

Agenda

- 1) Admin → extra BFS/DFS video
- 2) Review
- 3) Searching through Graph
 - Breadth First Search (BFS)
 - Depth First Search (DFS)
- 4) Shortest path in weighted graph
 - Dijkstra's Algorithm

Professor Hamlin
Class 15

Review : Graph: $G = (V, E)$

Vocab: adjacent, degree, walk, path, cycle, connected, subgraphs, weighted, directed,

Trees: parent, children, root, leaf, sibling, ancestors, descendants, forests

On computer: adjacency list, adjacency matrix

Graph isomorphism

Exercise: 1. Draw the following adjacency list as E-rooted tree

A: [B]

B: [A, D]

C: [D]

D: [B, C, E]

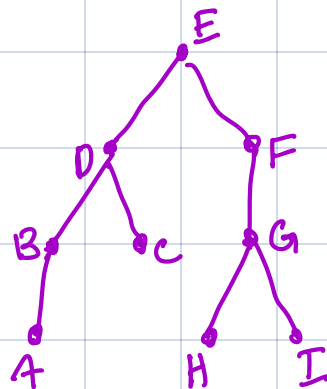
E: [D, F]

F: [G, E]

G: [F, H, I]

H: [G]

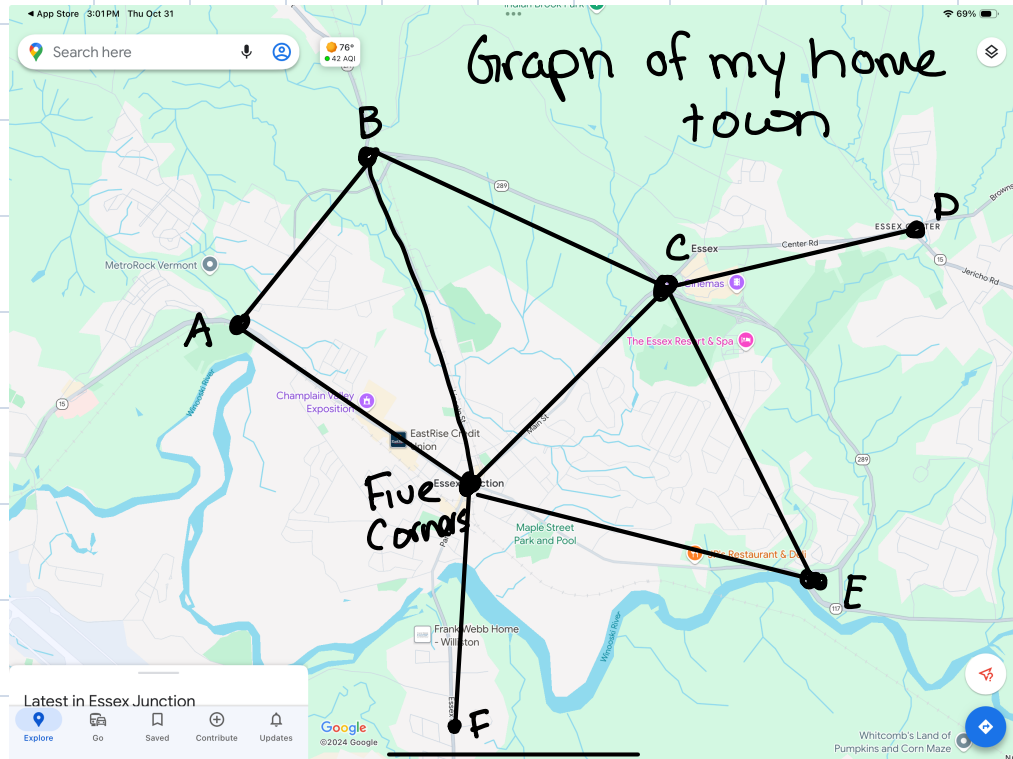
I: [G]



2. In the previous graph, what is the path between A & G?

A, B, D, E, F, G

Searching a graph



As a teen could go anywhere within 10 minutes of 5 corners. What could I get to?

A, B, C, E, F

D is "too far away"

This is easy for a human to do but harder for a computer.

Need to describe a method that works for all graphs: an algorithm

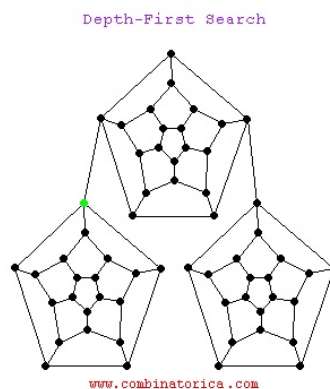
First algorithm...

Depth First Search (DFS)

(The brave algorithm)

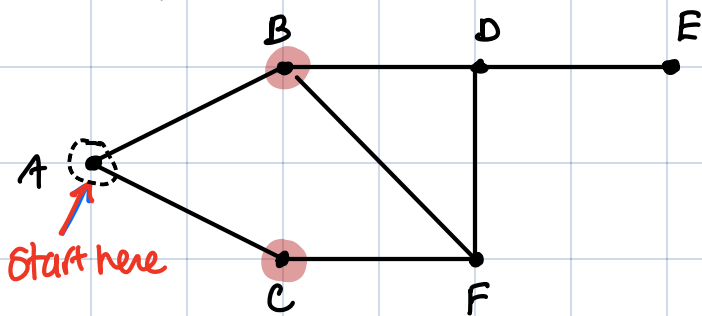
Intuition: keep going far away ^(via adjacent nodes) as possible before backing up and trying a different path.

<https://www3.cs.stonybrook.edu/~skiena/combinatorica/animations/search.html>



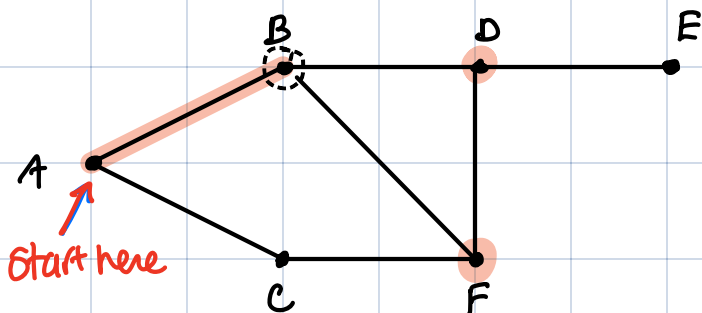
Formally, visit adjacent, unvisited vertices as long as possible, then back up one edge, look for another unvisited vertex to visit using same method.

