

## CS1800 day 3

### Admin:

- hw1 released today (due the following friday, as nearly all HWs are)
- tutoring groups

### Content:

- Two's complement (system to represent negative binary numbers)
- Overflow
- Floating point (system to represent non-whole numbers) (if time)

Whats the difference between operating in base-b and operating in base-b on a computer?

Computers store all values with the same number of bits

why? quicker / easier

Assume: a computer is using a 3-bit representation of values. How does it compute & store the following?

$$\begin{array}{r} (111)_2 \\ 7 \end{array} + \begin{array}{r} (001)_2 \\ 1 \end{array} = \begin{array}{r} (1000)_2 \\ 8 \end{array}$$

ONLY THESE THREE BITS CAN BE STORED

THIS IS DISCARDED

For today: assume we're working with values on a computer

- all values are N-digits  
(you'll be given this info in problem statement)
- discard the most significant (left-most) digits if needed  
(as shown in green on last slide)



LARGEST PLACE VALUE  
MAGNITUDE

## Number Systems:

Currently we're missing:

- negative values (e.g. -43)
- non-whole values (e.g. 321.12358)

Number systems:

### - Unsigned Integers:

can represent whole, non-negative numbers

everything we've done so far are unsigned integers (we just didn't cover name until now)

e.g.  $(110)_2 = 6$

### - Two's Complement:

can represent whole (potentially negative) numbers

(will study today)

### - Floating Point Values:

non whole-numbers

(will study today if time)

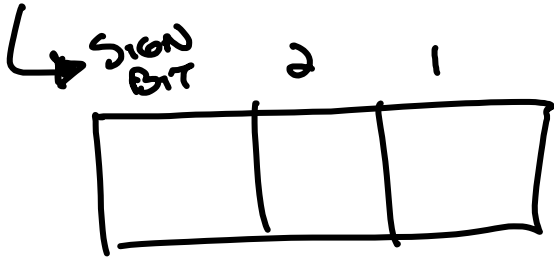
Sign bit\*:

A not-so-great number system for negative values

3 BIT "SIGN BIT"

1 IF NEGATIVE

0 OTHERWISE



$$(001)_2 = 1$$

$$(000)_2 = 0$$

$$(111)_2 = -3$$

$$(101)_2 = -1$$

# SIGN BIT: PROBLEMS

NO UNIQUE ZERO

$$(000)_2 = 0$$

$$(100)_2 = -0$$



OUR OPERATIONS YIELD INCORRECT RESULTS

$$(111)_2 + (001)_2 = (\text{DISCARD } X000)_2$$

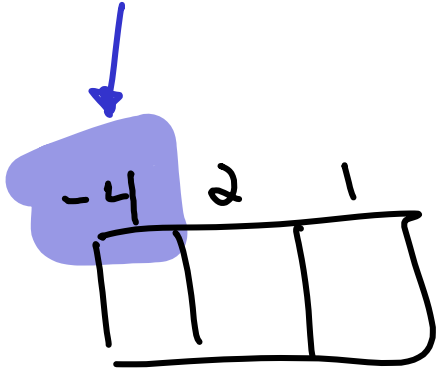
$$-3 + 1 = 0$$

0	2	1
1	1	0

$$\begin{array}{r} 111 \\ + 001 \\ \hline \end{array}$$

# Two's complement: A better way to store negative numbers

Big idea: the most significant (biggest) place value is negative, all others are positive



Example: 3-bit two's complement

-4	2	1	
1	0	0	-4
1	0	1	-3
1	1	0	-2
1	1	1	-1
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3

$$\begin{array}{r} 3+3 \\ + \quad 011 \\ \hline 110 \end{array}$$

-2

# TWO'S COMPLEMENT: PROBLEMS SOLVED

UNIQUE  
ZERO

$$(000)_2 = 0$$



OUR OPERATIONS YIELD CORRECT<sup>★</sup>  
RESULTS

$$(111)_2 + (001)_2 = (\overset{\text{DISCARD}}{X}000)_2$$

$$-1 + 1 = 0$$



★ Assumes that correct result may be represented (more later)

