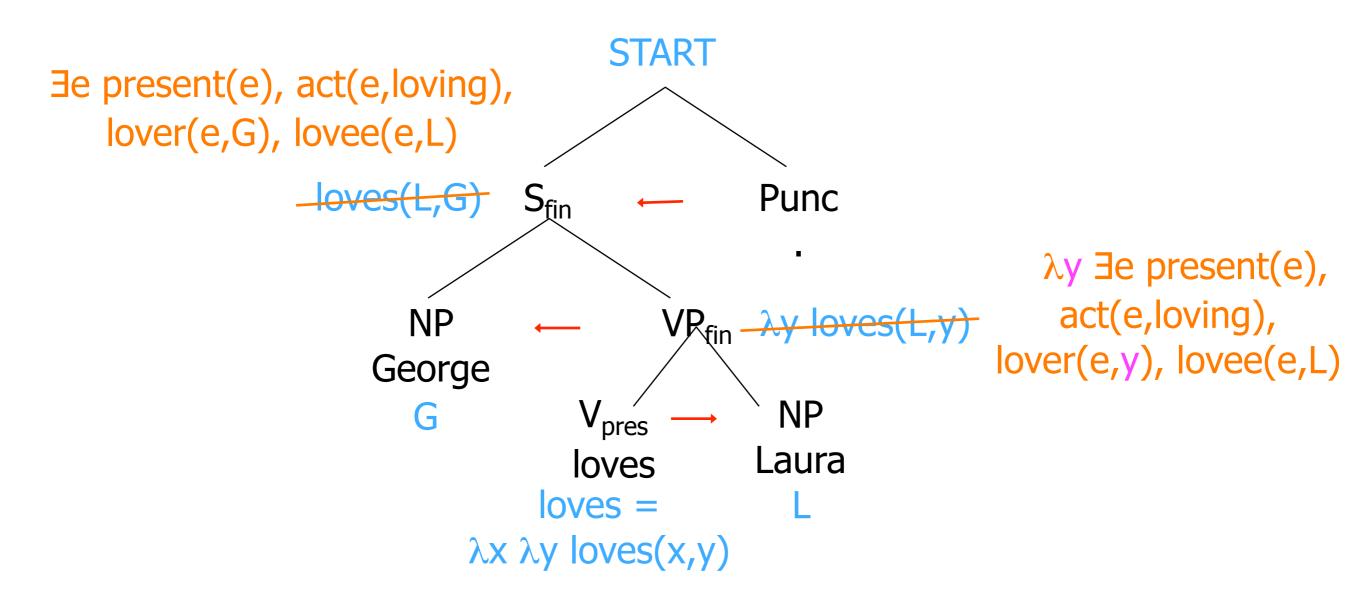


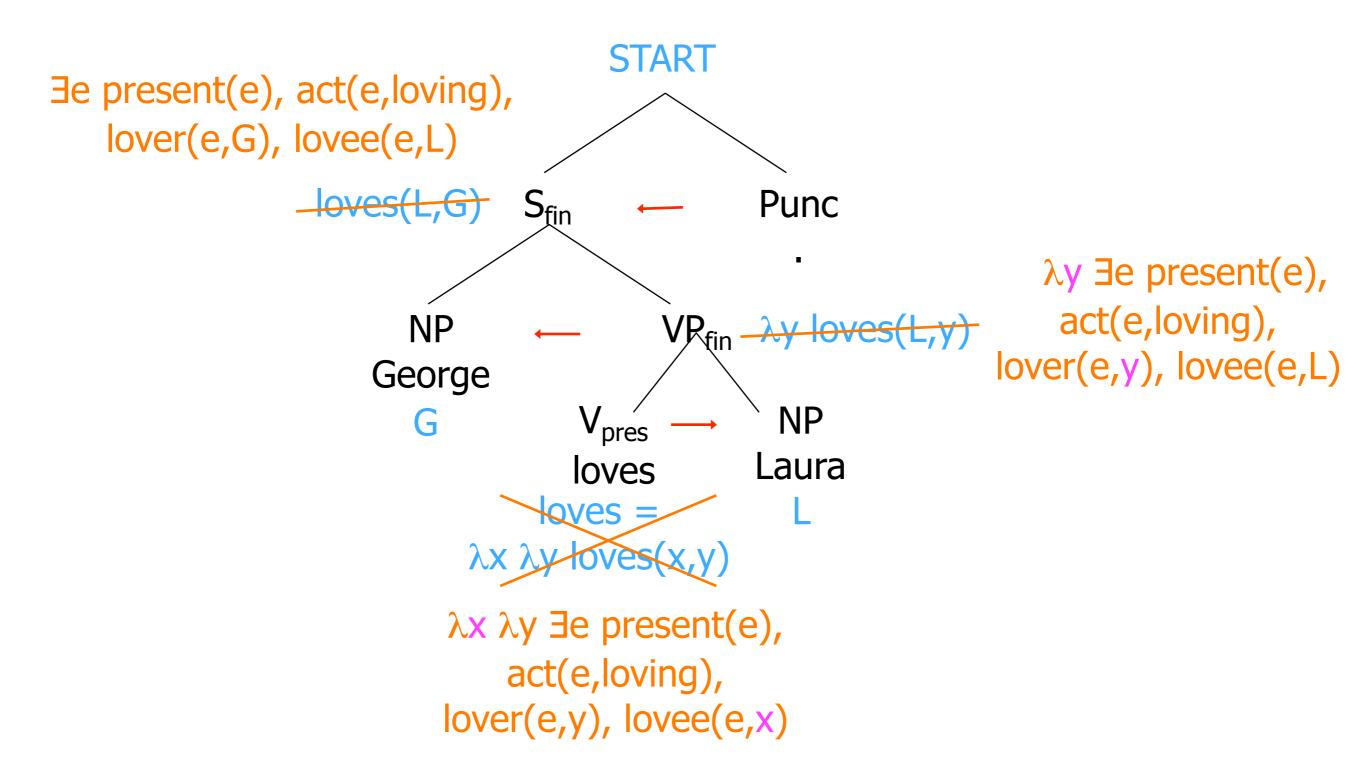


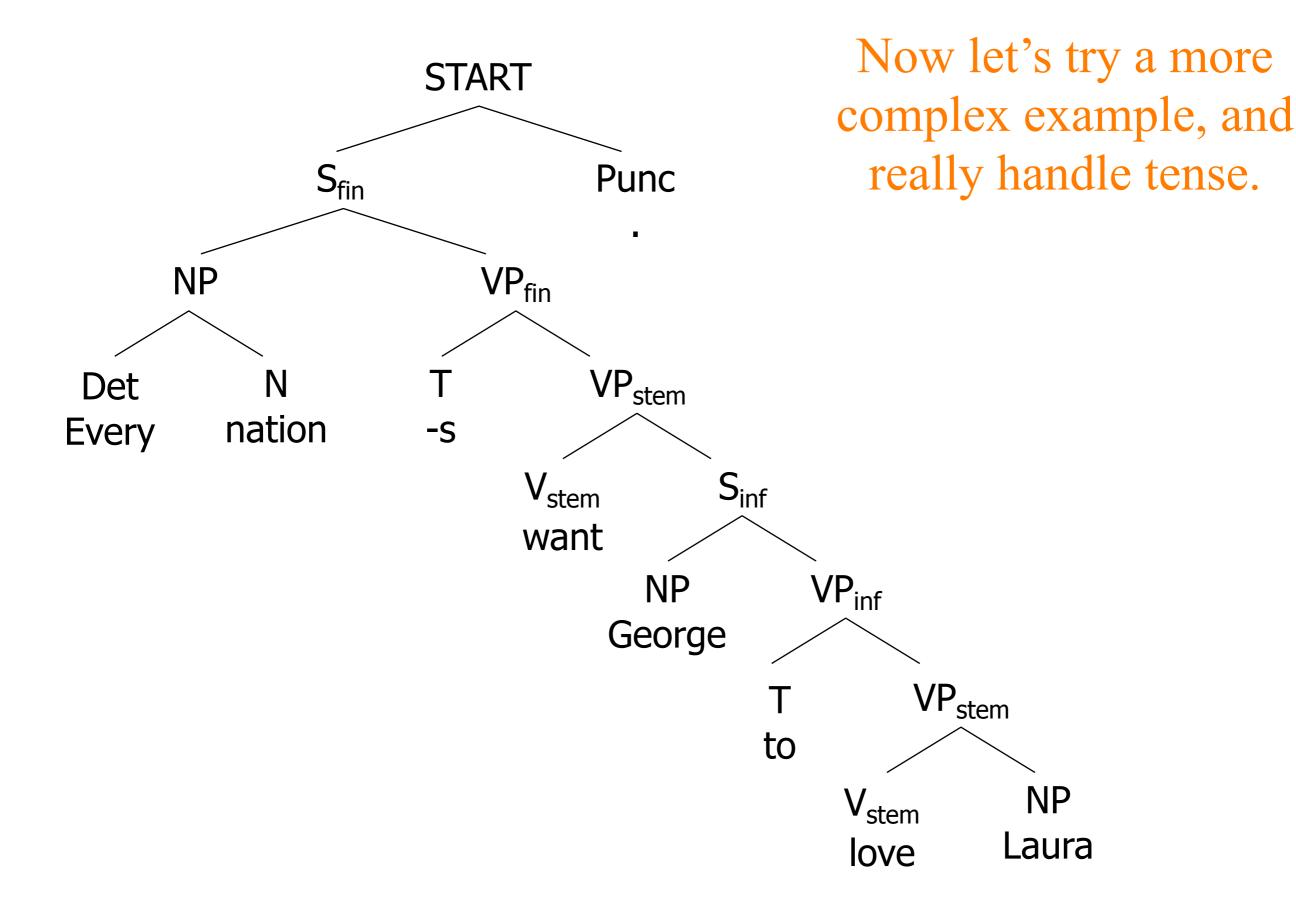


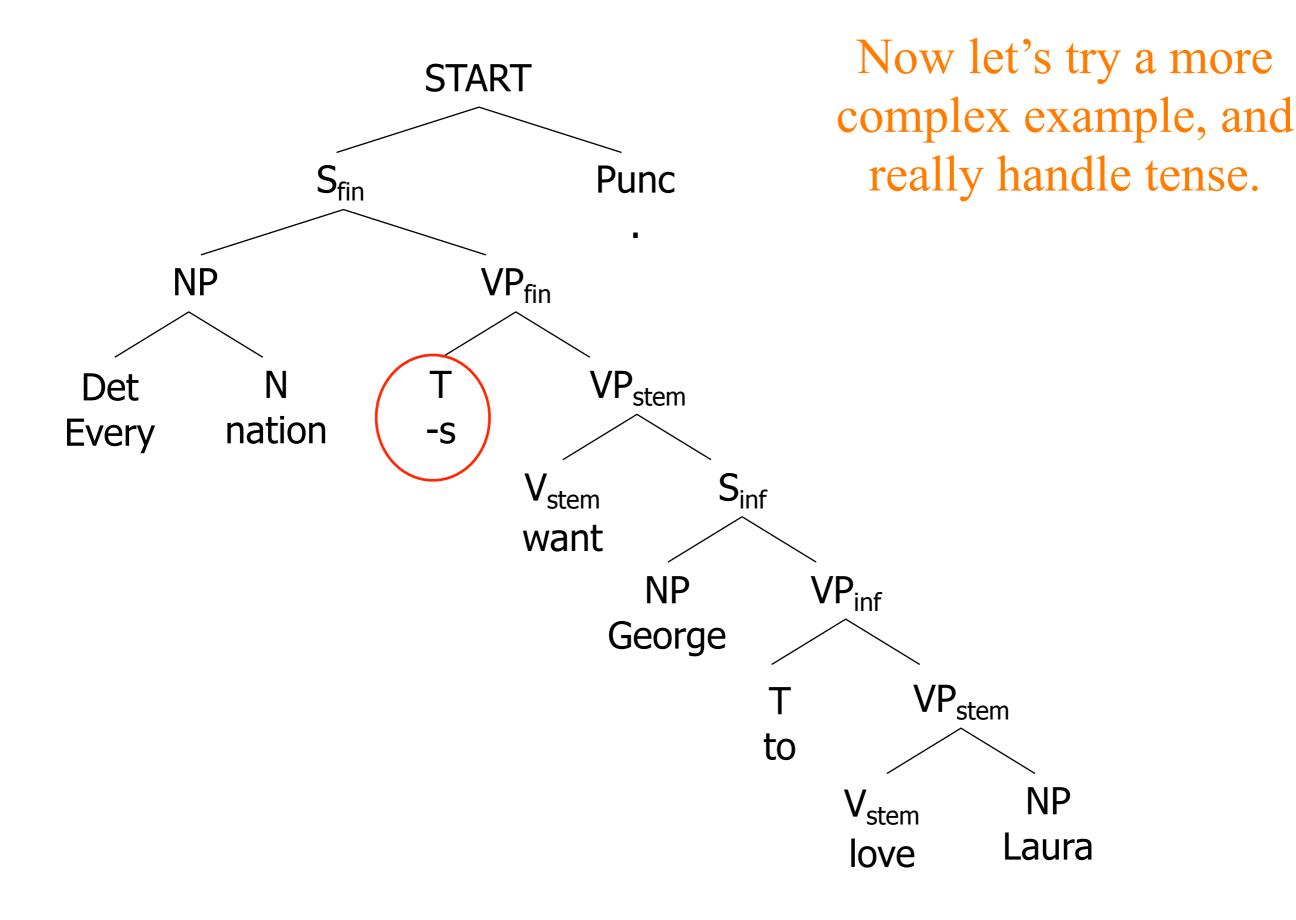
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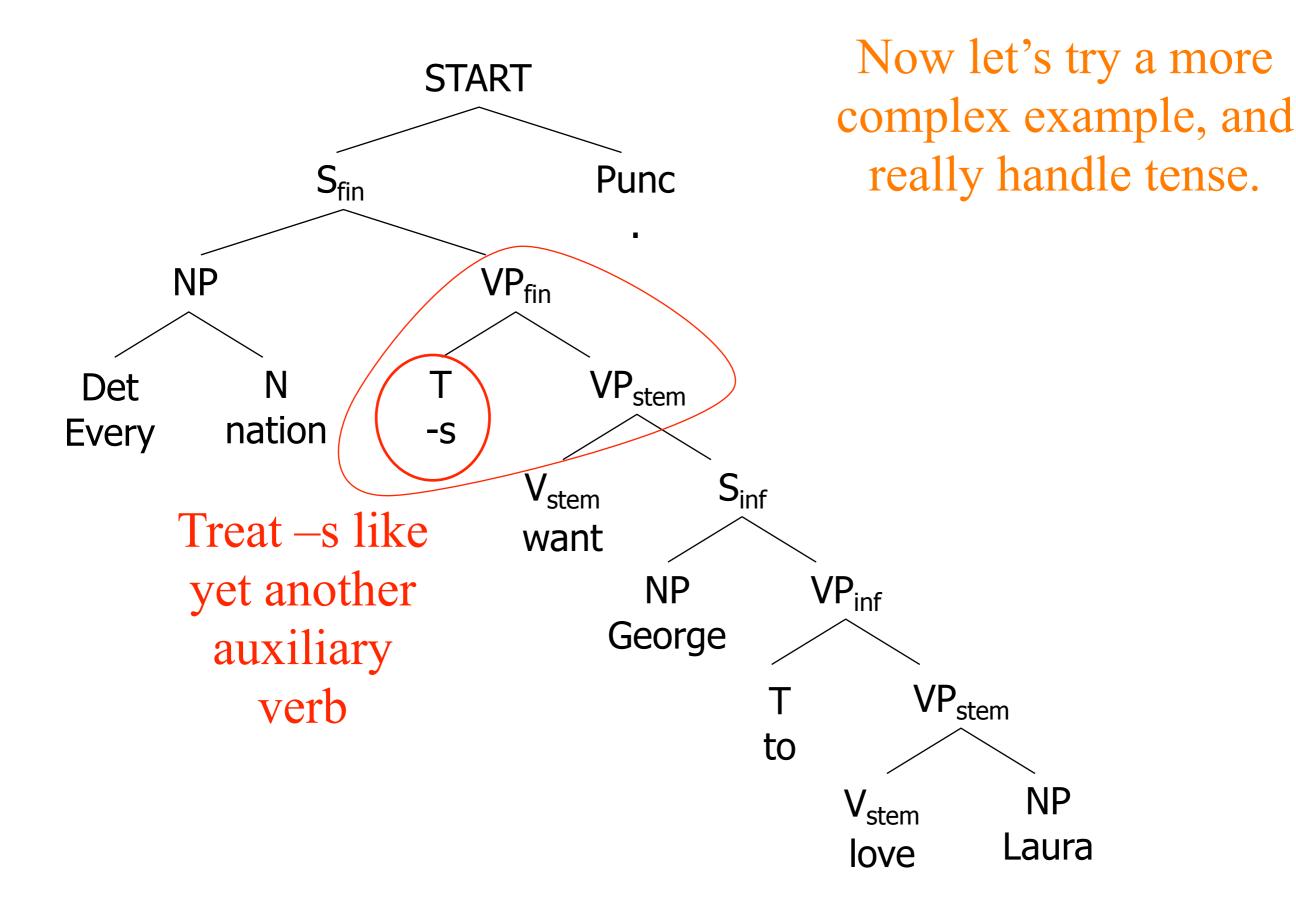