Example



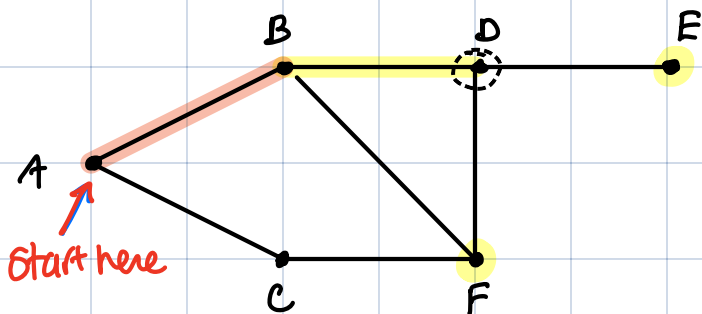
Visited
A

2 neighbours - B, C. Choose 1st alphabetically



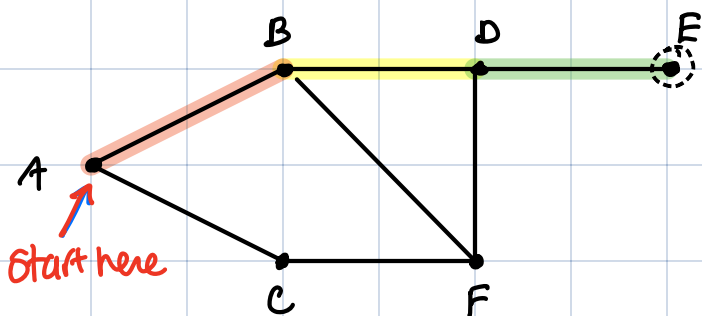
Visited
A, B

2 neighbours - D, F. Choose D



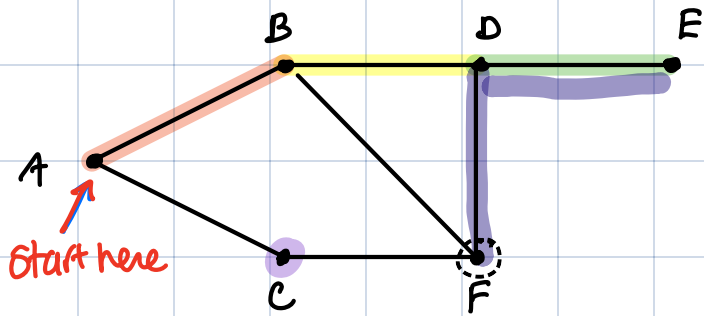
Visited
A, B, D

2 neighbours - E, F. Choose E



Visited
A, B, D, E

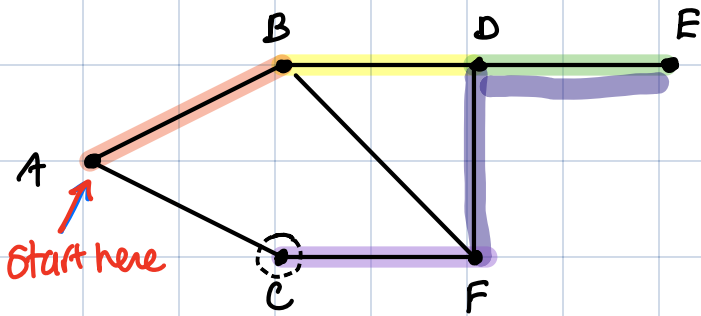
No unvisited neighbors, back-track to D,
F unvisited



Visited

A, B, D, E, F

Two neighbors: B, C, only C unvisited

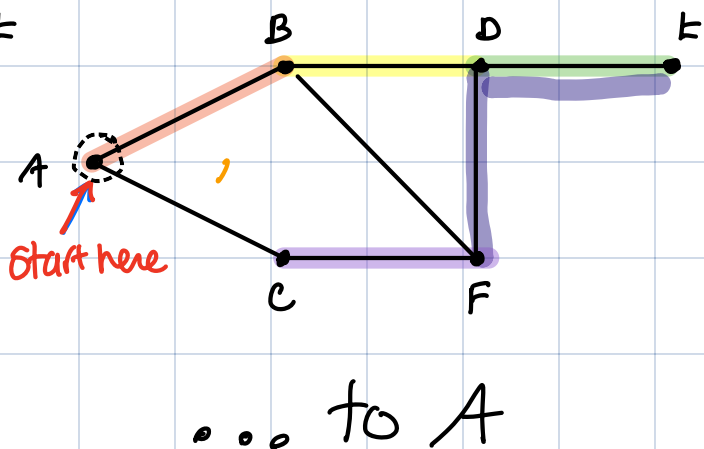
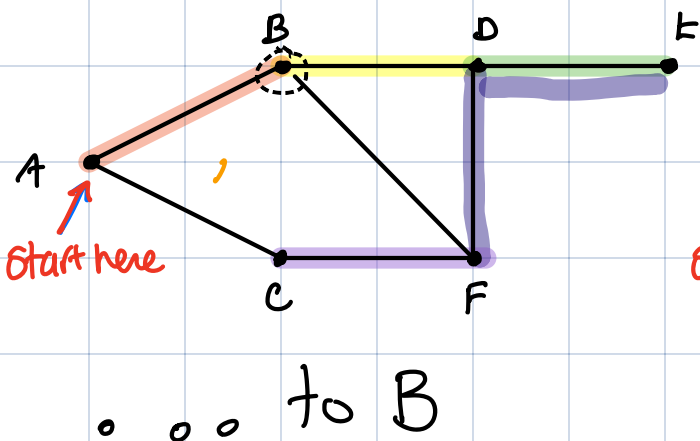
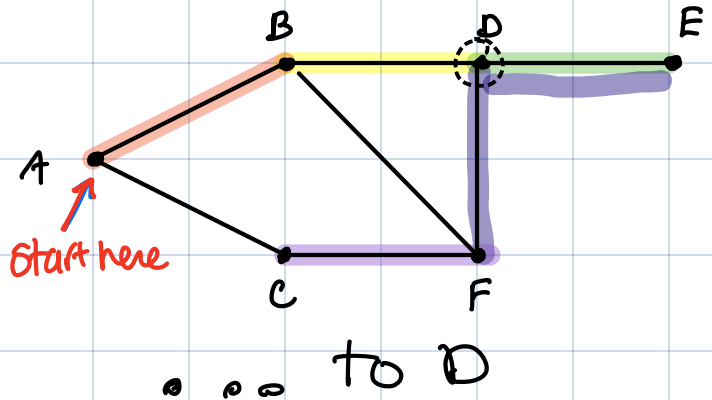
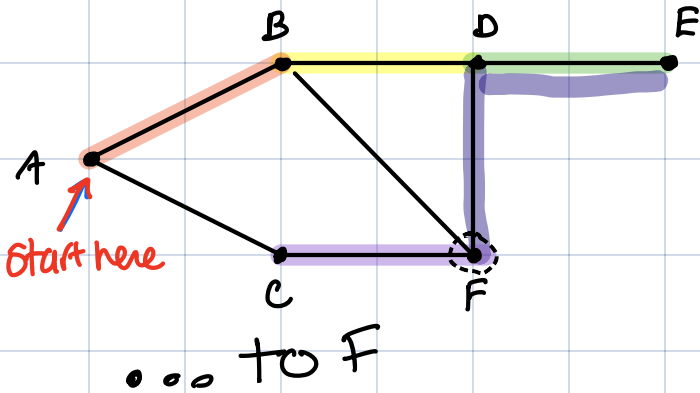