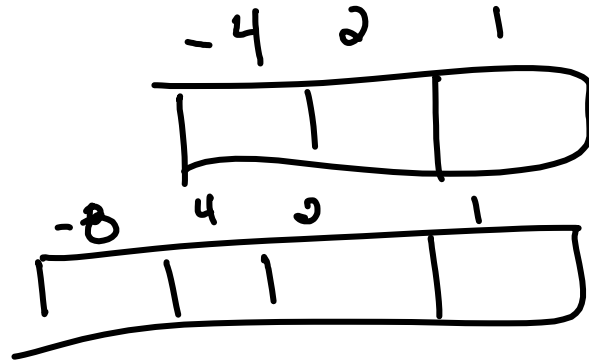


## In Class Activity:

If possible, convert each of the following values to the given number system. If not possible, justify why.

(Use guess-and-check as needed, a reliable decimal-to-2's-complement method coming shortly)

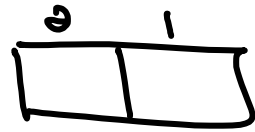
- 0 unsigned 2 bit integer
- 2 unsigned 3 bit integer
- 0 3 bit 2's complement
- 4 3 bit 2's complement
- 4 4 bit 2's complement
- 5 4 bit 2's complement
- 10 4 bit 2's complement
- 3 4 bit 2's complement



(++) What does the 2's complement idea look like in a base which isn't 2? Does it also have the properties we love so much in binary (unique zero, addition operations still work)?

0 unsigned 2 bit integer

$(00)_2$



-2 unsigned 3 bit integer: impossible!

- unsigned number systems only represent non-negative values

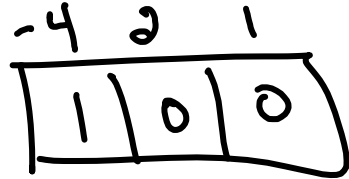
- each place value is positive, we may not represent a negative value

0 3 bit 2's complement

$(000)_2$

-4 3 bit 2's complement

$(100)_2$



-4 4 bit 2's complement

$$(1100)_2$$

5 4 bit 2's complement

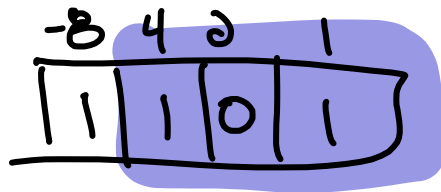
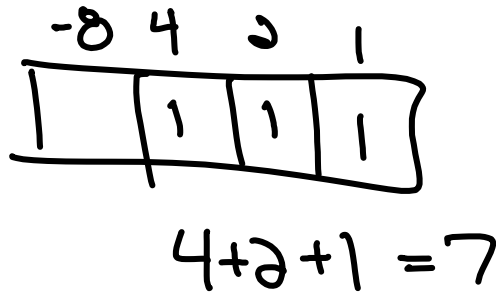
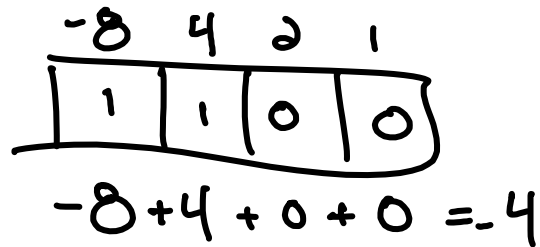
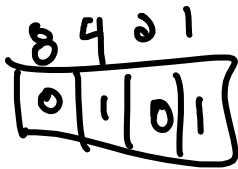
$$(0101)_2$$

10 4 bit 2's complement

10 is impossible! biggest value in  
4 bit 2's comp is  $(0111)_2 = 4 + 2 + 1 = 7$

-3 4 bit 2's complement

$$\begin{aligned} -8 + x &= -3 \\ x &= 5 \end{aligned}$$



R

A hand-drawn grid with 4 columns and 4 rows. The top-left cell contains a circled '0'. The grid is drawn with slightly curved lines, suggesting it might be a representation of a page or a specific type of grid.

G

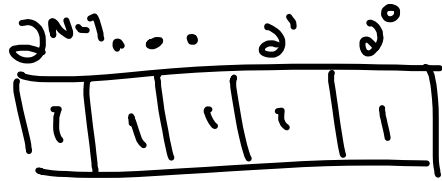
A hand-drawn grid with 4 columns and 4 rows. The top-left cell contains the number '255'. The grid is drawn with slightly curved lines, suggesting it might be a representation of a page or a specific type of grid.