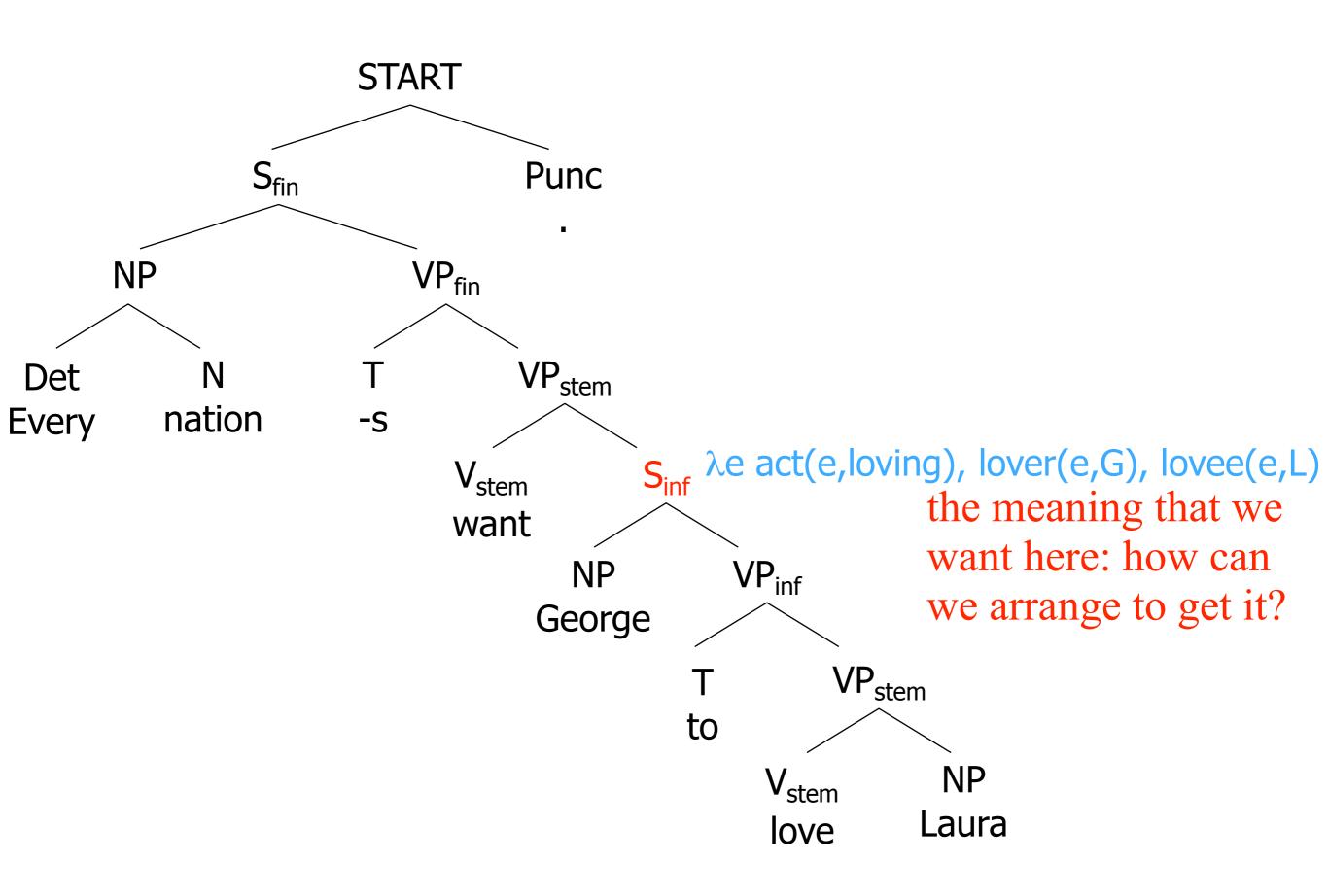


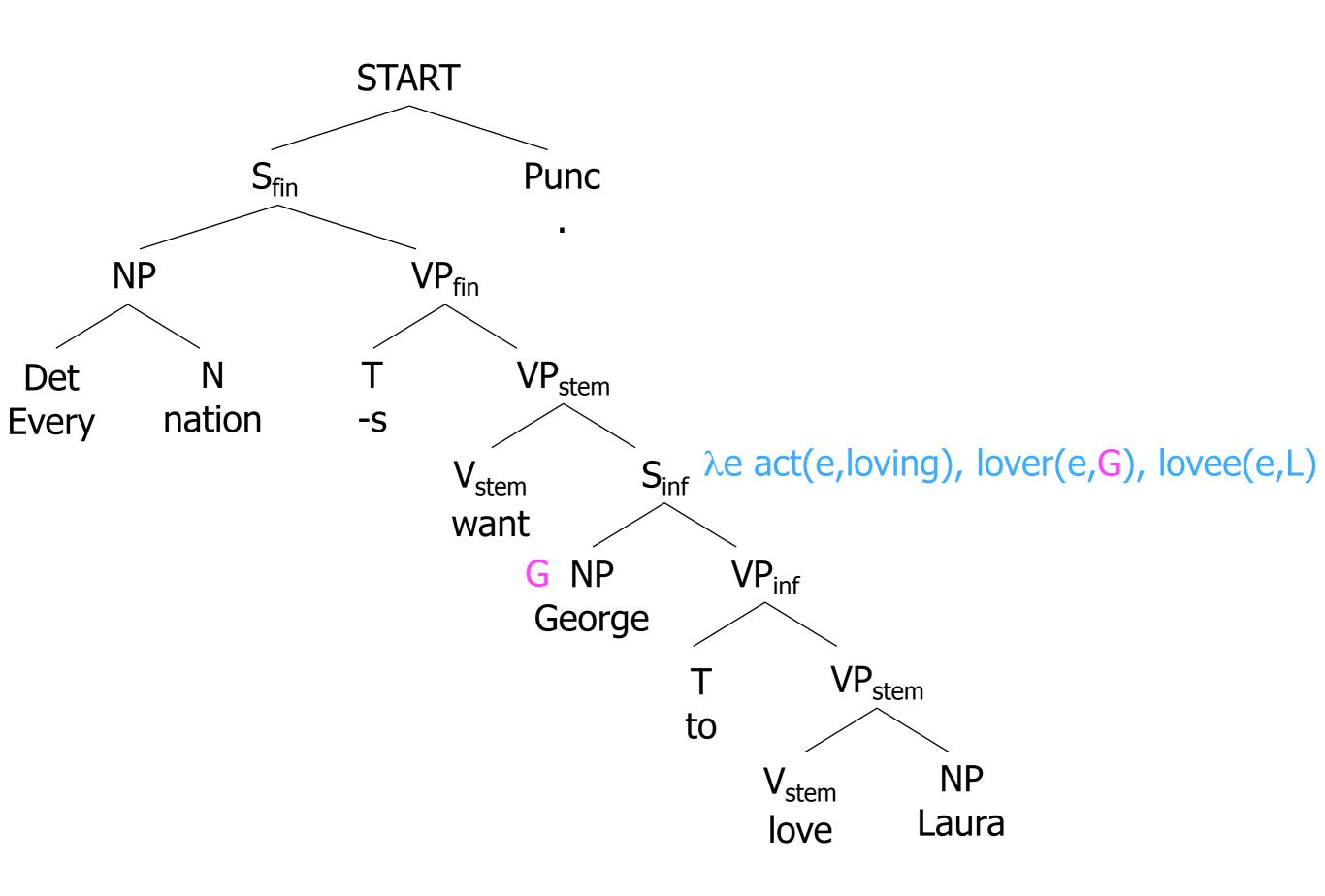


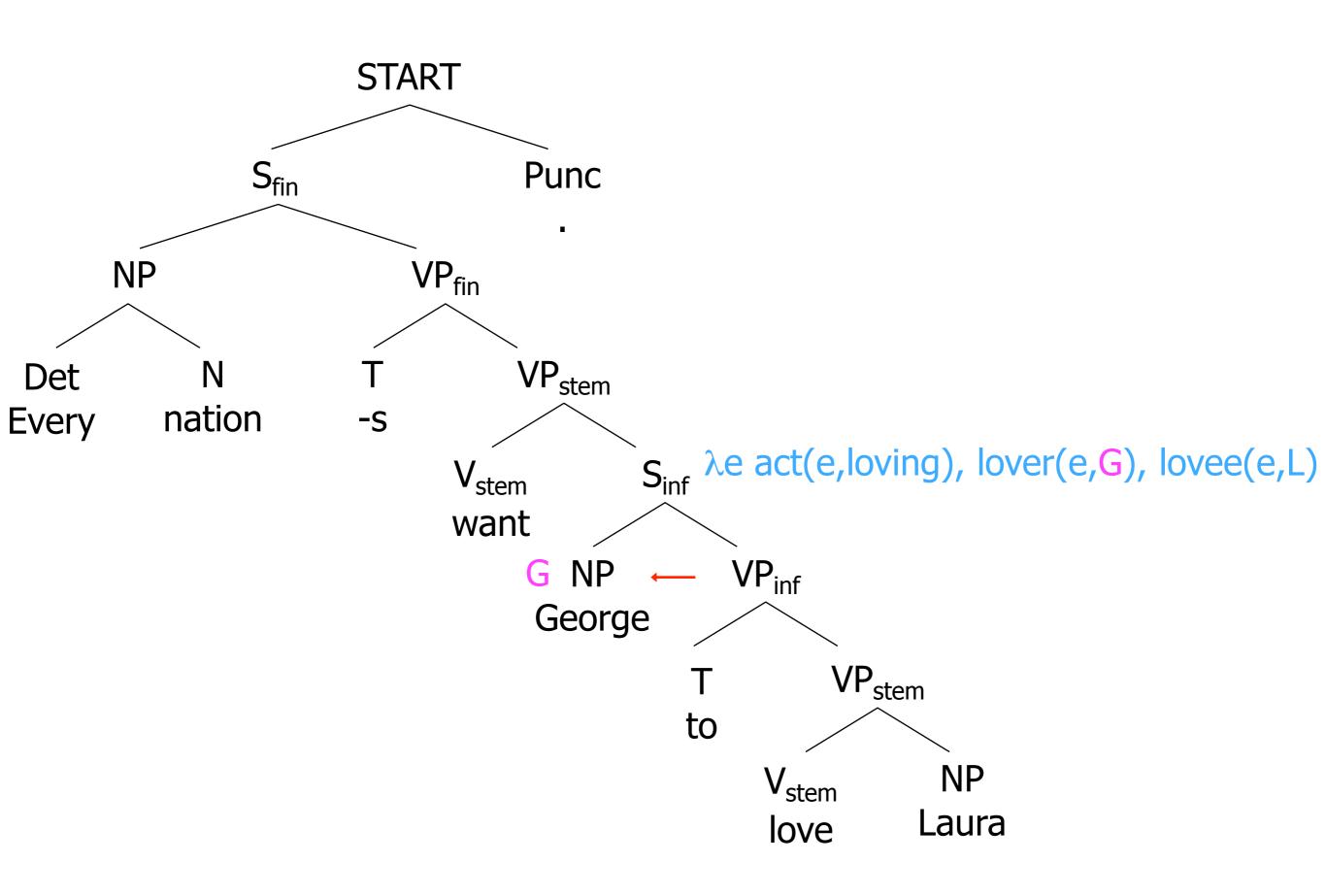


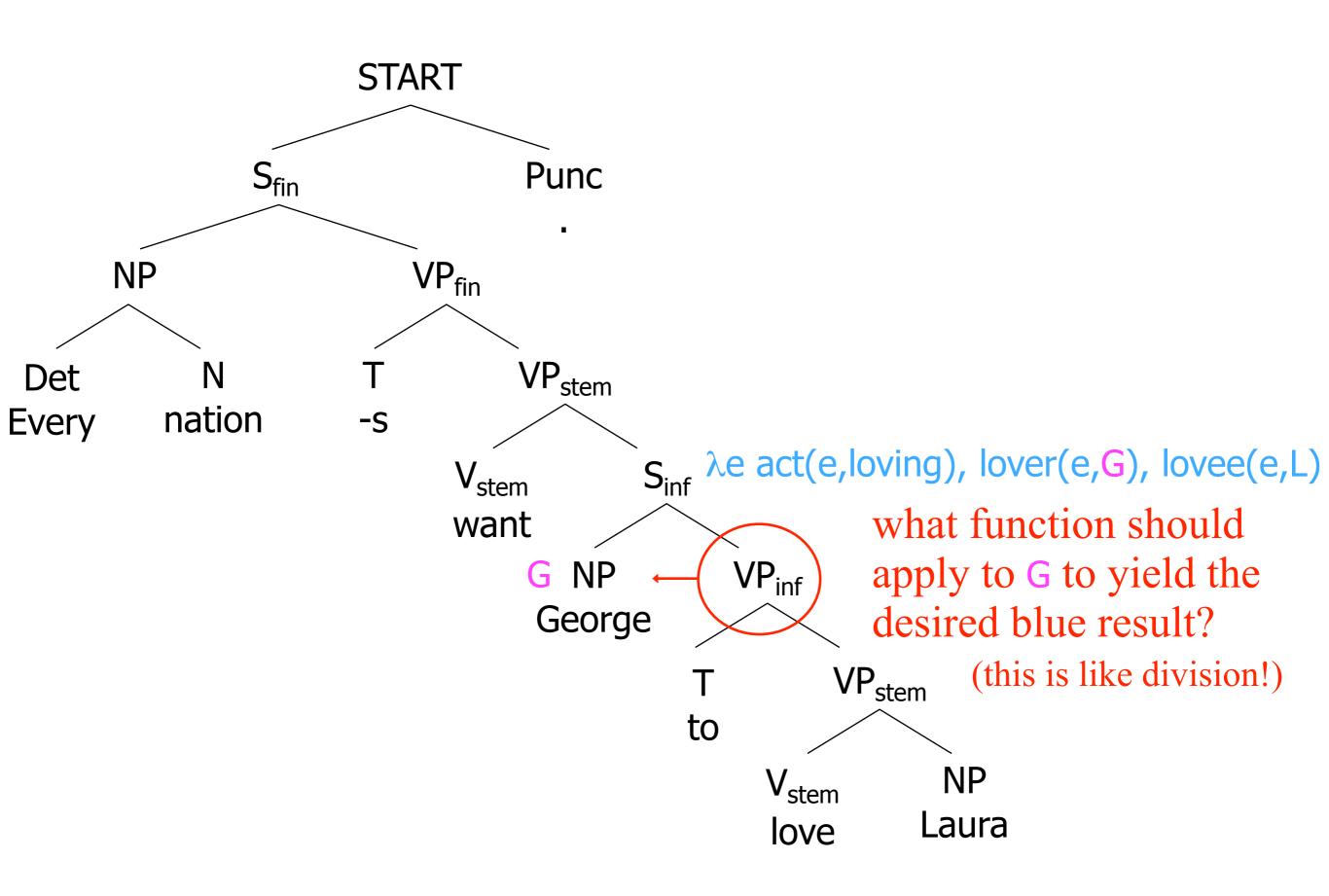


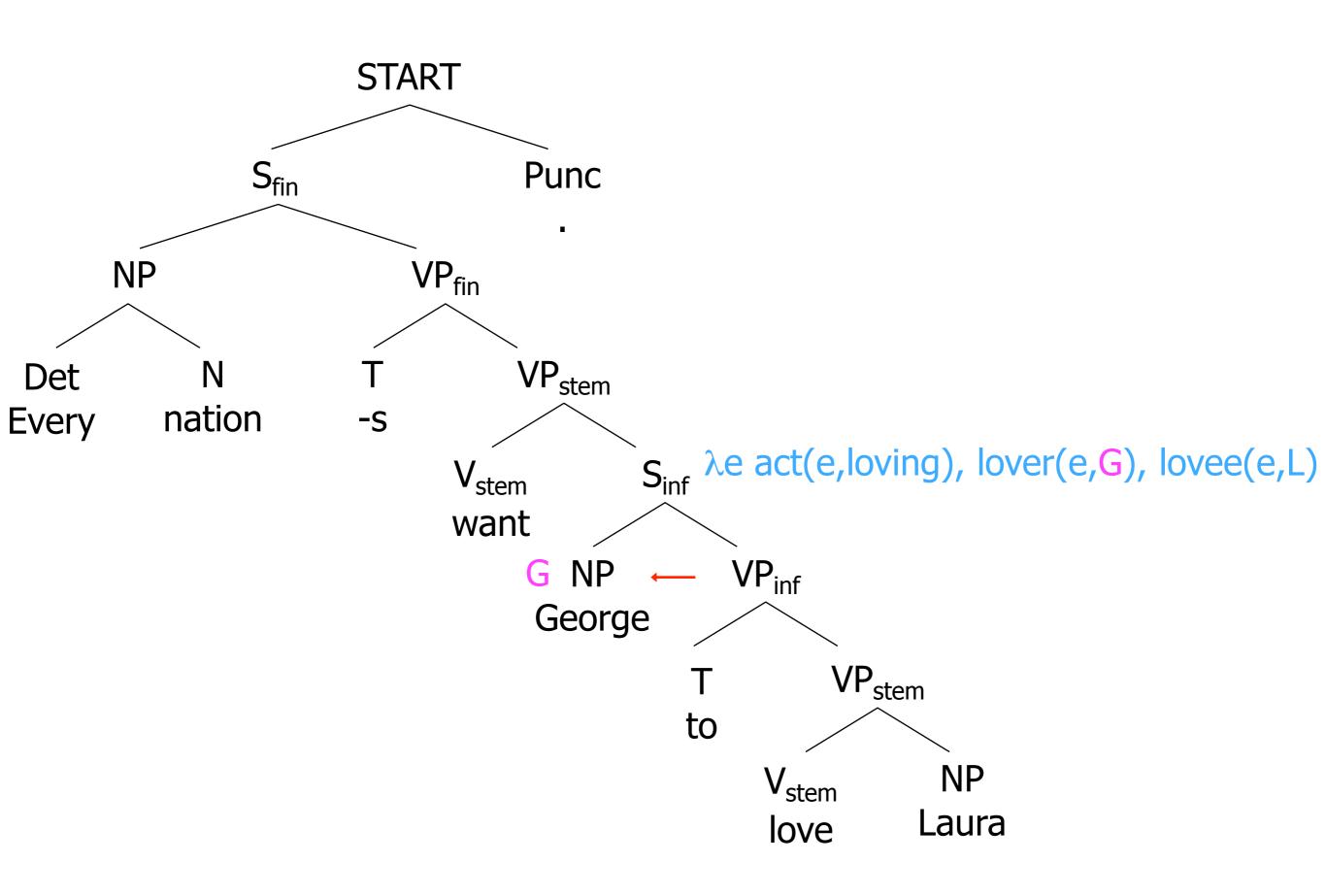


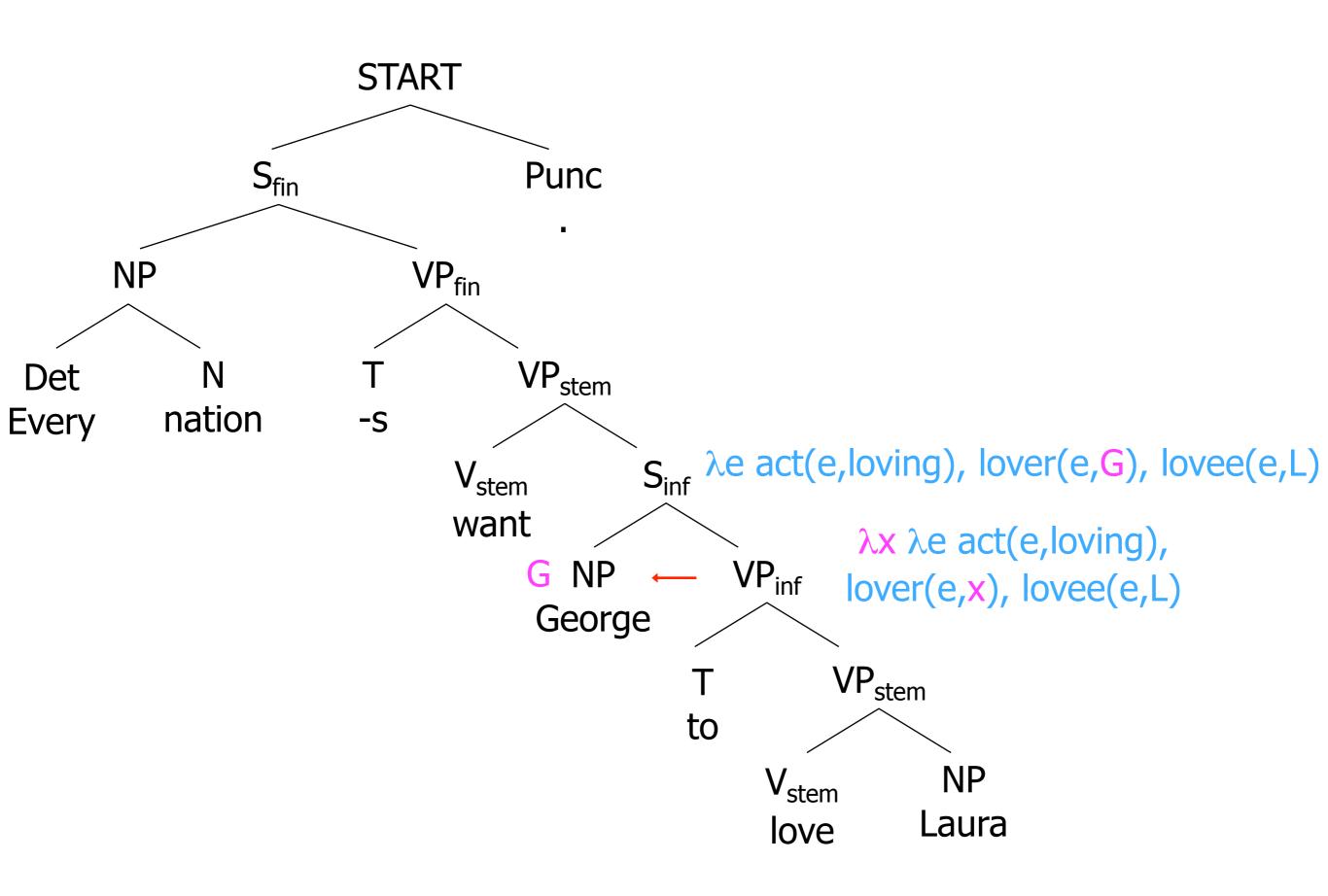


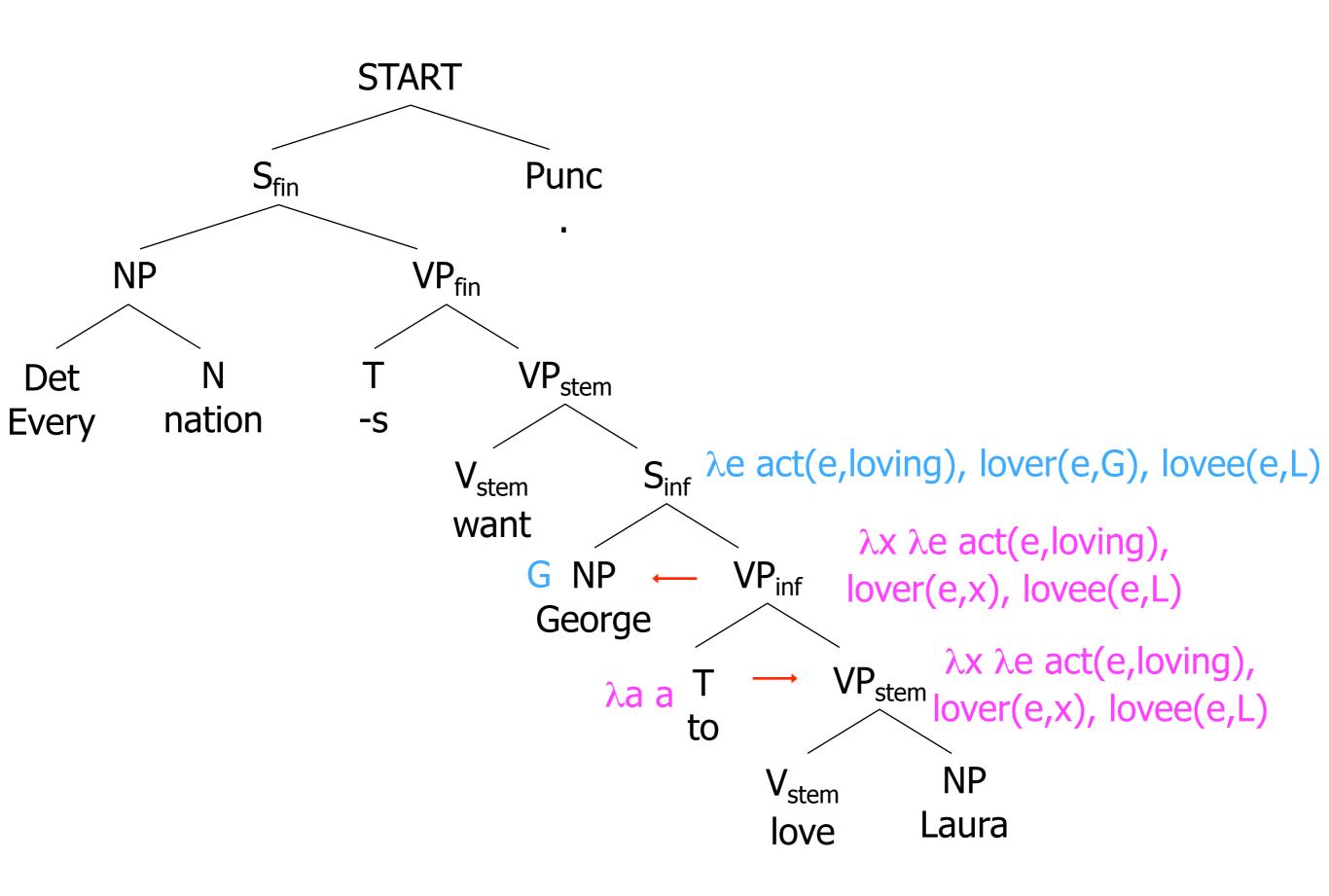