Visited

A, B, D, E, F, C

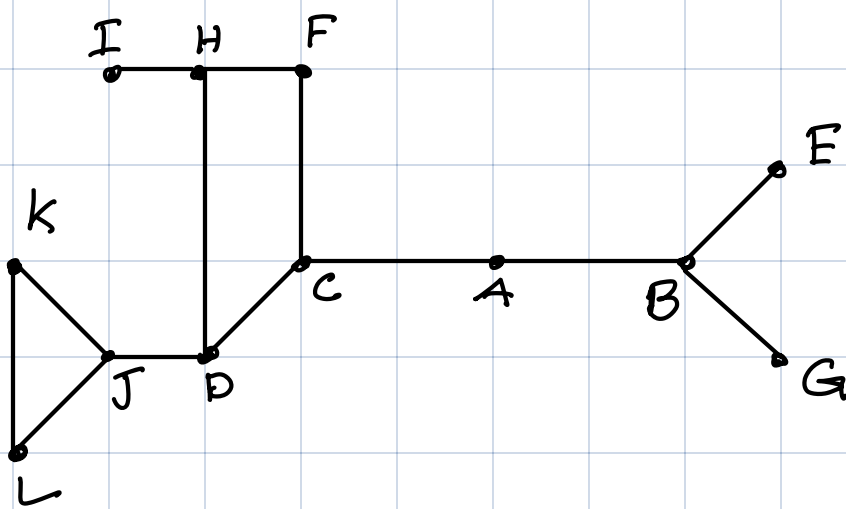
We may be done but Computer isn't!

One neighbor, A but visited, so backup...



Then the computer is done.

Exercise



1) DFS starting at A (order visit nodes in starting at A)
A B E G C D H F I J K L

2) DFS starting at H
H D C A B E G F J K L I

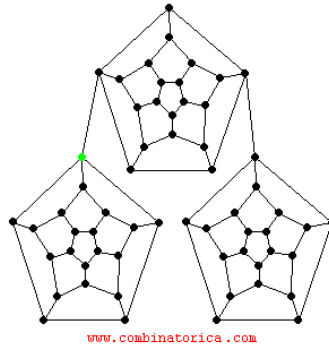
3) DFS starting at G
G B A C D H F I J K L E

Breadth First Search (BFS)

(the cautious algorithm)

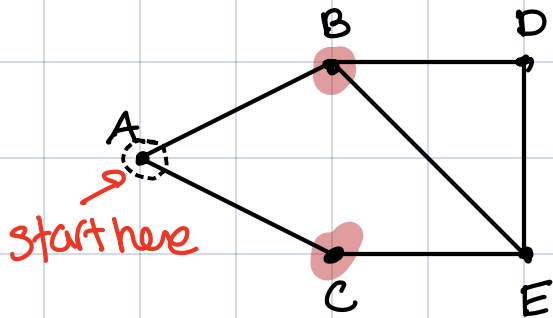
idea: explore all the close things before venturing out the next step

Breadth-First Search



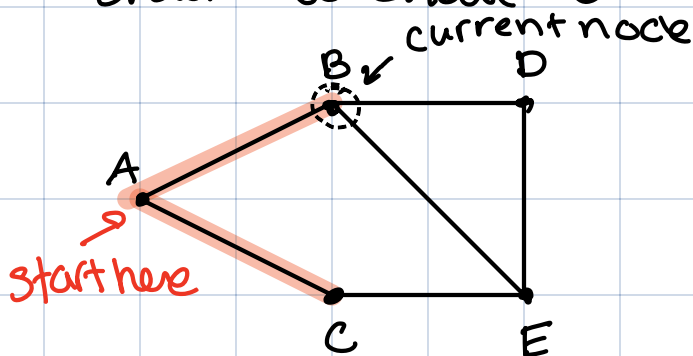
Formally: visit all vertices adjacent to the starting vertex then do the same from each of those vertices

Example



Visited
A

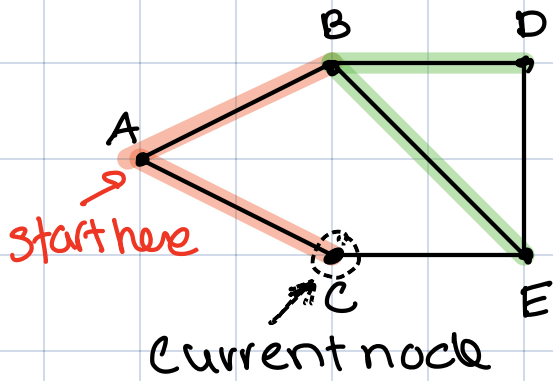
Neighbors are B, C, choose to visit in alphabetic order: so choose B.