B

A hand-drawn grid with 4 columns and 4 rows. The grid is drawn with slightly curved lines, suggesting it might be a representation of a page or a specific type of grid.

# What values can we represent with N bits?

Unsigned Integers



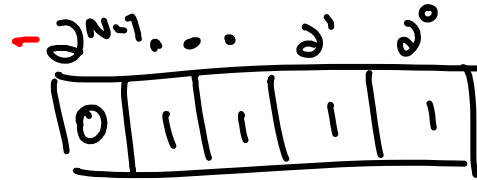
SMALLEST VALUE

$$0$$

LARGEST VALUE

$$2^N - 1$$

Two's Complement



SMALLEST VALUE

$$-2^{N-1}$$

LARGEST VALUE

$$2^{N-1} - 1$$

$$\begin{array}{cccccc}
 16 & 8 & 4 & 2 & 1 & \\
 \hline
 & 1 & 1 & 1 & 1 & \\
 \hline
 4 & 3 & 2 & 1 & 0 & 
 \end{array}$$

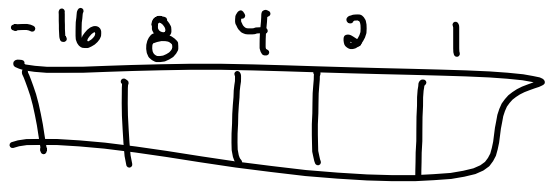
$$(1 + X)_{10} = (11111)_2 + (1)_2$$

$$= (100000)_2$$

$$= 2^5$$

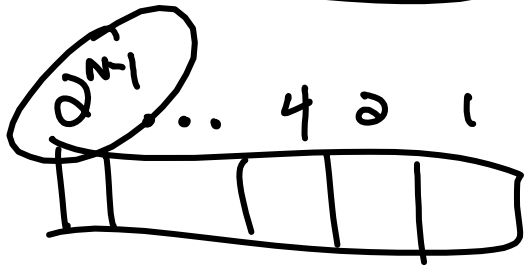
$$\begin{array}{r}
 11111 \\
 + 00001 \\
 \hline
 100000 \\
 \hline
 \end{array}$$

$$X = 2^5 - 1$$



$$N=5$$

$$-2^{5-1} = -16$$

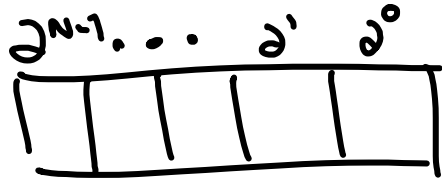


16	8	4	2	1
0	0	0	0	0



# What values can we represent with N bits? (representability)

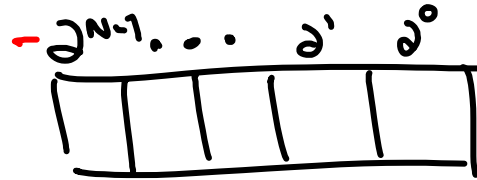
Unsigned Integers



SMALLEST VALUE  
0

LARGEST VALUE  
 $2^N - 1$

Two's Complement



SMALLEST VALUE  
 $-2^{N-1}$

LARGEST VALUE  
 $2^{N-1} - 1$

We can represent all whole values from smallest to largest (including smallest & largest)  
(we won't justify this)

# Overflow

$$\begin{array}{|c|c|c|} \hline -4 & 2 & 1 \\ \hline 1 & 1 & 1 \\ \hline \end{array} = -4 + 2 + 1 = -1$$

$$\begin{array}{|c|c|c|} \hline -4 & 2 & 1 \\ \hline 0 & 0 & 1 \\ \hline \end{array} = 1$$

Overflow: the outcome of an operation can't be represented in the given number system

example from earlier in lesson:

$$(111)_2 + (001)_2 = (\text{X}000)_2$$

$-1 + 1 = 0$

DISCARD

7 + 1 = 8 as 3 bit values

overflow since 8 can't be represented as a 3-bit value



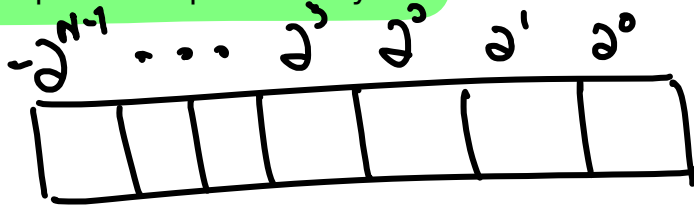
Common misconception:



There are times when we discard a bit but result is correct (no overflow occurs)

punchline: bit discard not relevant when determining overflow

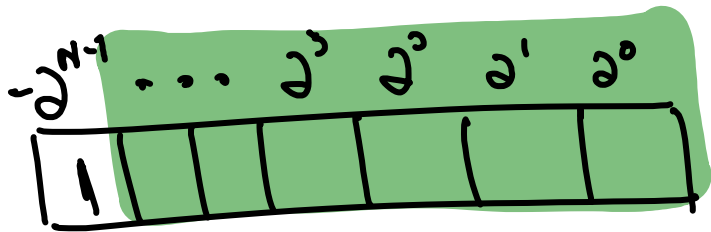
## Decimal to N-bit two's complement: preliminary



1. Validate that value can be represented as N-bit two's complement (see "representability")
2. If value is positive, its the same as N bit unsigned integer  
methods:
  - subtract largest power of two
  - Euclid's Division Algorithm

(if a non-negative number can be represented in n-bit 2's complement, its the same as n-bit unsigned)
3. If value is negative: see "x" method on next slide

## Decimal to N-bit two's complement: "x" method for negative representable values



DEFINE X AS  
VALUE OF  
LAST N-1 BITS  
(UNSIGNED)

$$\text{VALUE} = -2^{N-1} + X$$

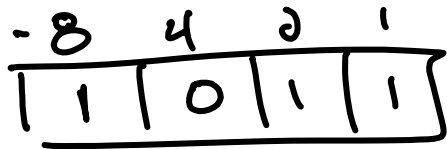
SO THAT

$$X = \text{VALUE} + 2^{N-1}$$

- Solve for X
- Represent X as N-1 bit unsigned int
- Append a leading 1 to indicate the  $-2^{N-1}$

# "X" method example

Represent ~~0~~<sup>-5</sup> in 4 bit 2's complement



$$-2^{N-1} = -2^3 = -8 \text{ SMALLEST}$$

$$2^{N-1} - 1 = 2^3 - 1 = 7 \text{ LARGEST}$$

$$-8 + X = -5$$

$$X = 3$$

## In Class Activity 2

If possible, express each of the following as a 6 bit two's complement value. Use the "x" method where possible.

$(111011)_2$

-5  
5  
~~3~~

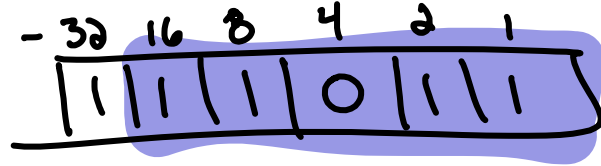
↑

IMPOSSIBLE

32 BIGGER THAN  
BIGGEST REPRESENTABLE

$$\text{SMALLEST} = -32$$

$$\text{LARGEST} = 31$$



$$-32 + x = -5$$

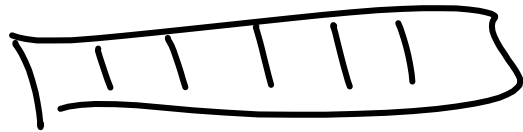
$$x = 27$$

$$= 16 + 11$$

$$= 16 + 8 + 3$$

$$= 16 + 8 + 2 + 1$$

5 IN 6-BIT 2'S COMP



$(000101)_2$

(floating point if time)



## Floating Point: Representing non-whole values

To express 12.345, rewrite it as:

$$12.345 = \underbrace{12345}_{\text{significand}} \times \underbrace{10^{-3}}_{\text{base}}^{\text{exponent}}$$

ALSO KNOWN AS  
MANTISSA

big idea: the significand and exponent will always be whole values and we can store those!

A few notes about the "base"

- isn't the same base the number system for significand & exponent number system  
(you can use base 10, as shown, and still store significand & exponent in binary)
- no need to store floating point base per individual value

IF TIME! NUMPY DOCUMENTATION