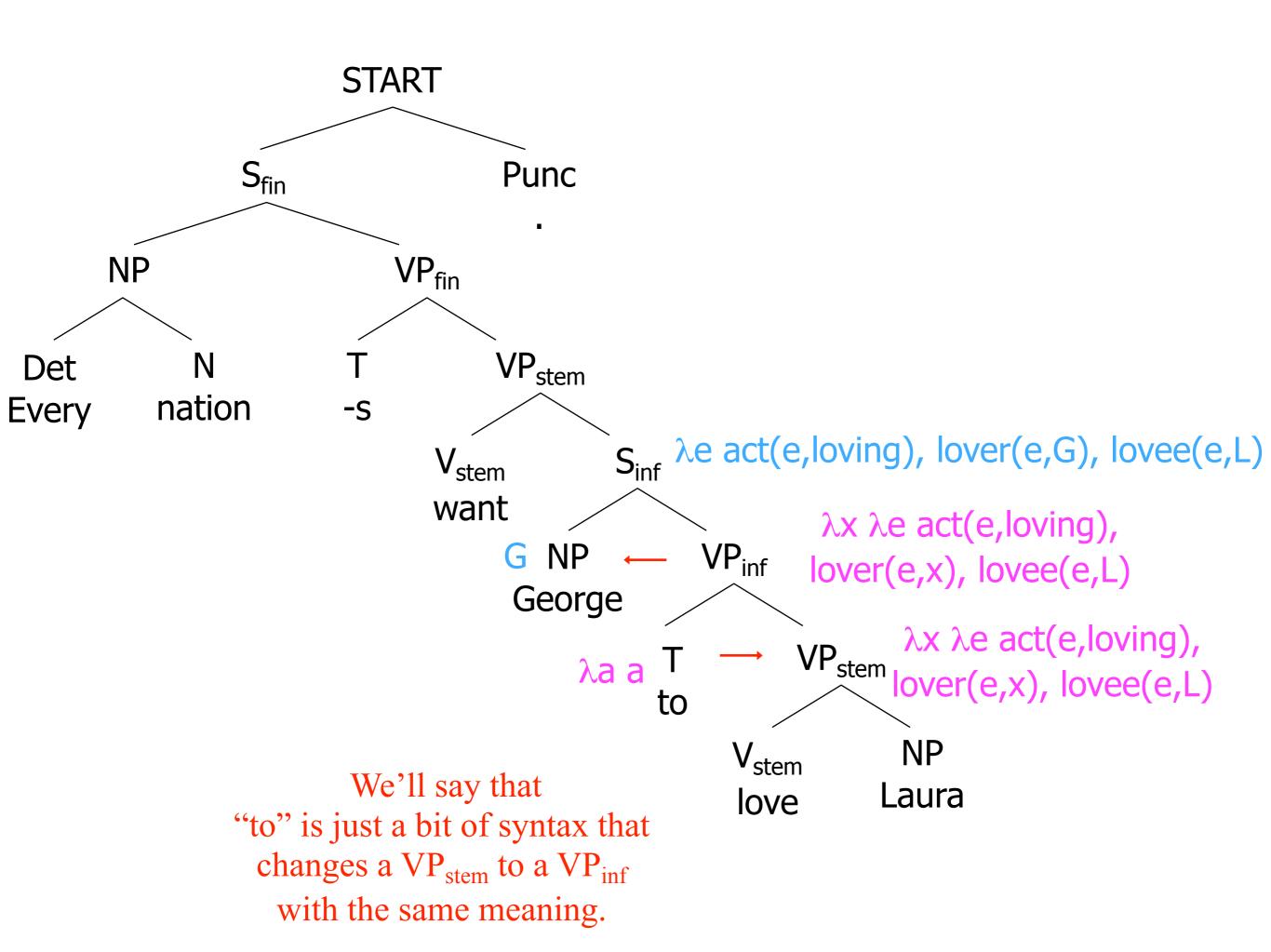


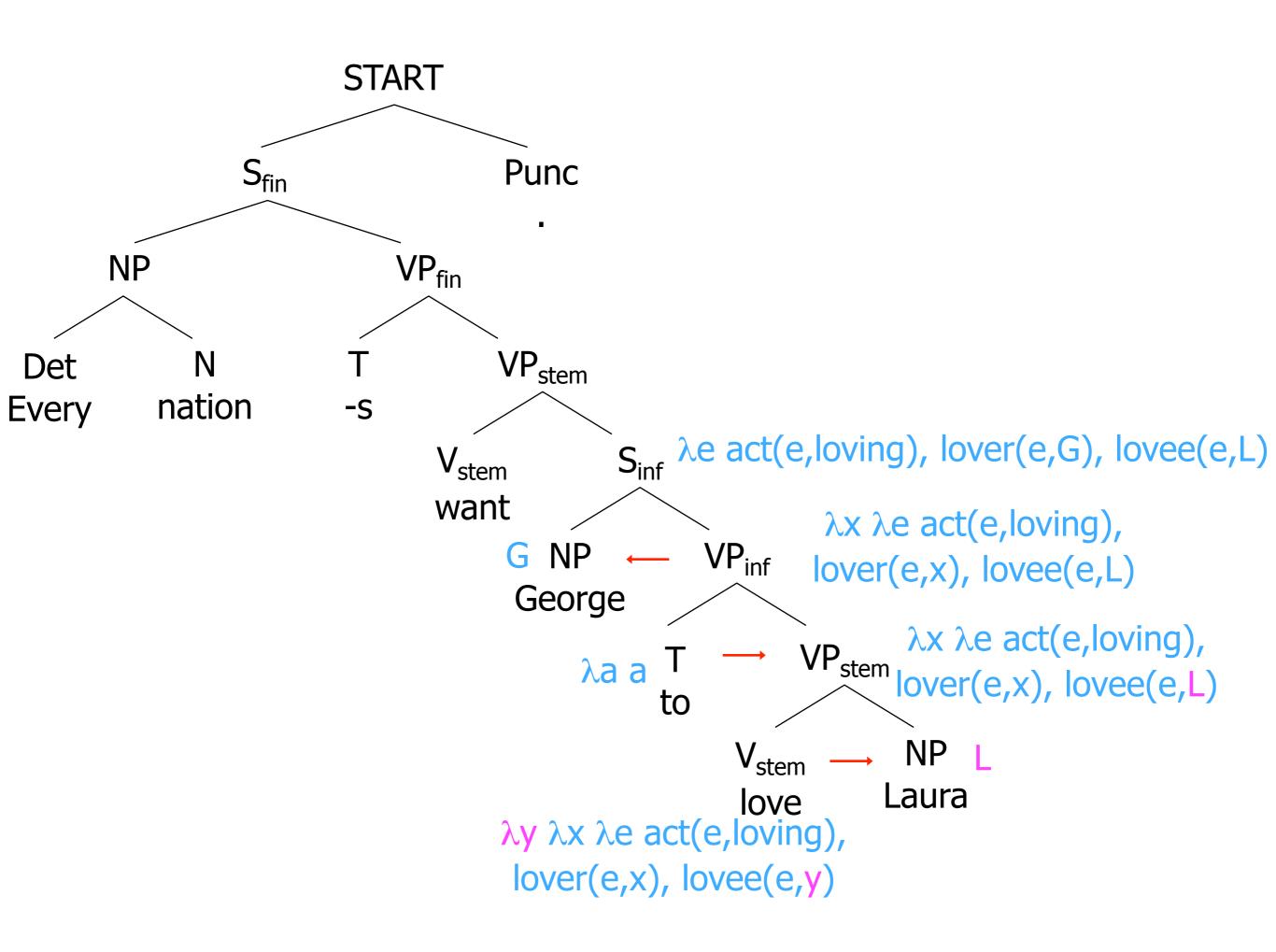


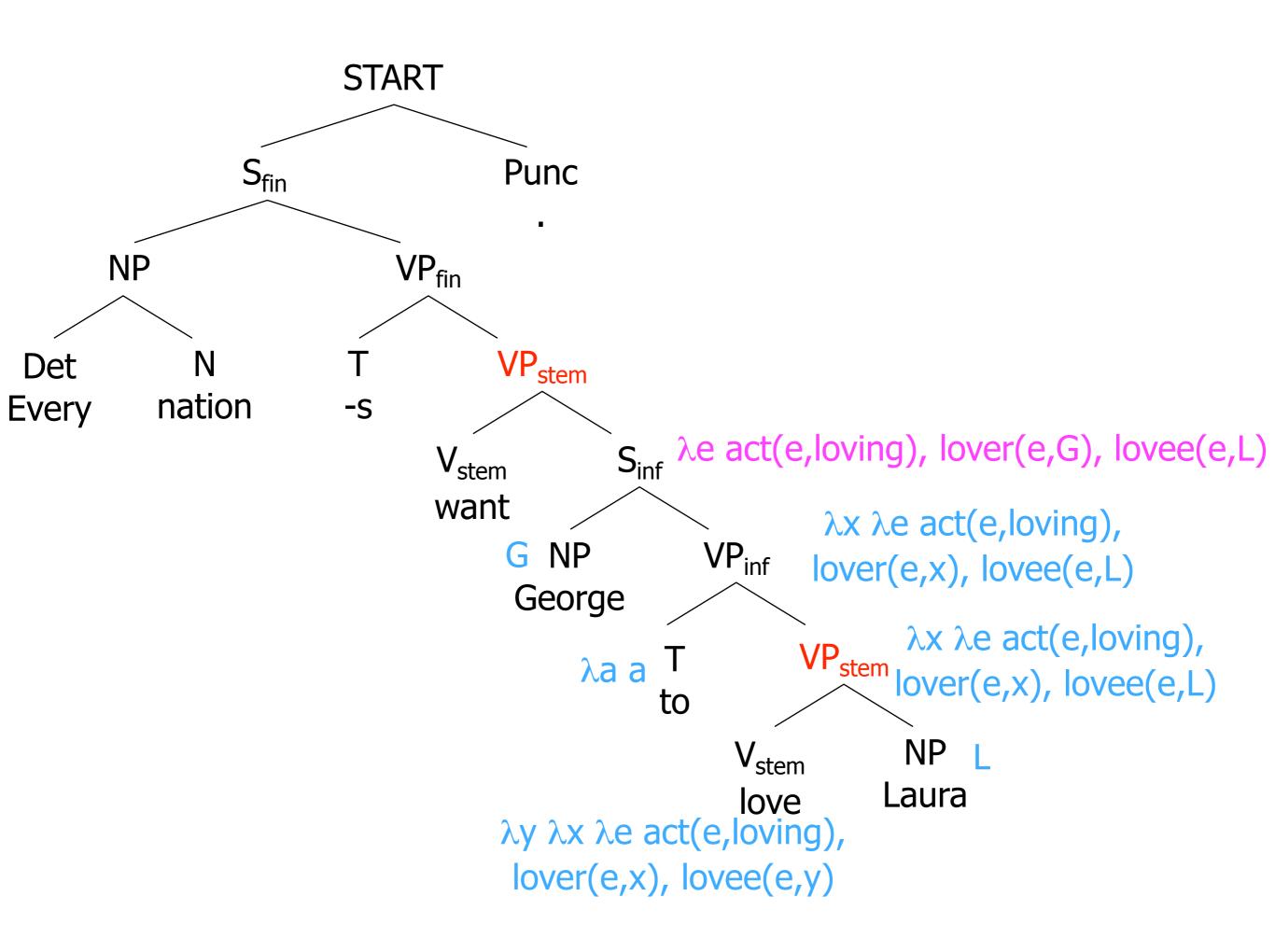


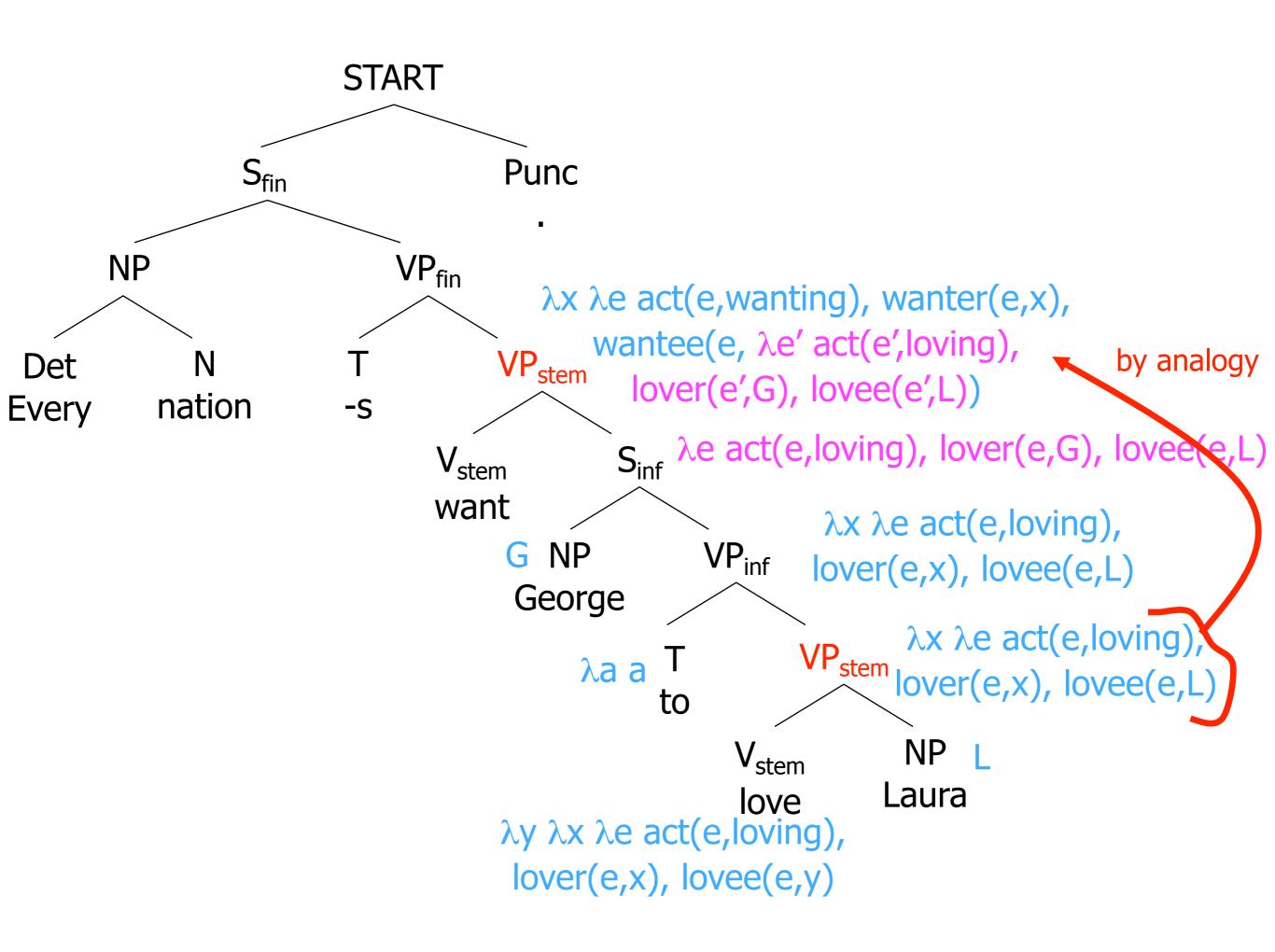


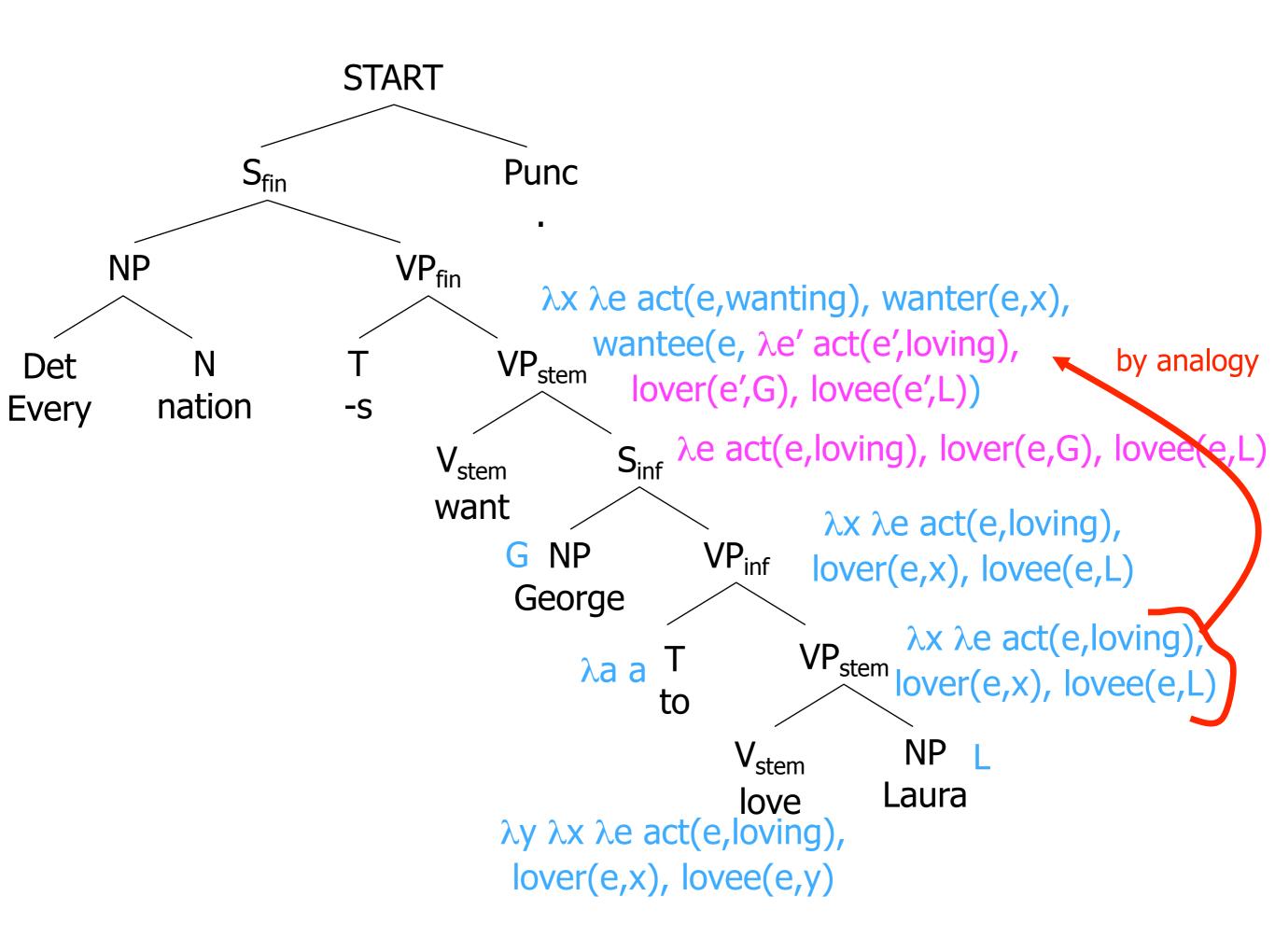


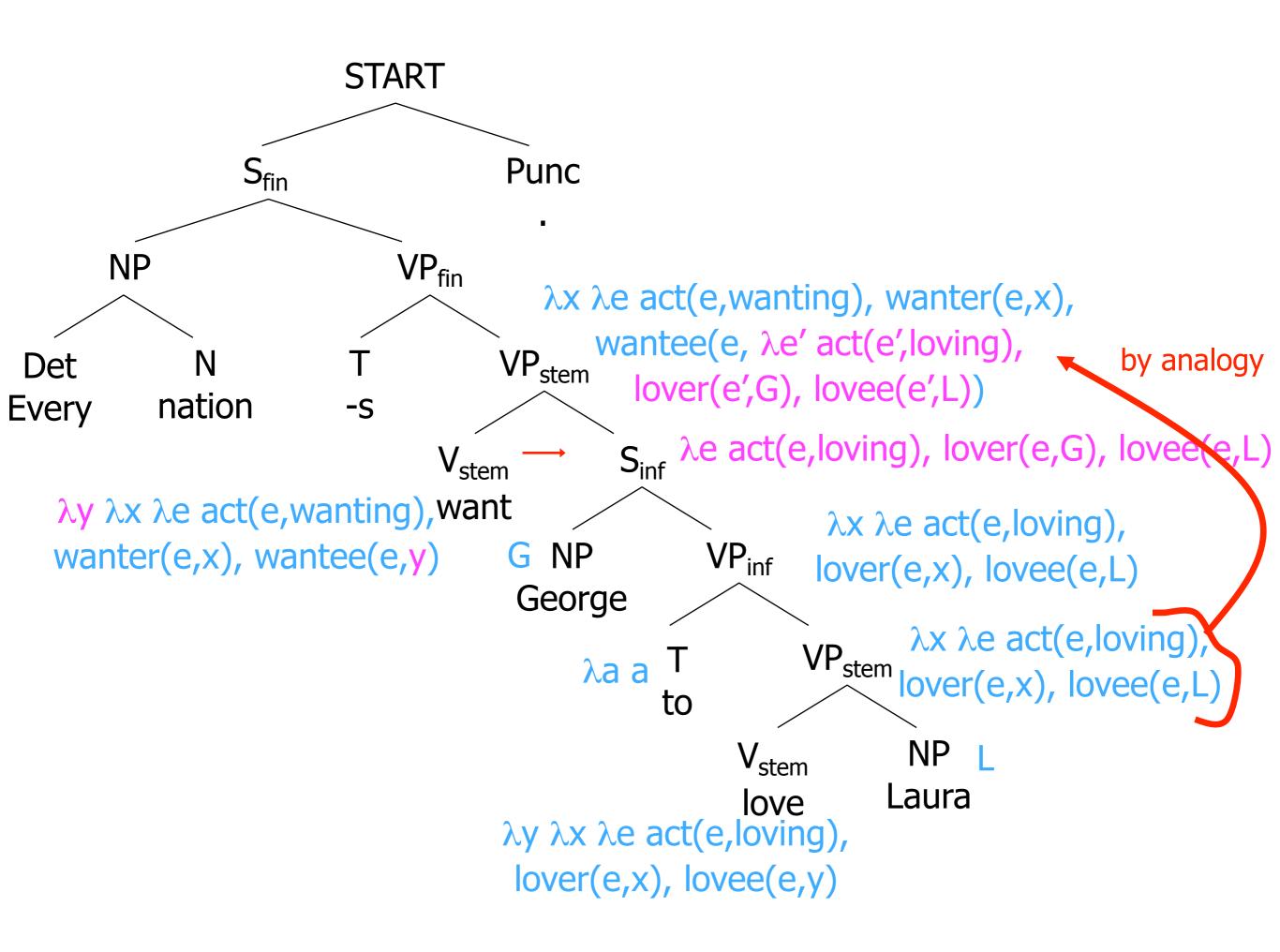


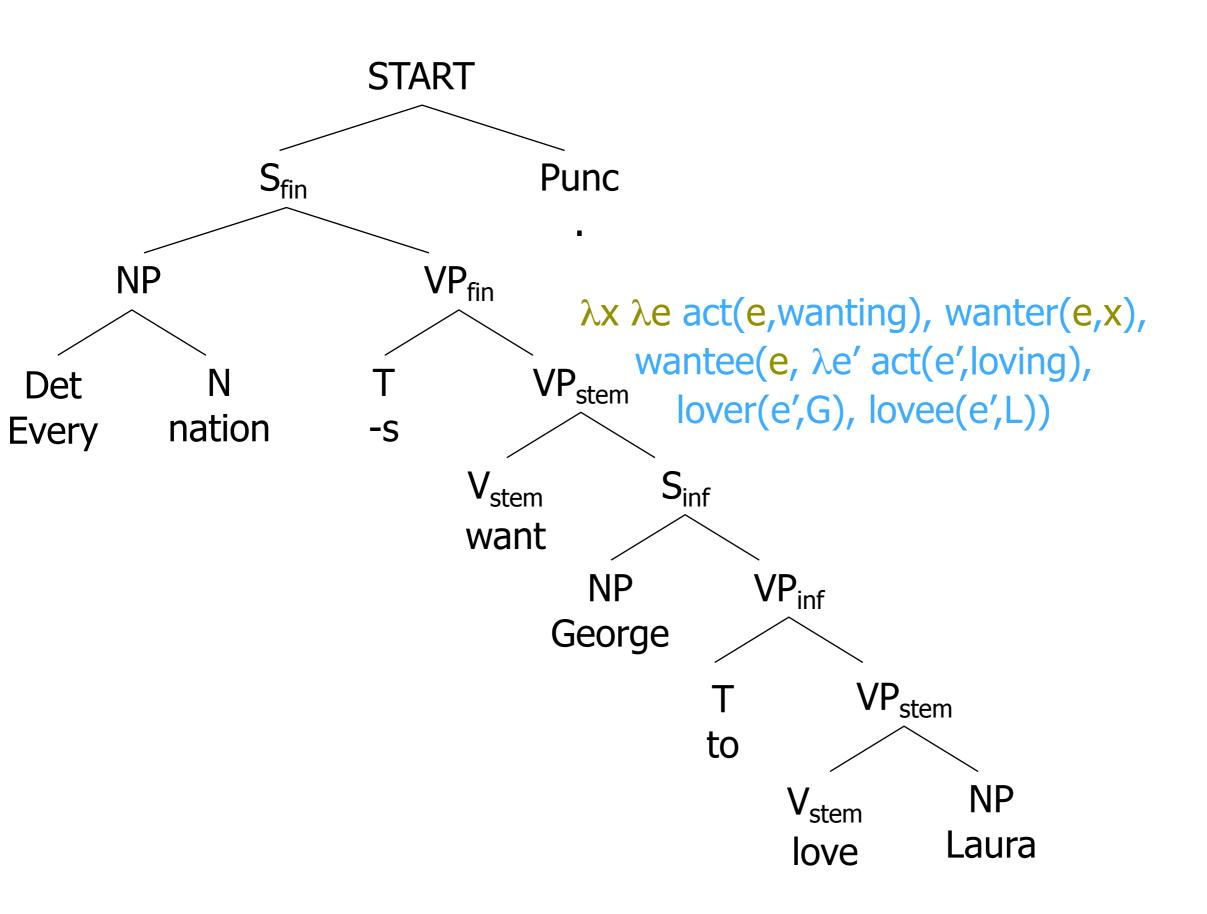