Visited
A B C
↑

B has neighbors D, C, E. But have already

visited C, so add D, E

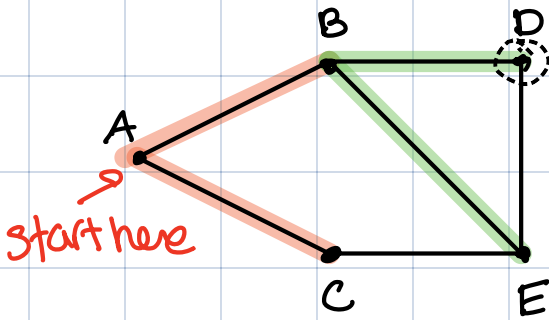


Visited

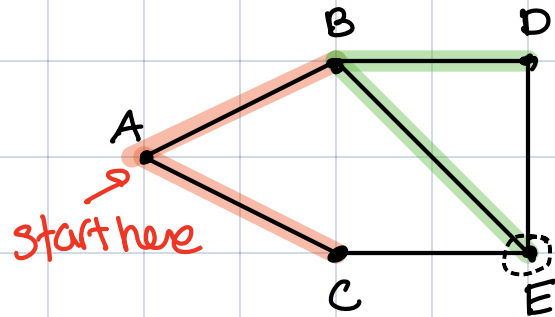
A B C D E
 ↑

we may be done but
the computer is not!

C has no unvisited neighbors so continue to...



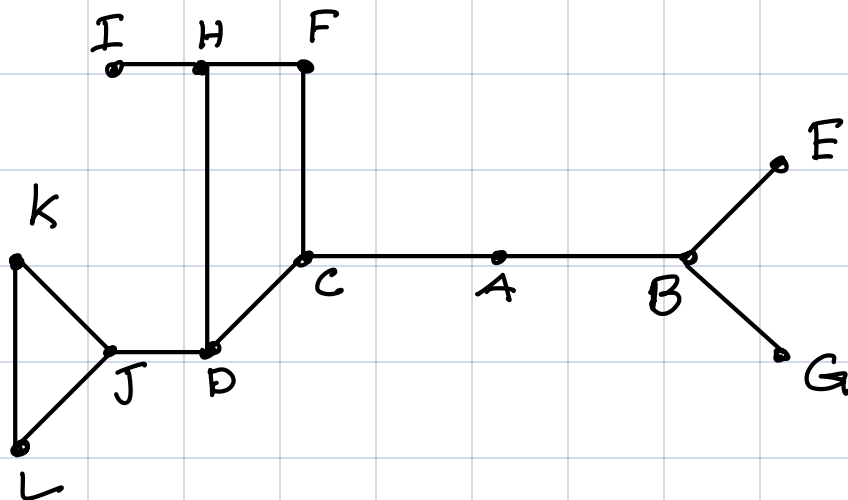
... to D



... to E

Now computer is done!

Exercise



1) BFS starting at A

A B C E G D F H J I K L
* * * * * * * * * *

2) BFS starting at H

H D F I C J A K L B E G
* * * * * * * * * *

3) BFS starting at G

G B A E C D F H J I K L
* * * * * * * * * *

Why BFS/DFS (again)

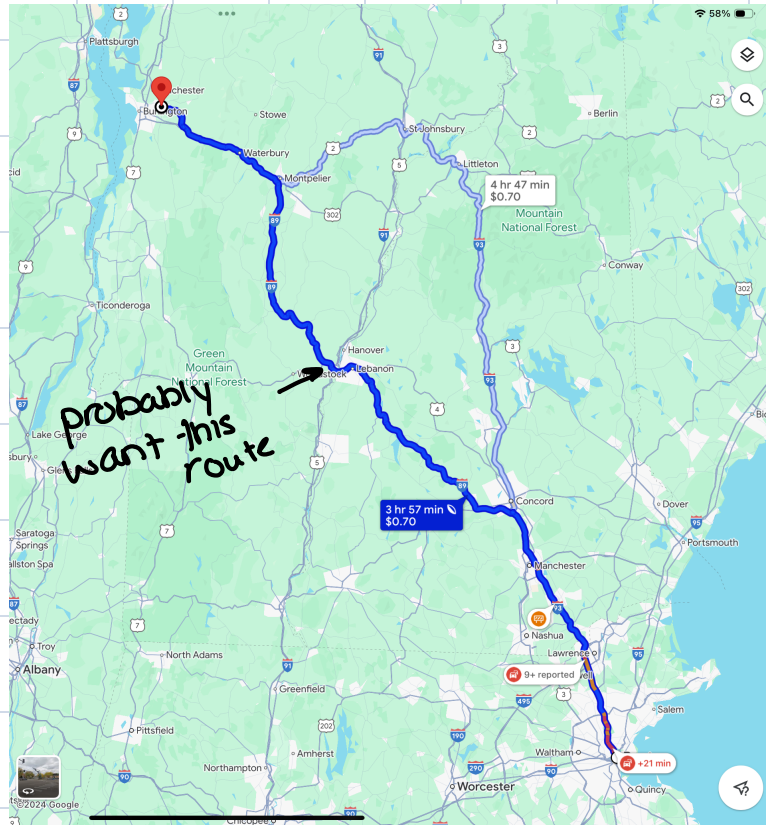
→ DFS/BFS gives you a connected subgraph
(where can I get by taking only one
airline)

→ DFS can detect cycles in a graph
(we bump into an already visited
neighbor)

→ BFS orders nodes by how far away they
are (1-hop, 2-hops, etc)

→ They can be written recursively ☺

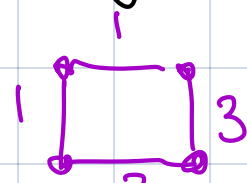
Now a different type of problem: how to get to point A to point B the fastest/cheapest/etc.



Known as "^{cheapest} ~~shortest~~ path" problem - can be solved using Dijkstra's algorithm

(cheapest path isn't always the one w/ least # of edges)

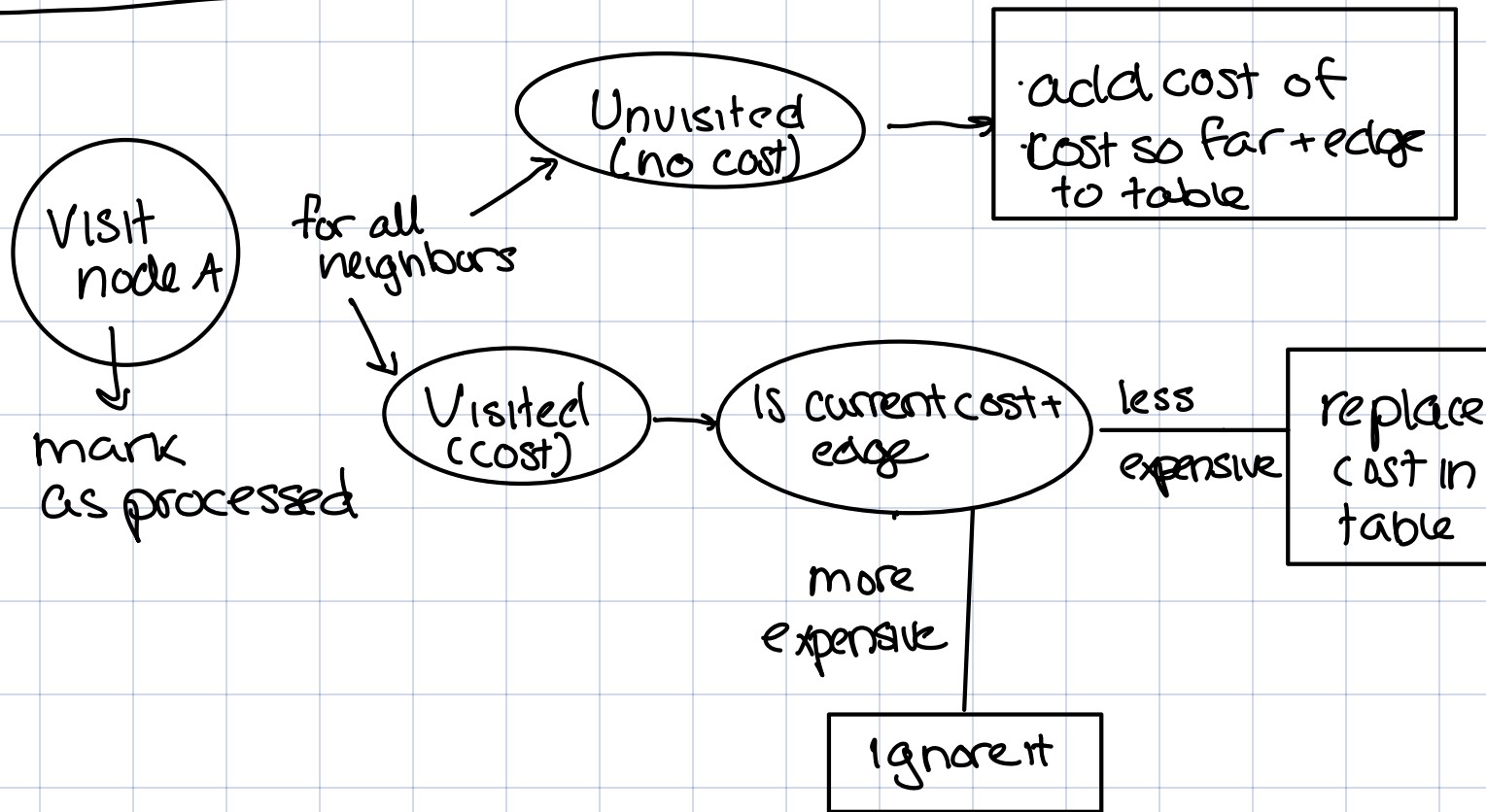
Formally: - weighted graph w/ positive edges
- what is least cost from A to B if we add up the edges along the path



Approach: keep track of cheapest path so far
in table

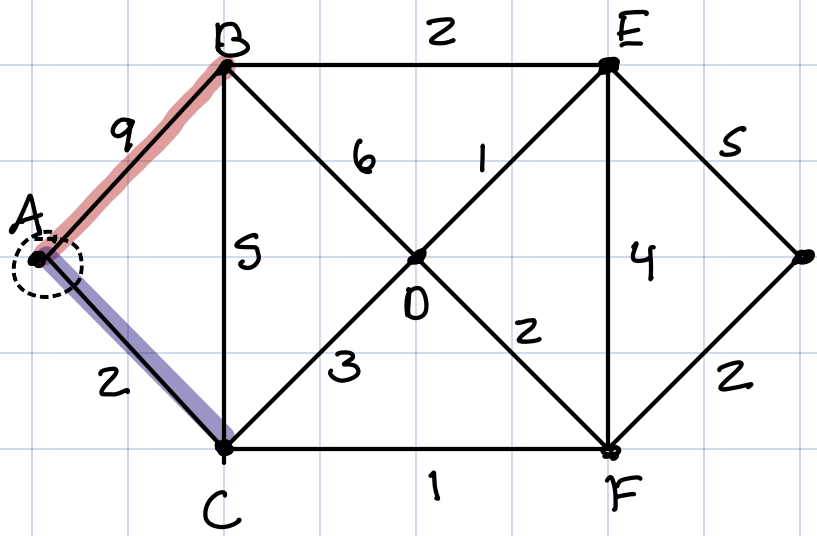
Warning flow chart below is for reference
will build up intuition w/ example