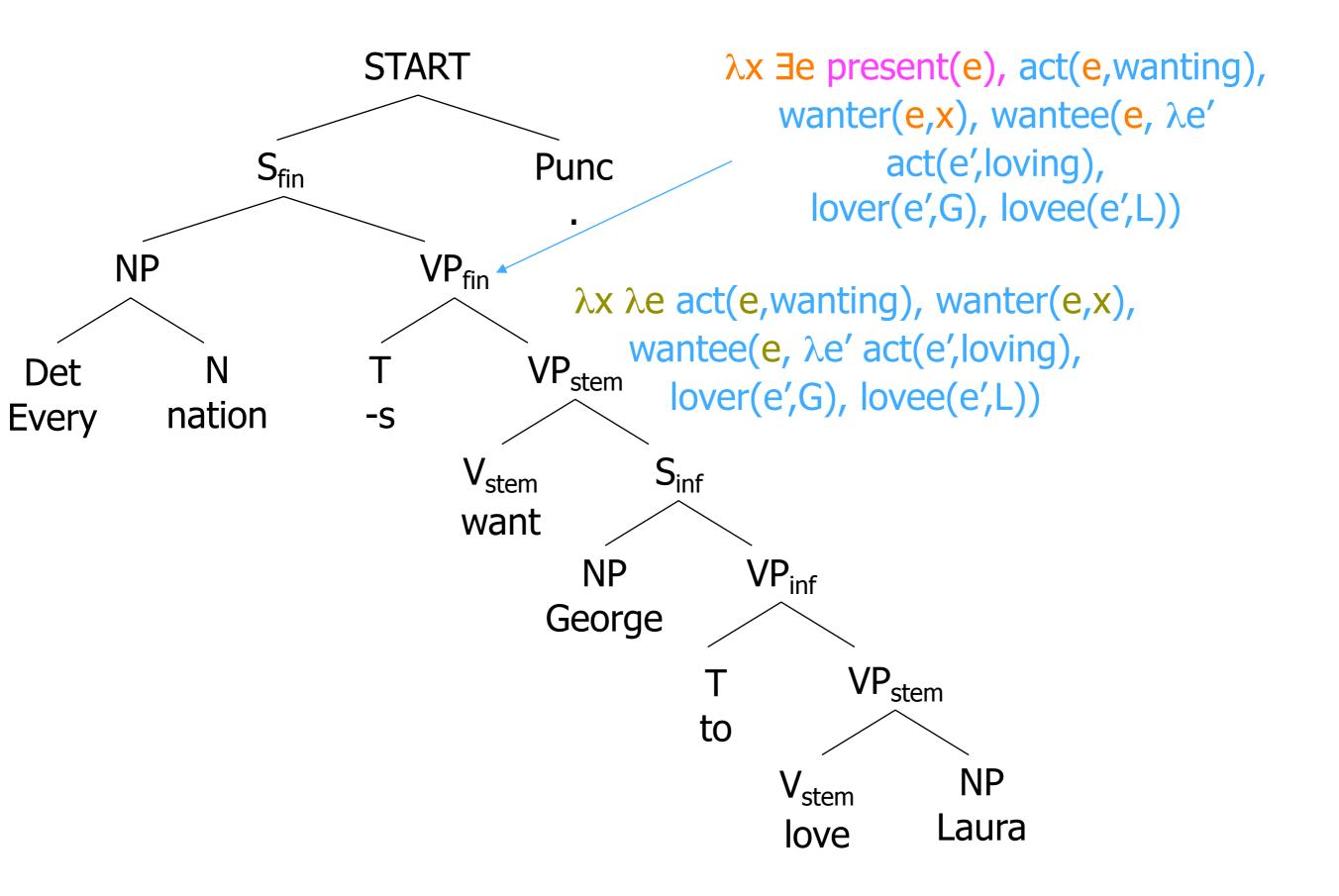


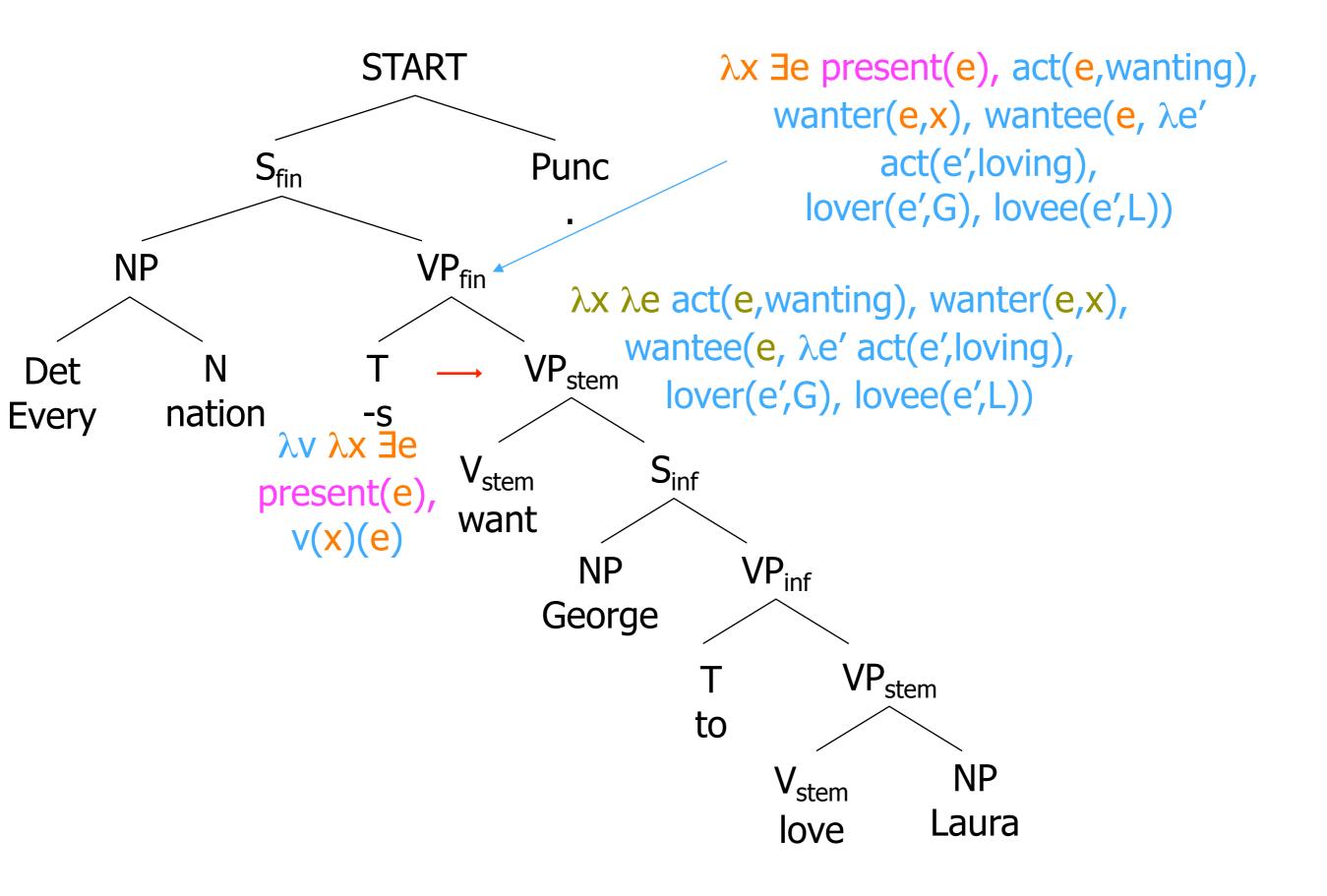


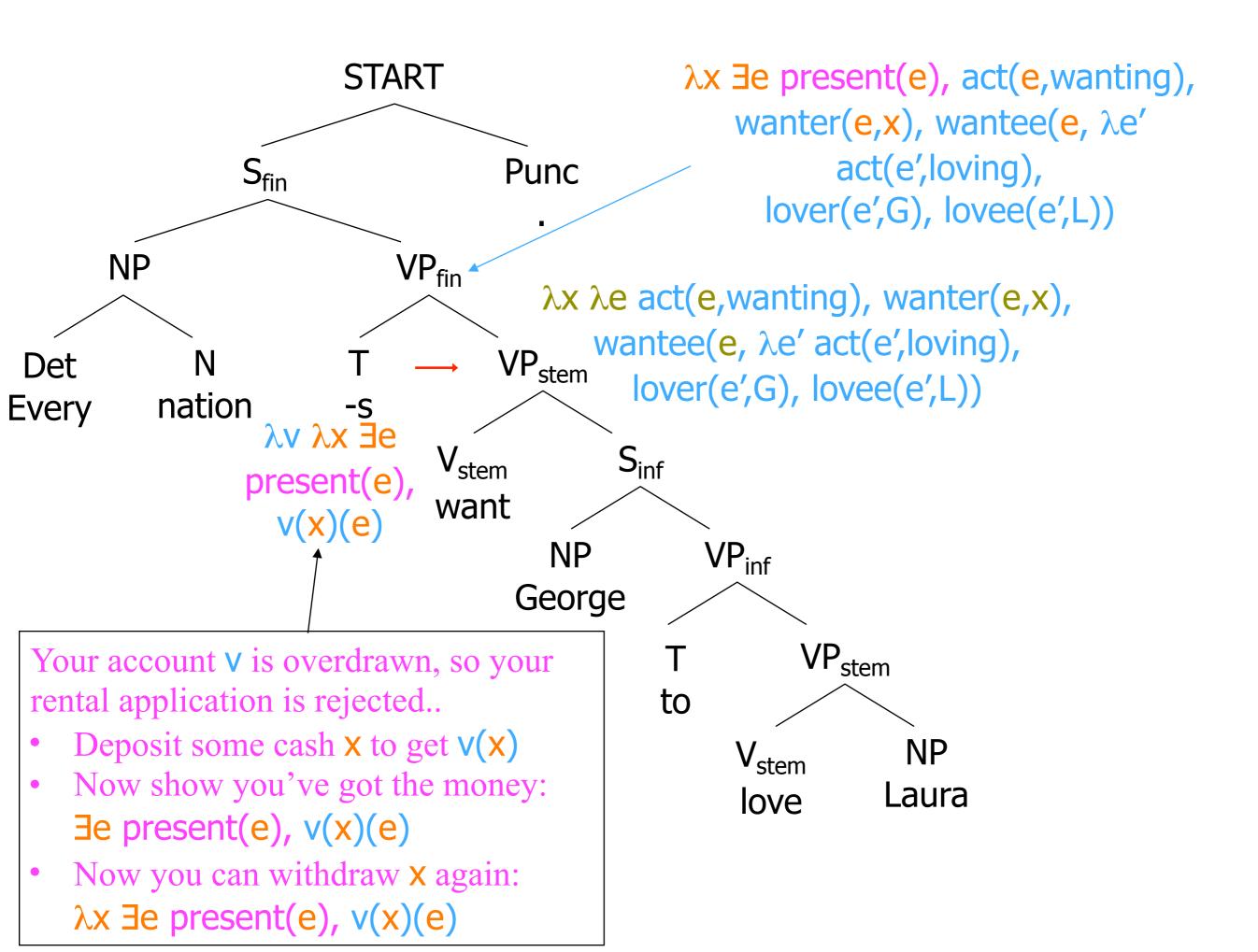


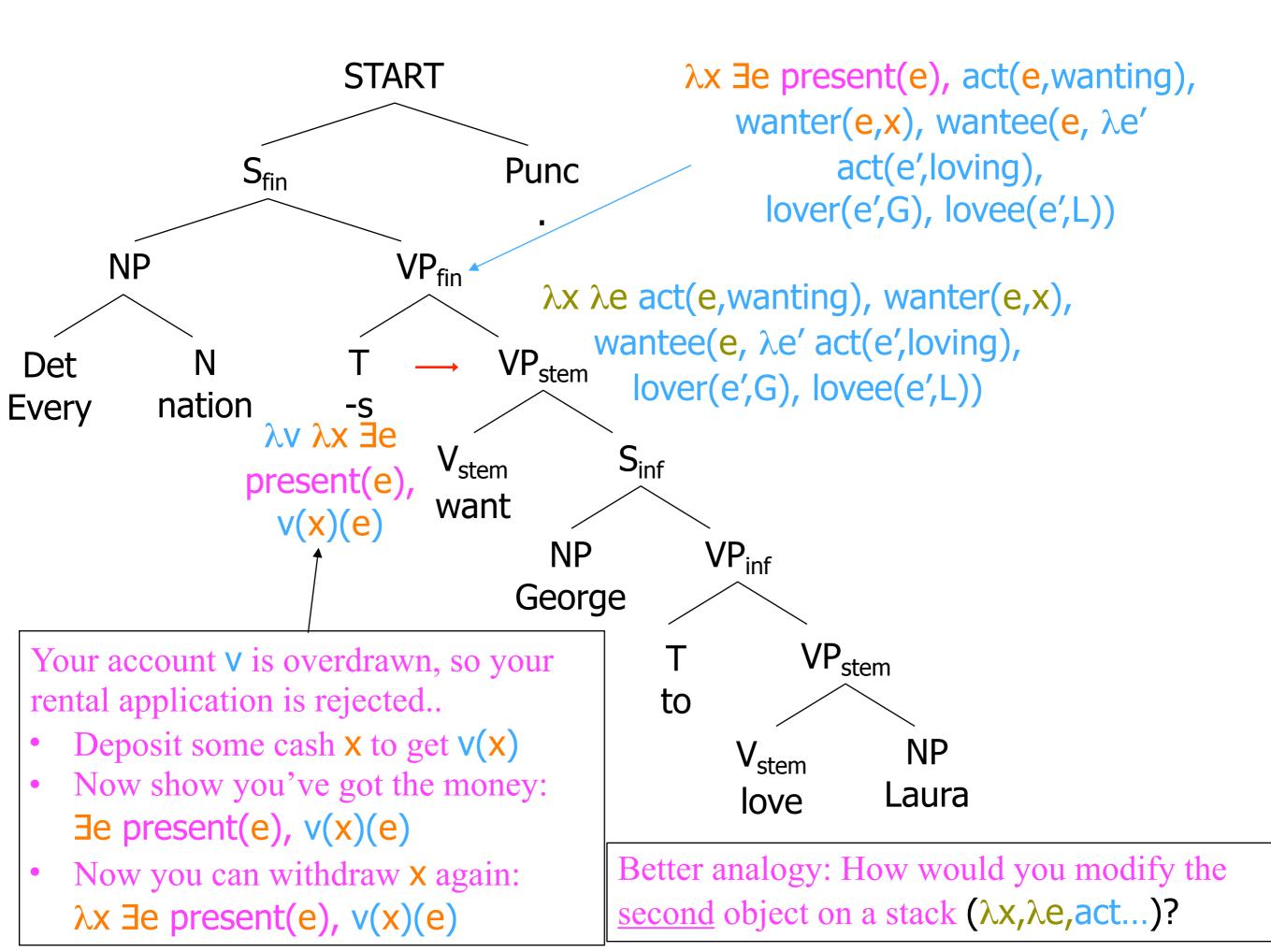


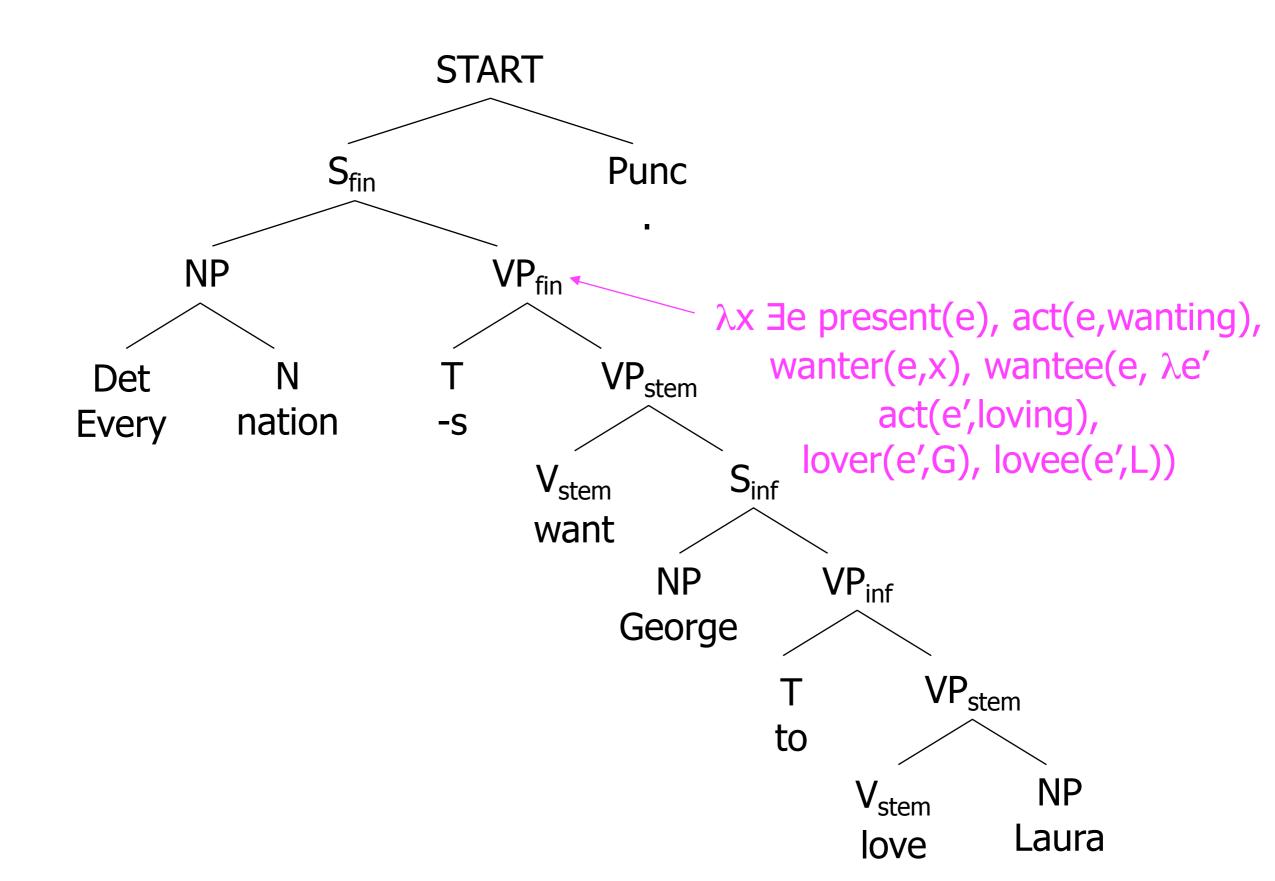


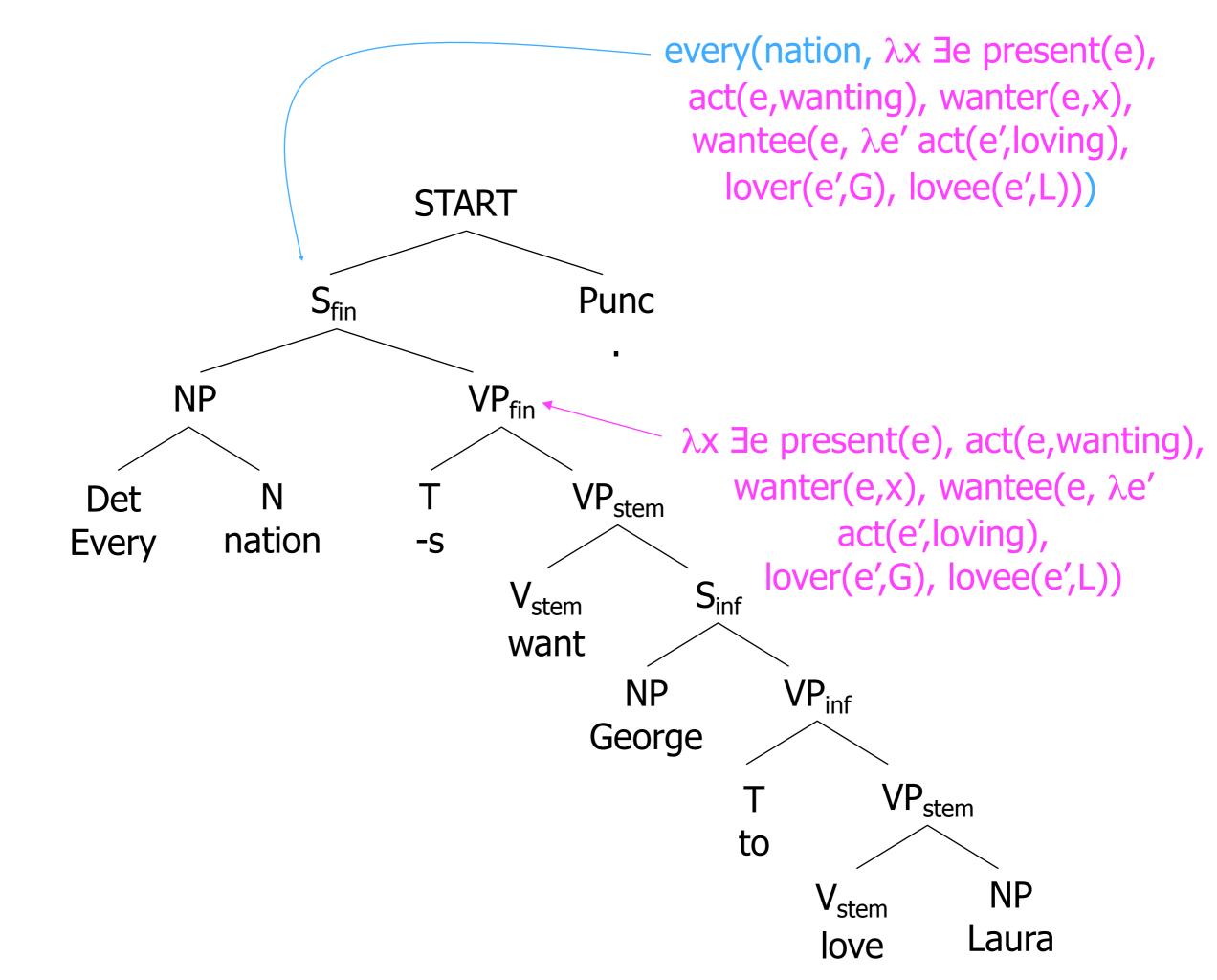


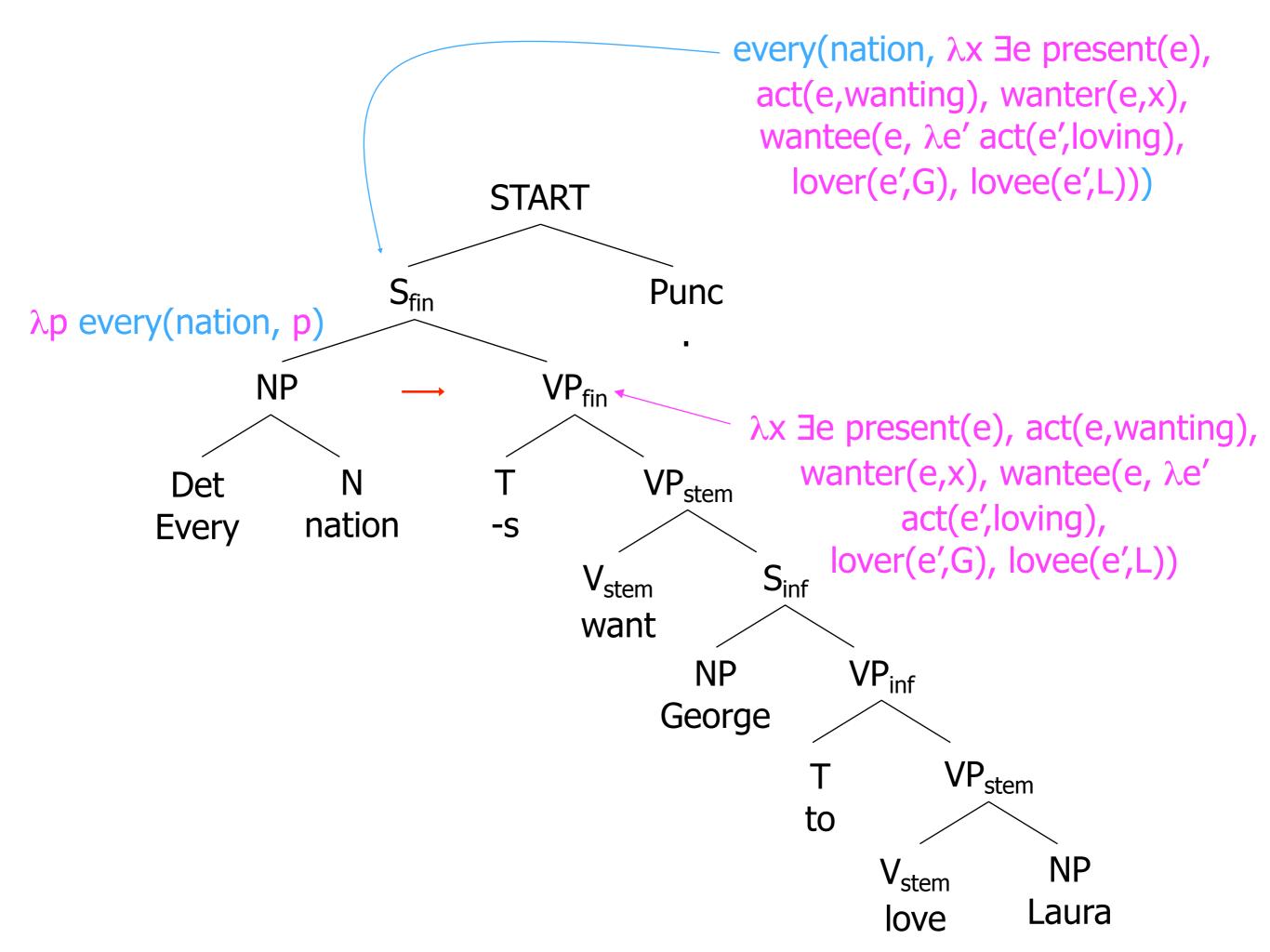


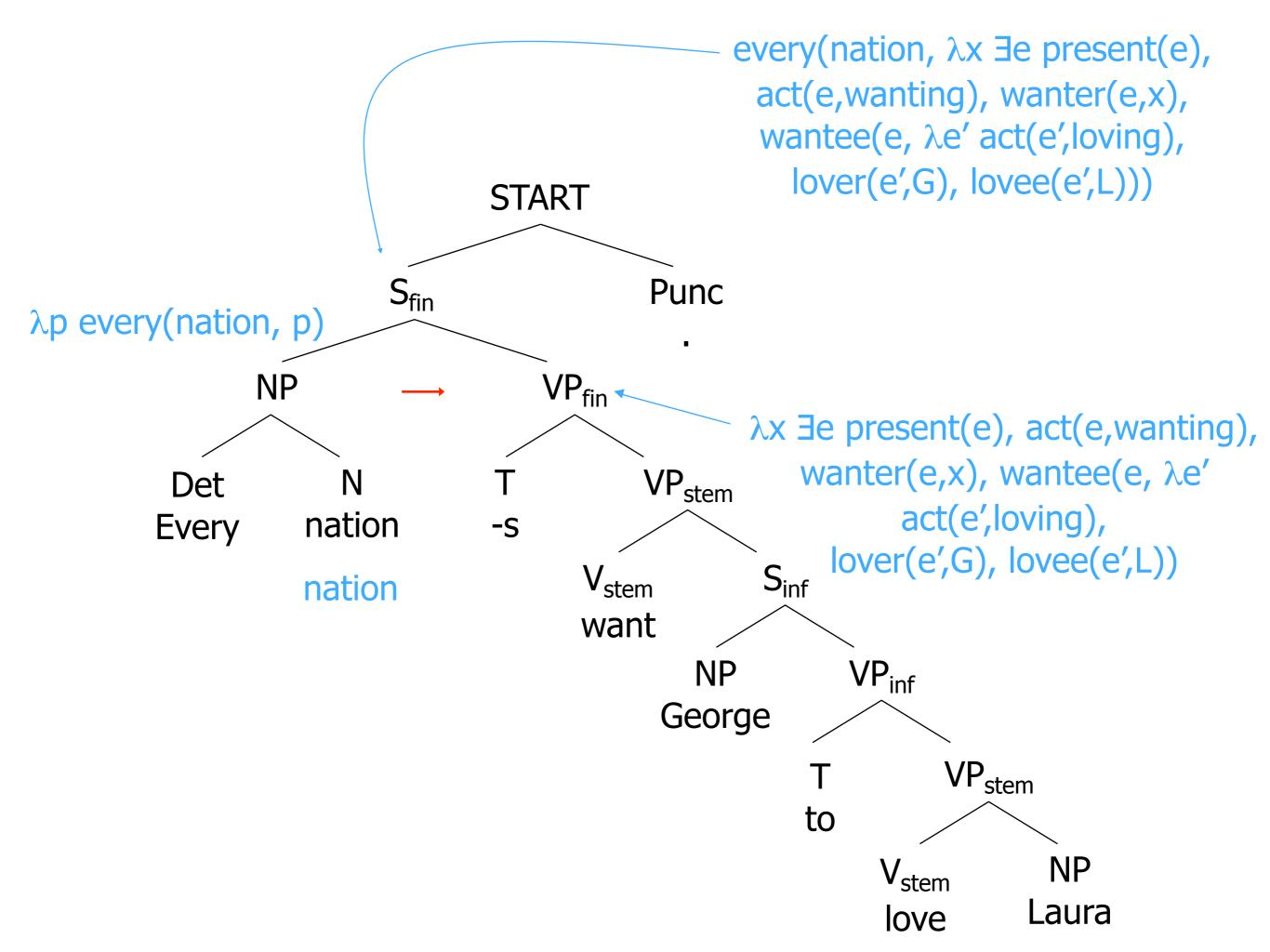