For reference



Then go to next node not processed
yet that has minimum cost

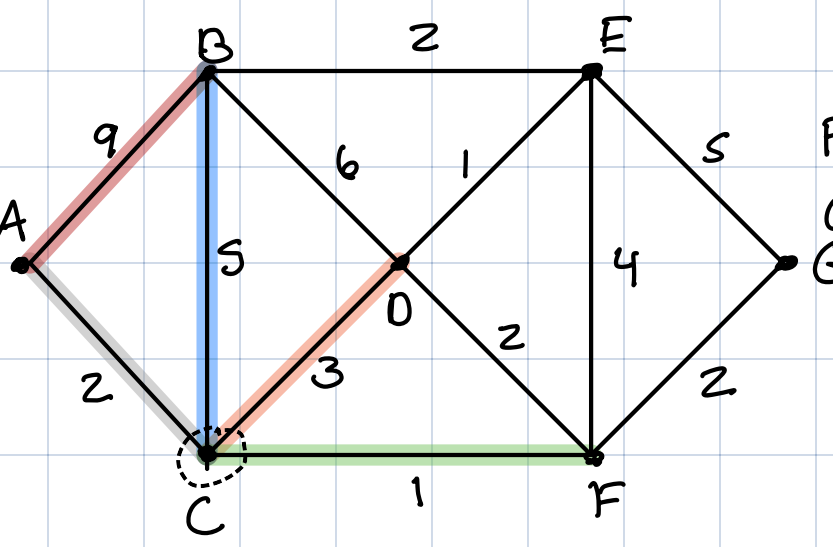
Cheapest path from A to G:



	A	B	C	D	E	F	G
Proc	X						
Cost	0	9	2				

1. Mark A as processed
2. Examine all neighbors
no cost? add to table

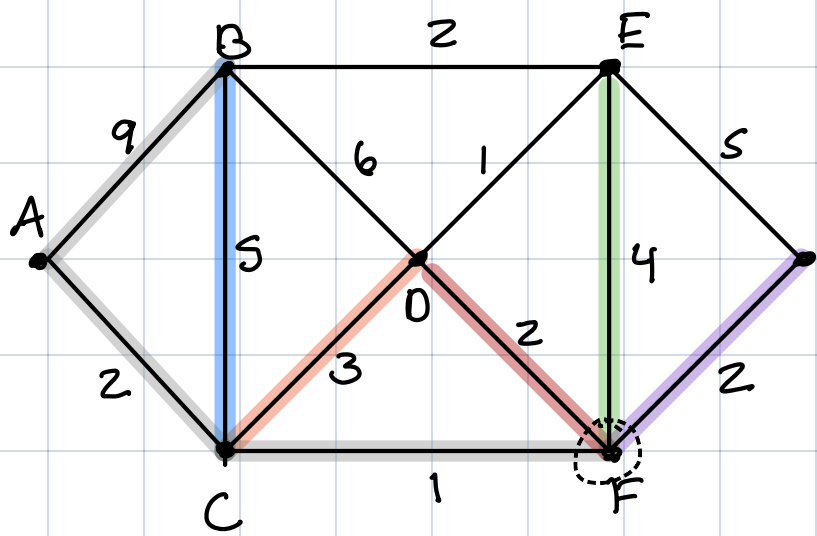
Next vertex is
 unvisited
 lowest cost
 so C



	A	B	C	D	E	F	G
Proc	X		X				
Cost	0	9	2	5		3	

1. Mark vertex as processed
2. Examine all neighbors
no cost? add to table
cost? lower, add to table
higher, ignore

Cost for D is $2+3$
 Cost for F is $2+1$
 Cost for B is $2+5$
 $7 < 9$, update
 Next vertex is F

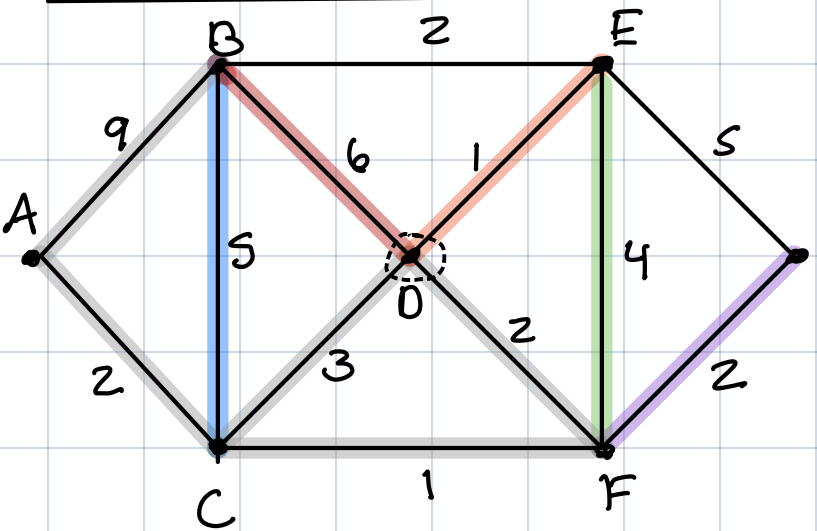


	A	B	C	D	E	F	G
Proc	X		X			X	
Cost	0	7	2	5	7	3	5

Cost for D is $3+2$
 $5=5$, ignore
 Cost for E is $3+4=7$
 Cost for G is $3+2$

1. Mark vertex as processed
2. Examine all neighbors
 no cost? add to table
 cost? lower, add to table
 higher, ignore
 or =