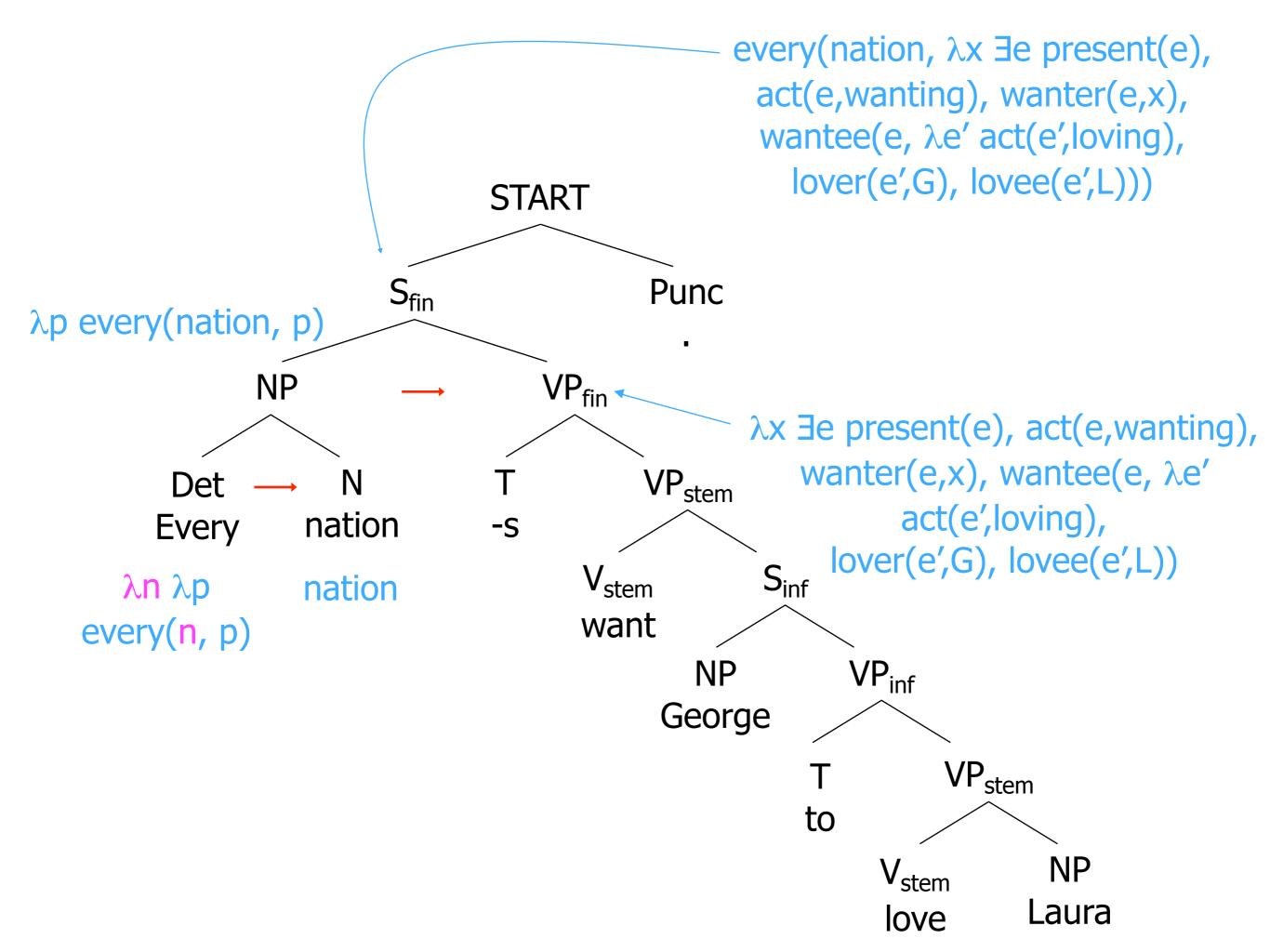


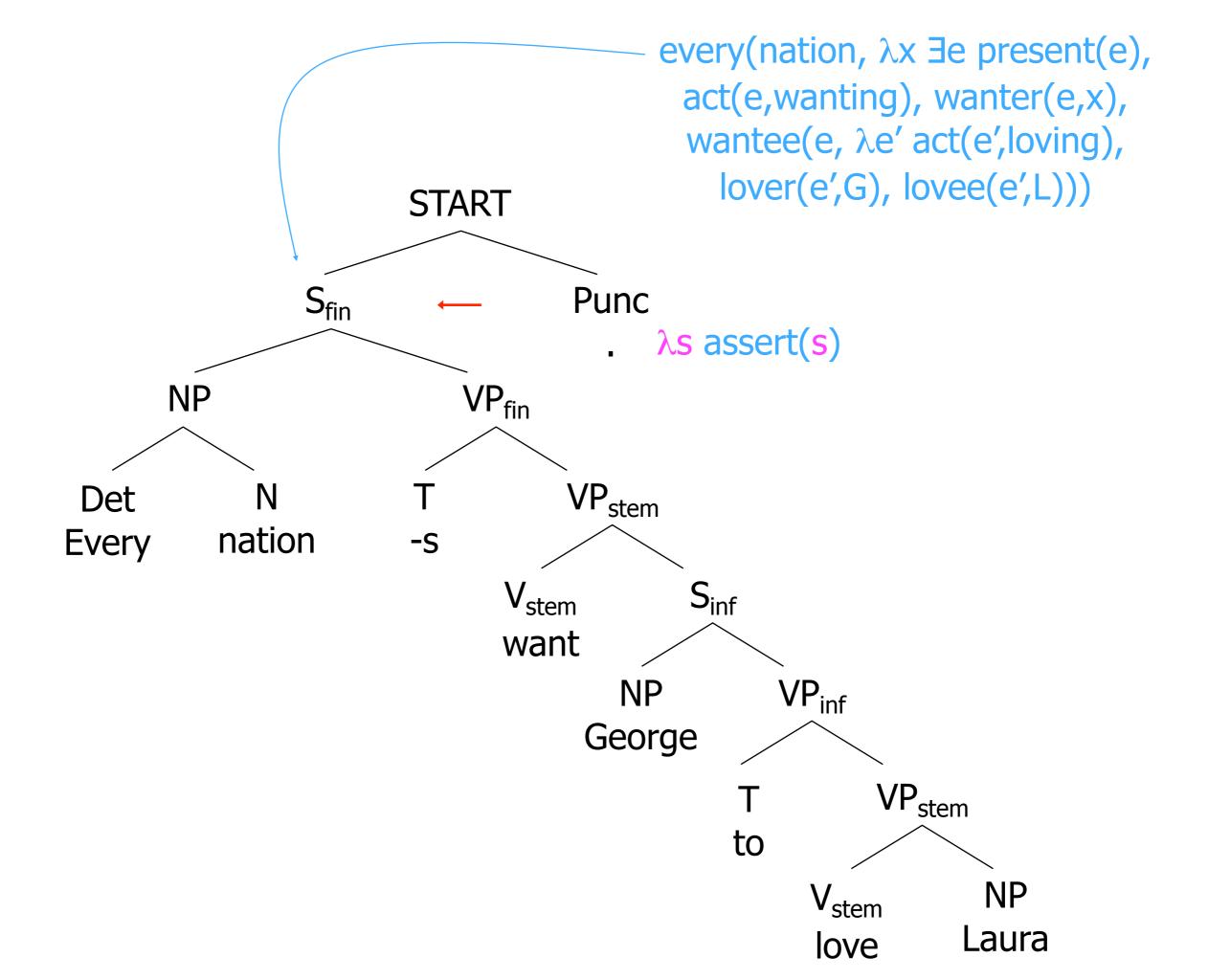




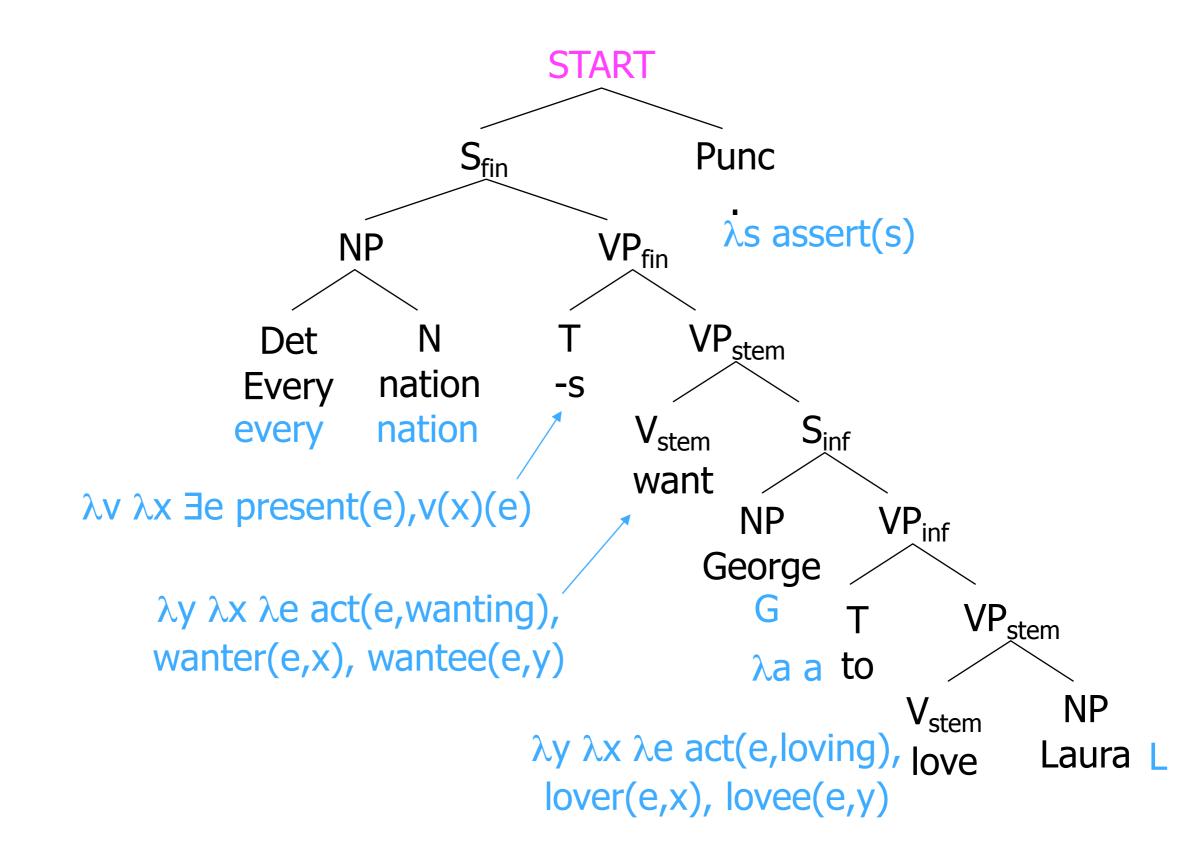




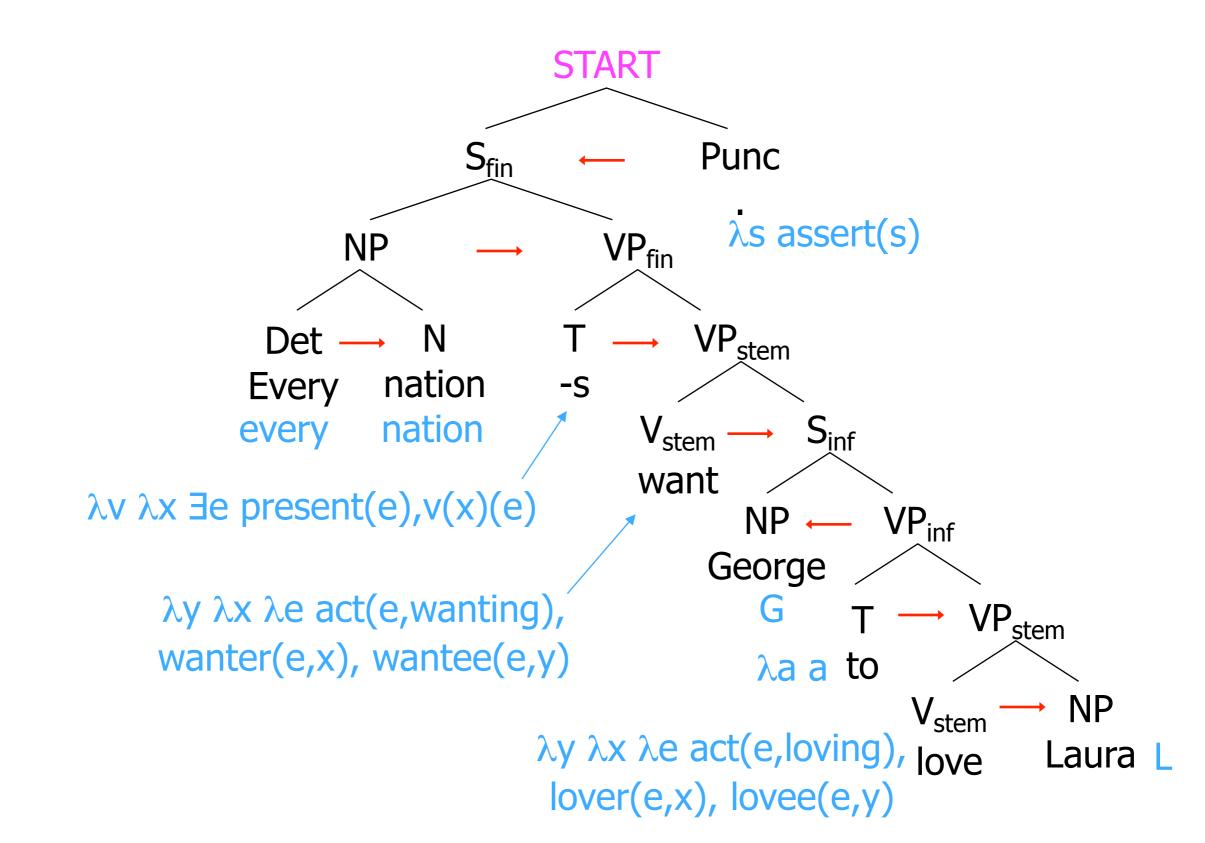




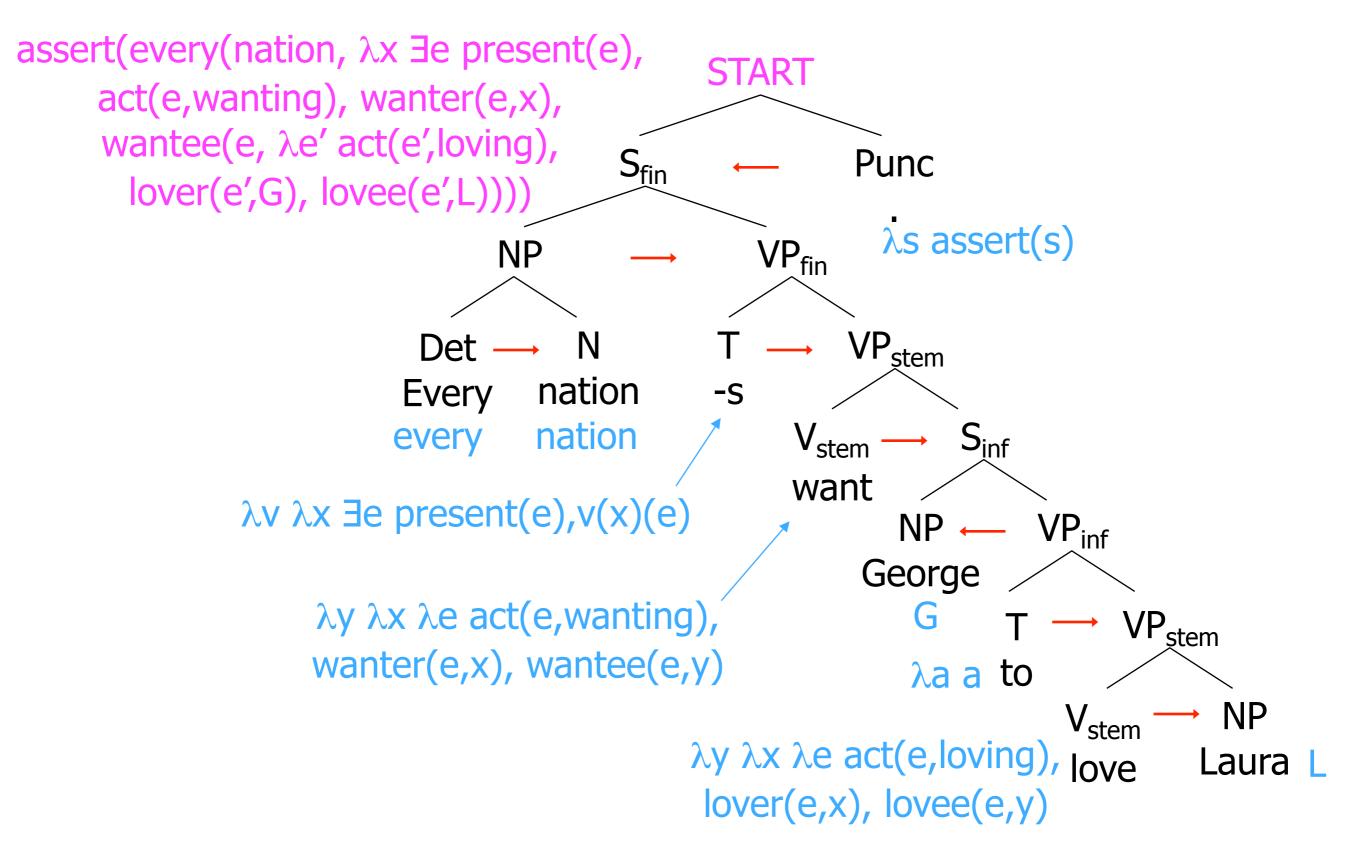
In Summary: From the Words



In Summary: From the Words



In Summary: From the Words



Other Fun Semantic Stuff: A Few Much-Studied Miscellany

Temporal logic

- Gilly <u>had swallowed</u> eight goldfish before Milly reached the bowl
- Billy said Jilly was pregnant
- Billy said, "Jilly <u>is</u> pregnant."

Generics

- Typhoons arise in the Pacific
- Children must be carried

Presuppositions

- The king of France is bald.
- Have you stopped beating your wife?

Pronoun-Quantifier Interaction ("bound anaphora")

- Every farmer who owns a donkey beats <u>it</u>.
- If you have a dime, put <u>it</u> in the meter.
- The woman who every Englishman loves is <u>his</u> mother.
- I love my mother and <u>so</u> does Billy.

In Summary

- How do we judge a good meaning representation?
- How can we represent sentence meaning with first-order logic?
- How can logical representations of sentences be **composed** from logical forms of words?
- Next time: can we train models to recover logical forms?