next vertex is D
 D & G have same value but we go alphabetically

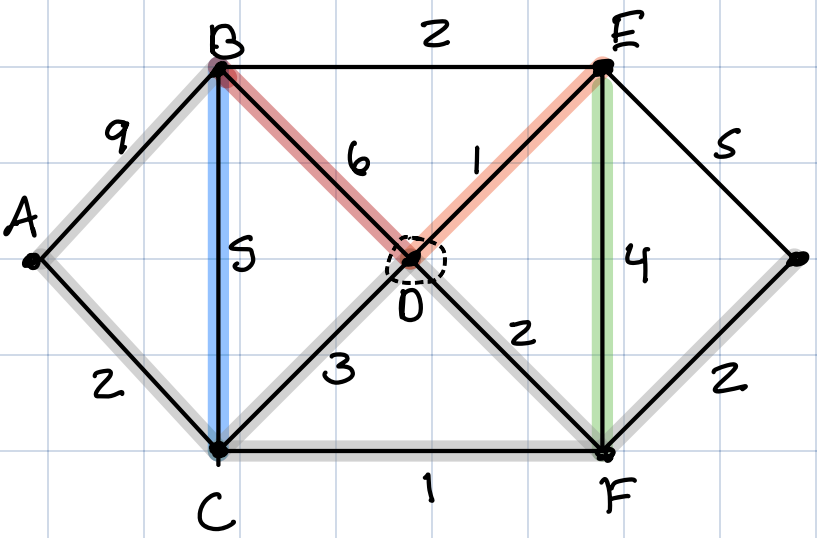


	A	B	C	D	E	F	G
Proc	X		X	X		X	
Cost	0	7	2	5	7 ⁶	3	5

Cost for B is $5+6$,
 Cost for E is $5+1$
 $6 < 7$

1. Mark vertex as processed
2. Examine all neighbors
 no cost? add to table
 cost? lower, add to table
 higher, ignore
 or =

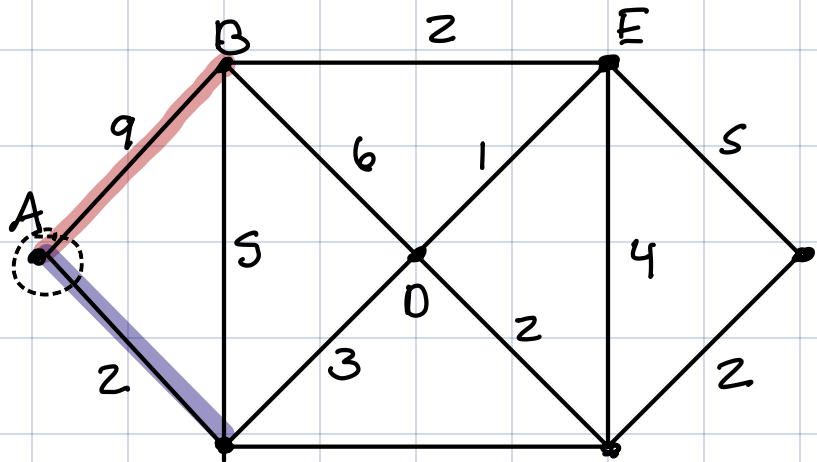
Next node is G!



	A	B	C	D	E	F	G
Proc	X		X	X		X	X
Cost	0	7	2	5	6	3	5

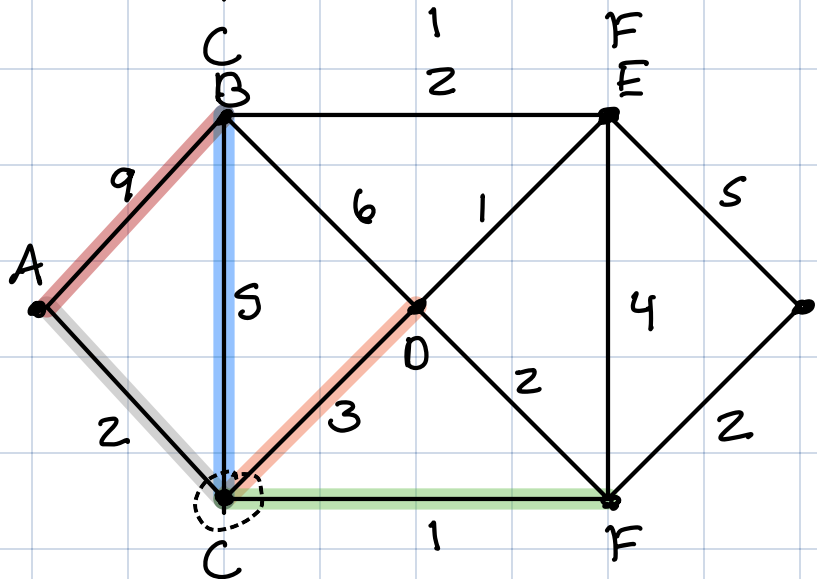
1. Mark vertex as processed
2. Examine all neighbors
 no cost? add to table
 cost? lower, add to table
 higher, ignore
 or =

We now know our shortest path to G cost 5, but what is that path?

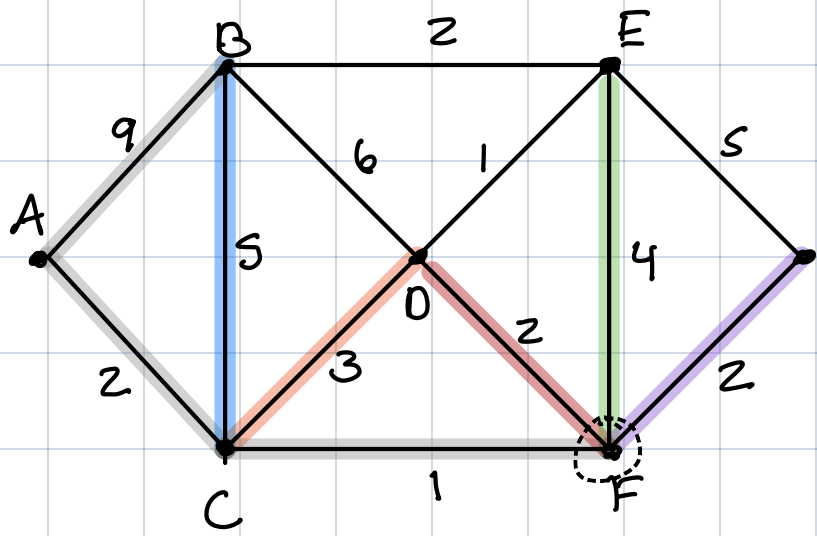


	A	B	C	D	E	F	G
Proc	X						
Cost	0	A:9	A:2				

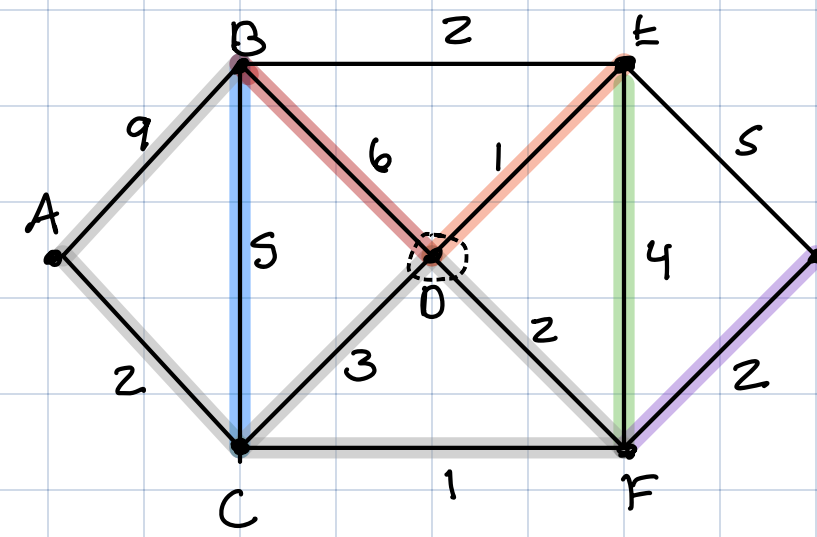
Keep track of proceeding vertex



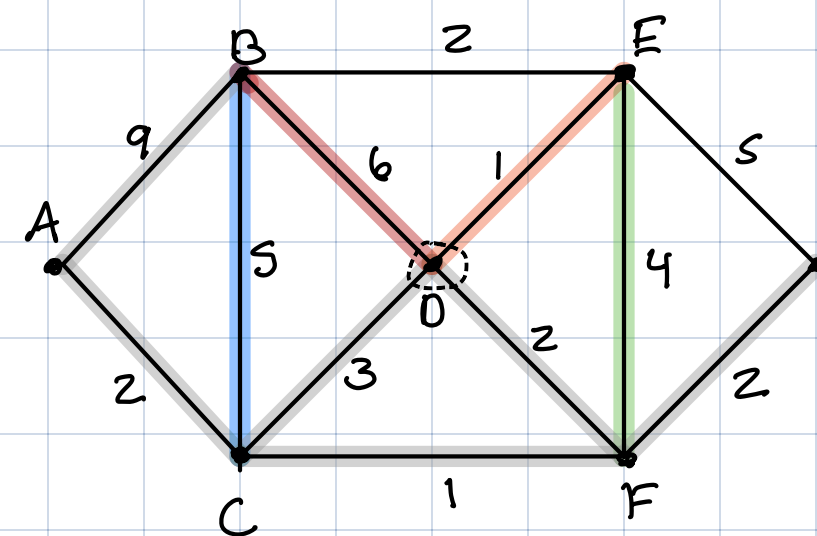
	A	B	C	D	E	F	G
Proc	X		X				
Cost	0	C:7	A:2	C:5		C:3	



	A	B	C	D	E	F	G
Proc	X		X			X	
Cost	0	C:7	A:2	C:5	F:7	C:3	F:5

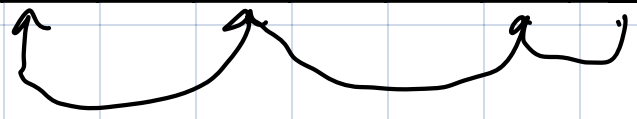


	A	B	C	D	E	F	G
Proc	X		X	X		X	
Cost	0	C:7	A:2	C:5	D:6	C:3	F:5



	A	B	C	D	E	F	G
Proc	X		X	X		X	X
Cost	0	7	A:2	C:5	6	C:3	F:5

	A	B	C	D	E	F	G
Proc	X		X	X		X	X
Cost	0	7	A:2	C:5	6	C:3	F:5



A → C → F → G

More backward through table to find path

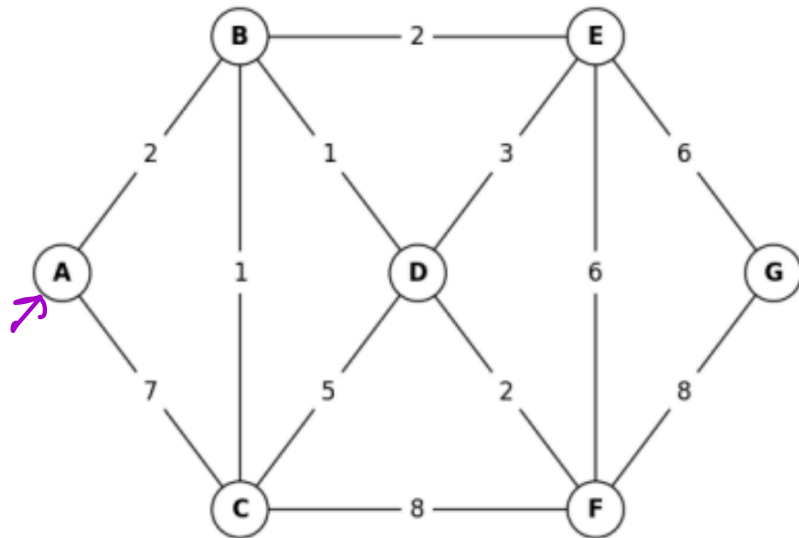
On Hw/Exam

iteration	node visited	A	B	C	D	E	F	G
0	A	start:0	A: 9	A: 2	none	none	none	none
1	C	start:0	C: 7	A: 2	C: 5	none	C: 3	none
2	F	start:0	C: 7	A: 2	C: 5	F: 7	C: 3	F: 5

The path with min weight is: $G \leftarrow F \leftarrow C \leftarrow A$

Example posted in detail on course web-page

Exercise



Shortest path from A to G

iteration	visited	A	B	C	D	E	F	G
0	A	0	A:2	A:7	-	-	-	-
1	B	0	A:2	B:3	B:3	B:4	-	-

